

Advanced Data Structures

Lecture 07: Packed and Compressed Hash Tables

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New Topic: External Memory Hash Tables

- now hash tables
- first packed and compressed hash table
- presented in January '23 at ALENEX

Motivation

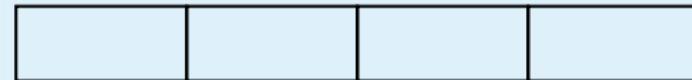
Setting

- static hash table for objects of variable size
- storing objects in external memory
- ideally retrieve objects in single I/O
- very small internal memory data structure

Objects of Variable Size



External Memory



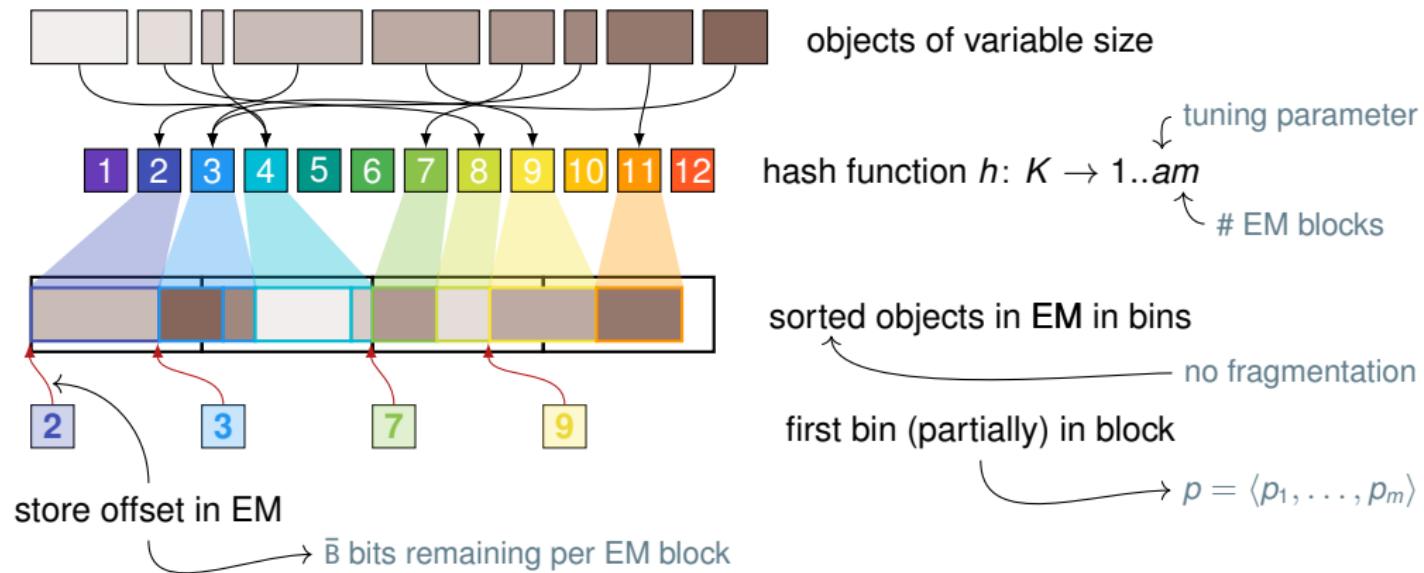
- only blocks of size B bits can be transferred
- one I/O per block transfer

Space-Efficient Object Stores from Literature

- objects of size 256 bytes
- blocks of size 4096 bytes
- internal space I_b (bits/block)
- (*) consecutive I/O

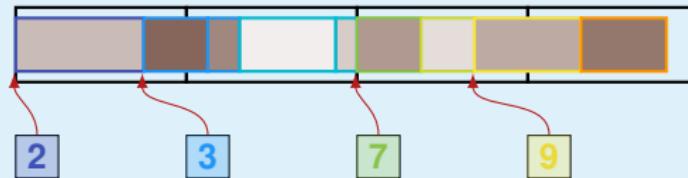
	Method	I_b	load factor	I/Os
fixed	Larson et al. [LR85]	96	<96 %	1
	SILT SortedStore [Lim+11]	51	100 %	1
	Linear Separator [Lar88]	8	85 %	1
	Separator [GL88; LK84]	6	98 %	1
	Robin Hood [Cel88]	3	99 %	1.3
	Ramakrishna et al. [RT89]	4	80 %	1
	Jensen, Pagh [JP08]	0	80 %	1.25
	Cuckoo [Aza+94; Pag03]	0	<100 %	2
	PaCHash , $a = 1$	2	100 %	2*
variable	PaCHash , $a = 8$	5	100 %	1.13*
	SILT LogStore [Lim+11]	832	100 %	1
	SkimpyStash [DSL11]	32	$\leq 98 \%$	8
	PaCHash , $a = 1$	2	99.95 %	2.06*
	PaCHash , $a = 8$	5	99.95 %	1.19*

PaCHash Overview



Finding Blocks

Query Algorithm



- $b_x = h(x)$
- find first i with $p_i \leq b_x$
- if $p_i = b_x$ let $i = i - 1$
- find first j with $p_j > b_x$
- return $i..(j - 1)$

Elias-Fano Coding

- given k monotonic increasing integers in $1..u$
 - store $\log k$ MSBs encoded in bit vector
 - store $\log(u/k)$ LSBs plain
 - $k(2 + \log(u/k)) + 1 + o(k)$ bits in total
- predecessor in $O(k)$ time

Lemma: Space with Elias-Fano Coding

When using Elias-Fano coding [Eli74; Fan71] to store p , the index needs $2 + \log a + o(1)$ bits of internal memory per block.

Predecessor Query in PaCHash Internal Memory

Lemma: Expected Predecessor Time

When using Elias-Fano coding to store p , the range of blocks containing the bin of an object x can be found in expected constant time.

Proof (Sketch)

- consider $\lceil \log m \rceil$ MSB
- let bin b_x have MSBs equal to u
- expected size $\mathbb{E}(Y_u)$ of all bins with MSB u that are $< b_x$ is

$$\begin{aligned}
 & \sum_{y \in S} |y| \cdot \mathbb{P}(h(y) \text{ w/ MSB } = u; h(y) < h(x)) \\
 & \leq \sum_{y \in S} |y| \cdot \mathbb{P}(h(y) \text{ w/ MSB } = u) \\
 & = \frac{1}{m} \sum_{y \in S} |y| = \frac{m\bar{B}}{m} = \bar{B}
 \end{aligned}$$

- number of entries to scan is $\mathbb{E}(Y_u)/\bar{B} = 1$

Loading Blocks from External Memory

Lemma: Additional Blocks Loaded

Retrieving an object x of size $|x|$ from a PaCHash data structure loads $\leq 1 + |x|/\bar{B} + 1/a$ consecutive blocks from the external memory in expectation.

Proof (Sketch)

- expected size of bin $b_x = h(x)$

$$\begin{aligned}\mathbb{E}(|b_x|) &= |x| + \sum_{y \in S, y \neq x} |y| \mathbb{P}(y \in b_x) \\ &\leq |x| + \sum_{y \in S} |y| \mathbb{P}(y \in b_x) \\ &= |x| + \sum_{y \in S} |y| \cdot \frac{1}{am} = |x| + \frac{\bar{B}}{a}\end{aligned}$$

Proof (Sketch, cnt.)

- expected number of blocks overlapped by b_x

$$\begin{aligned}\mathbb{E}(X) &= 1 + (\mathbb{E}(|b_x|) - 1)/\bar{B} \\ &= 1 + \frac{|x|}{\bar{B}} + \frac{1}{a} - 1/\bar{B}\end{aligned}$$

- $\mathbb{P}(\text{bin and block border align}) = 1/\bar{B}$

Experimental Evaluation

Hardware and Software

- Intel i7 11700 (base clock speed: 2.5 GHz)
- 1 TB Samsung 980 Pro NVMe SSD
- Ubuntu 21.10 (Kernel 5.13.0)
- io_uring for I/O operations
- GCC 11.2.0 (-O3 -march=native)
- $B = 4096$ bytes

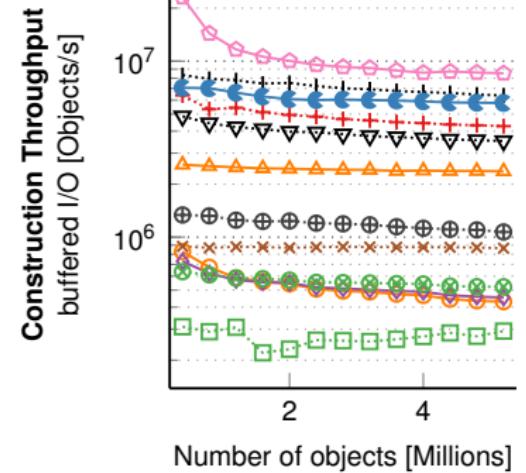
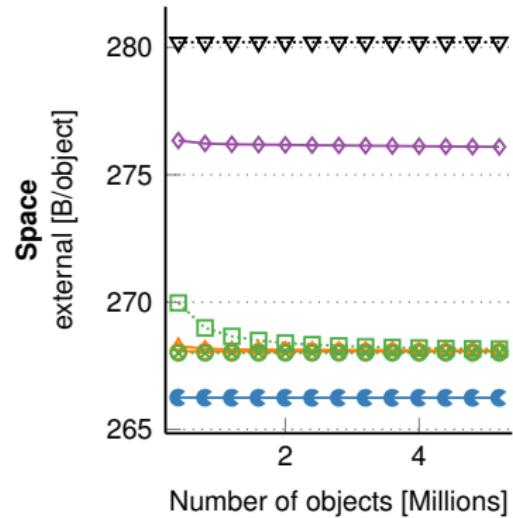
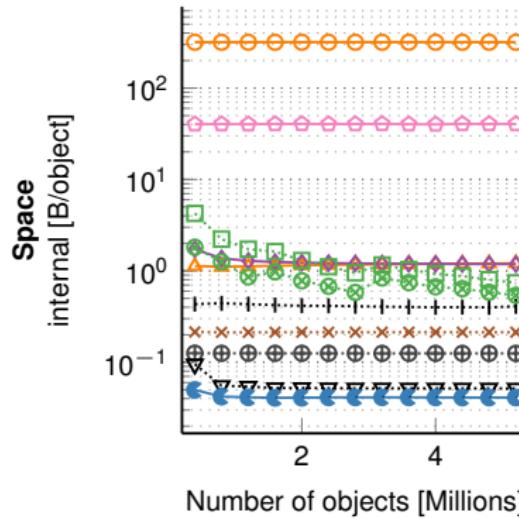
Objects

- here only **fixed size**
- more in the paper (very similar results)

Competitors

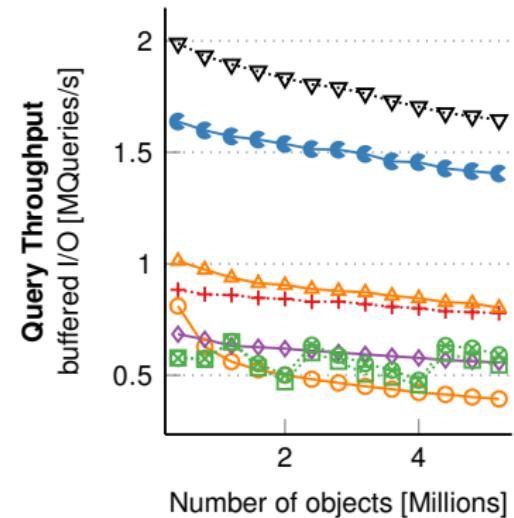
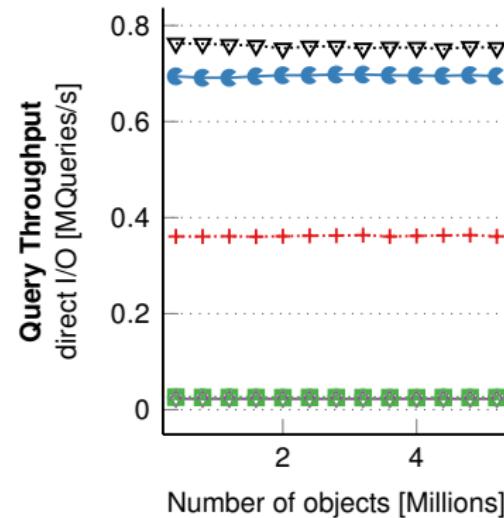
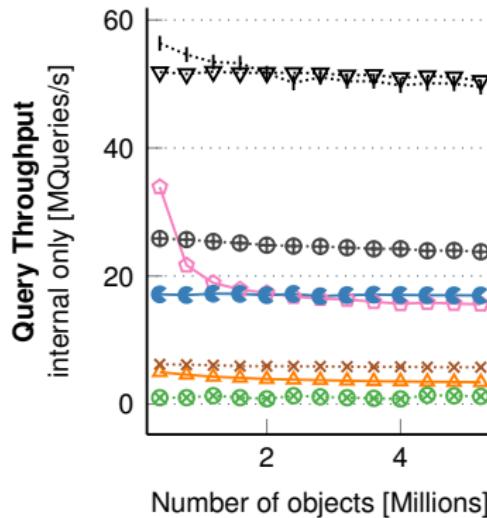
- LevelDB [[Goo21](#)]
- RocksDB [[Fac21](#)]
- SILT [[Lim+11](#)].
- `std::unordered_map`
- RecSplit [[EGV20](#)]
- CHD [[BBD09; CR+12](#)]
- PTHash [[PT21](#)]

Construction



- ⊕ CHD (16-perfect) [BBD09]
- + Cuckoo (here)
- △ LevelDB (Static part) [Goo21]
- LevelDB [Goo21]
- PTHash [PT21]
- RecSplit [EGV20]
- SILT (Static part) [Lim+11]
- SILT [Lim+11]
- ▽ Separator (here)
- ◊ std::unordered_map
- ◆ RocksDB [Fac21]

Queries



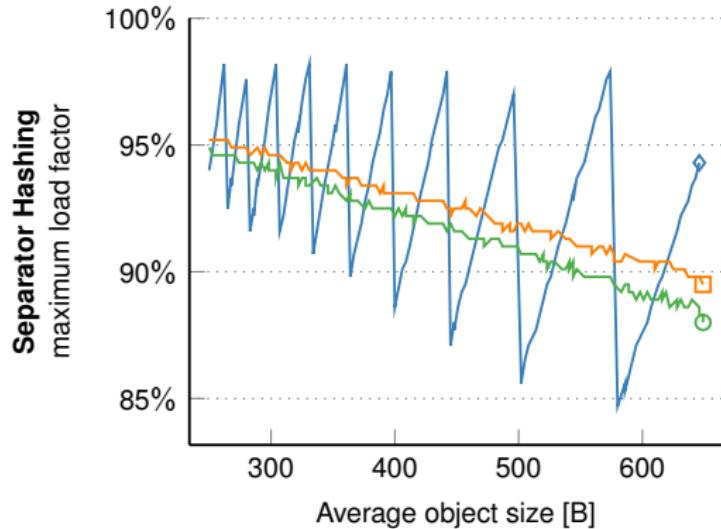
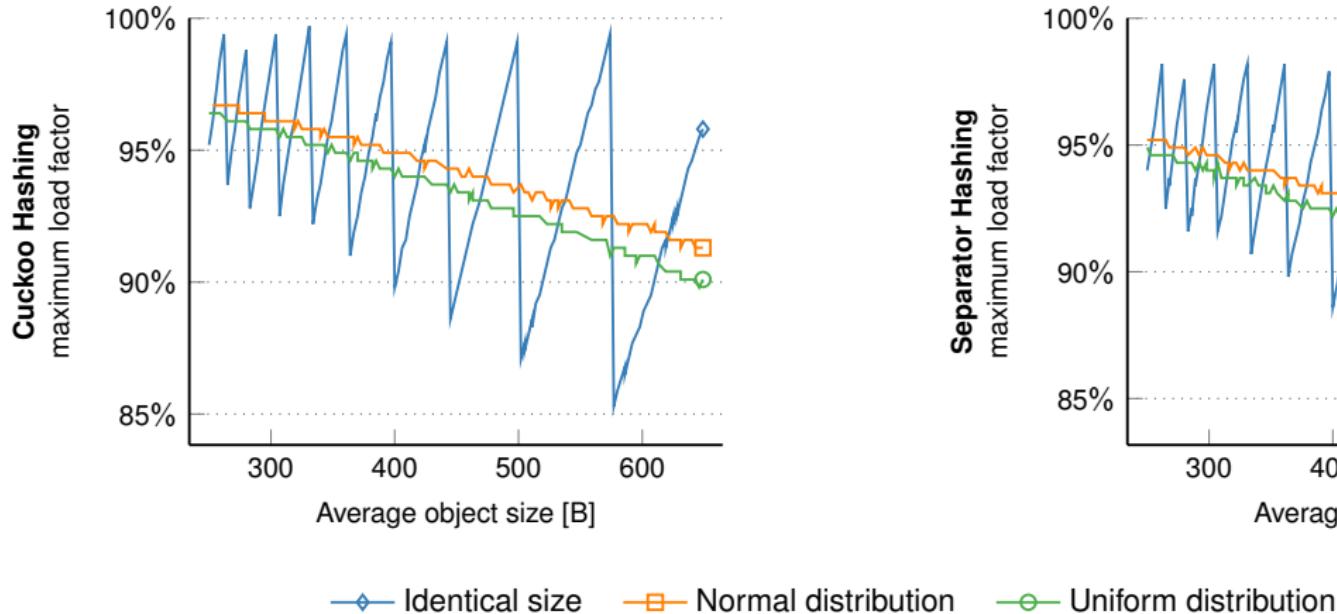
Legend:

- CHD (16-perfect) [BBD09]
- Cuckoo (here)
- LevelDB (Static part) [Goo21]

Legend:

- LevelDB [Goo21]
- PTHash [PT21]
- PaCHash (here)
- RecSplit [EGV20]
- RocksDB [Fac21]
- SILT (Static part) [Lim+11]
- SILT [Lim+11]
- Separator (here)
- std::unordered_map

Maximum Load Factor of Competitors



Alternative Internal Memory Data Structures

Lemma: Space with Succincter

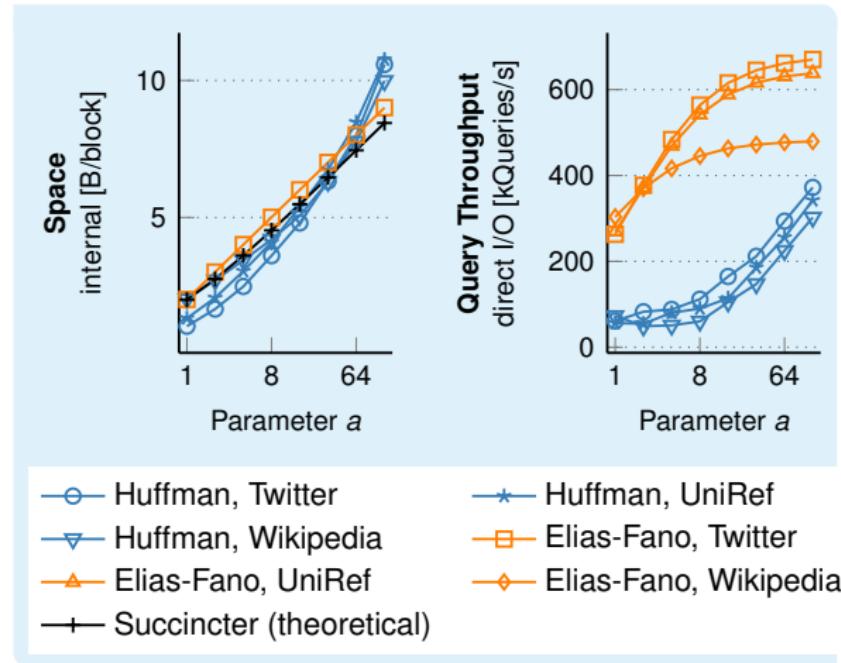
When using Succincter [Pat08] to store p , the index needs $1.44 + \log(a+1) + o(1)$ bits of internal memory per block.

Structure of Bit Vector

- runs of 0s and 1s
- sometimes additional 1s

Entropy Encoding

- encode positions directly
- compress bit vector using Huffman codes
- encode blocks of size 8, 16, 32, or 64



Conclusion and Outlook

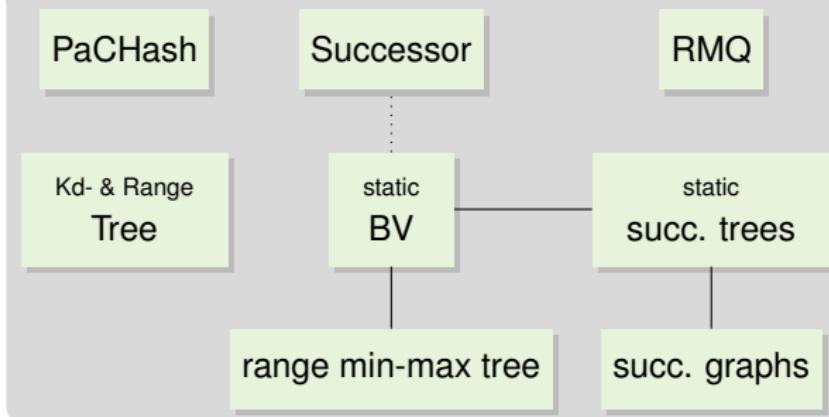
This Lecture

- PaCHash

Next Lecture

- more on hashing

Advanced Data Structures



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