Karlsruhe Institute of Technology

Advanced Data Structures

Lecture 10: Retroactive Data Structures (cnt.) and Minimal Perfect Hashing

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PINGO

https://pingo.scc.kit.edu/489786

Recap: Retroactive Data Structures

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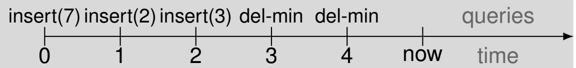
Operations

- $\texttt{INSERT}(t, \mathit{operation})$: insert operation at time t
- $\texttt{DELETE}(t)$: delete operation at time t
- $\texttt{QUERY}(t, \mathit{query})$: ask *query* at time t

- for a priority queue updates are
  - insert
  - delete-min
- time is integer ℹ for simplicity otherwise use order-maintenance data structure



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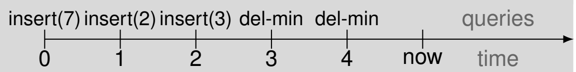

Definition: Partial Retroactivity

`QUERY` is only allowed for $t = \infty$ ℹ now

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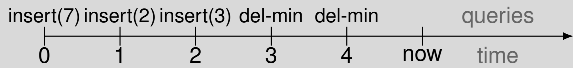

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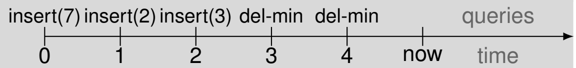

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`QUERY` is only allowed for $t = \infty$ ℹ now

Definition: Full Retroactivity

`QUERY` is allowed at any time t

Definition: Nonoblivious Retroactivity

`INSERT`, `DELETE`, and `QUERY` at any time t but also identify changed `QUERY` results

Priority Queues: Partial Retroactivity (1/6)

- priority queue with
  - insert
  - delete-min
- delete-min makes PQ non-commutative

Lemma: Partial Retroactive PQ

A priority queue can be partial retroactive with only $O(\log m)$ overhead per partially retroactive operation

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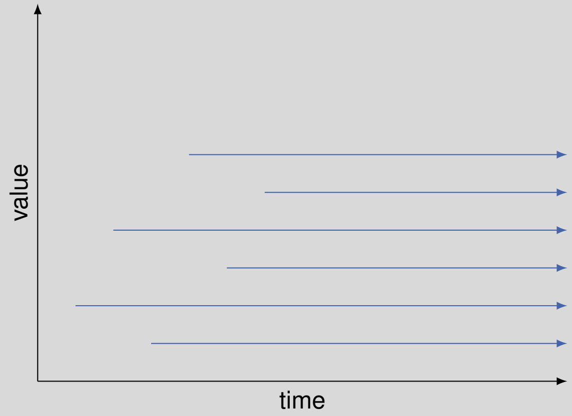

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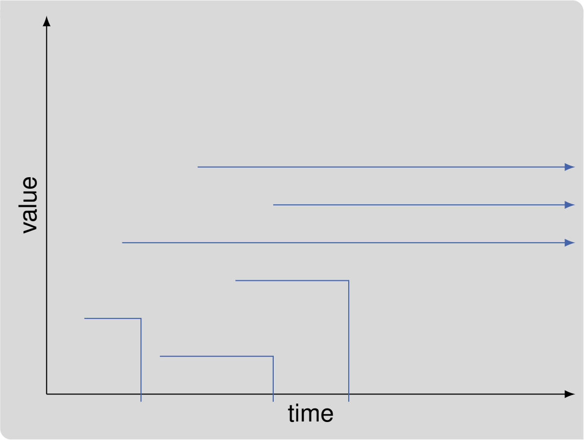

Priority Queues: Partial Retroactivity (2/6)

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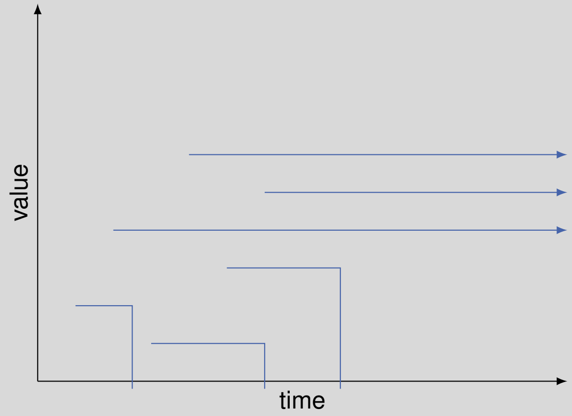
  - `INSERT(t,`insert`(i))`



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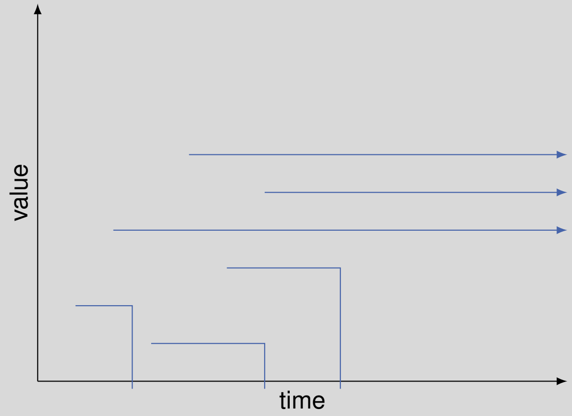

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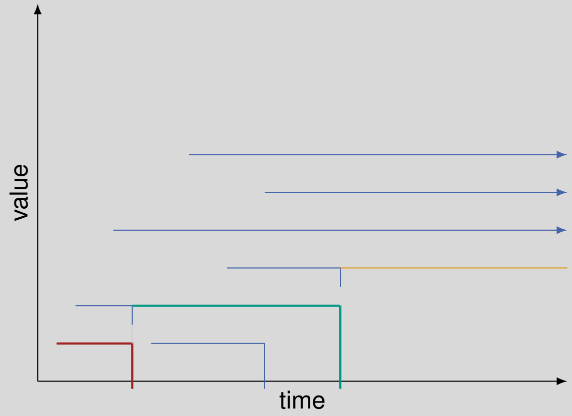

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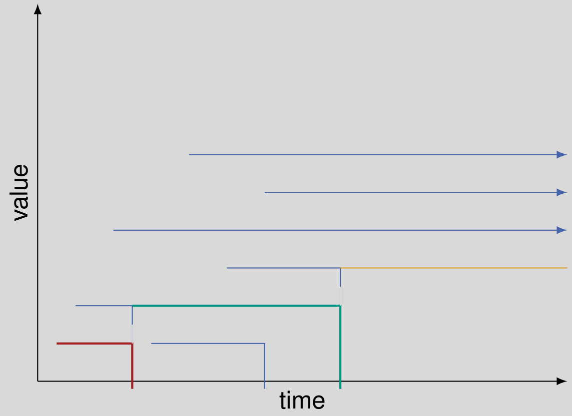

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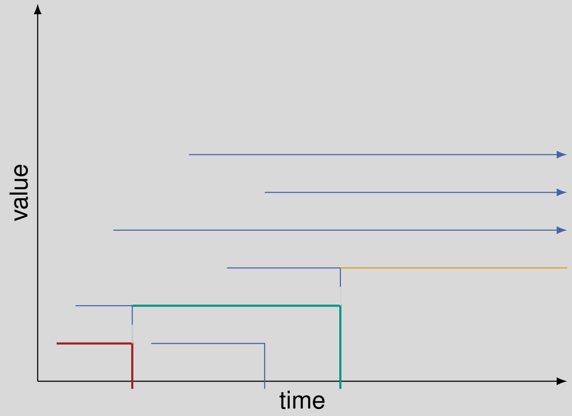
- insert deleted minimum right after deletion



Priority Queues: Partial Retroactivity (3/6)

- let Q_t be elements in PQ at time t

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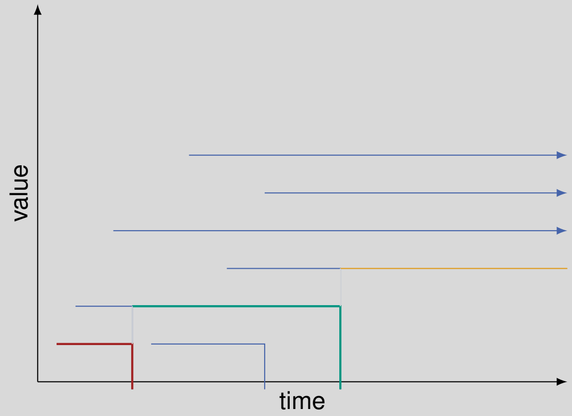
- what value inserts $\texttt{INSERT}(t, \text{insert}(v))$ in Q_∞
- values is $\max\{v, v' : v' \text{ deleted at time} \geq t\}$
- maintaining deleted elements is hard can change a lot



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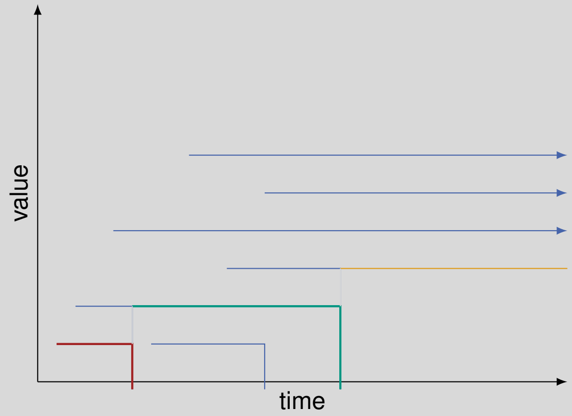
A time t' is a bridge if $Q_{t'} \subseteq Q_\infty$

- all elements present at t' are present at t_∞



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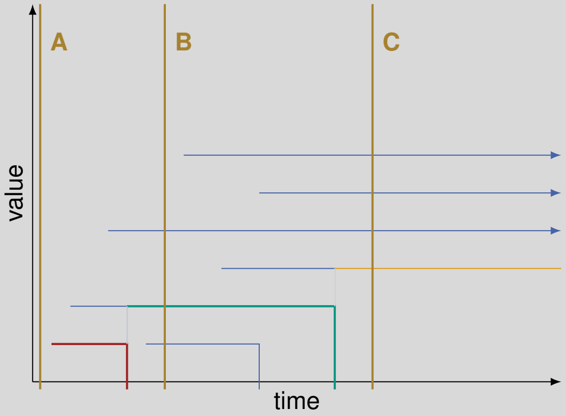

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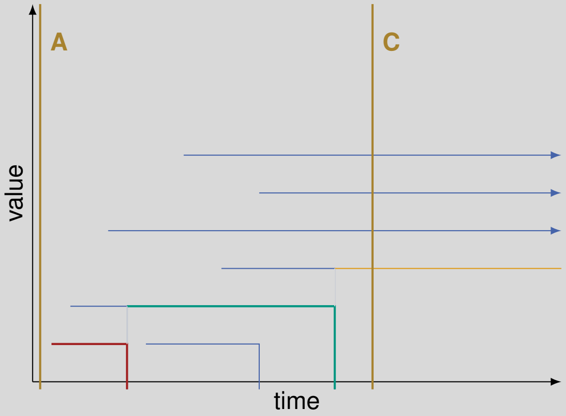

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Priority Queues: Partial Retroactivity (4/6)

Lemma: Deletions after Bridges

If time t' is closest bridge preceding time t, then

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Proof (Sketch)

- $\max\{v' \notin Q_\infty : v' \text{ inserted at time } \geq t'\} \in \{v' : v' \text{ deleted at time } \geq t\}$
  - if maximum value is deleted between t' and t
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  - if v' is deleted at some time $\geq t$
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- update all BBSTs in $O(\log n)$ time

Priority Queues: Partial Retroactivity (6/6)

Lemma: Partial Retroactive PQ

A priority queue can be partial retroactive with only $O(\log m)$ overhead per partially retroactive operation

- requires three BBSTs
- updates need to update all BBSTs

Hashing (1/2)

- $h\colon \{0, \ldots, u-1\} \to \{0, \ldots, m-1\}$

- n objects
- from universe $U = \{0, \ldots, u-1\}$
- hash table of size m ℹ m close to n
- $m \ll u$

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Definition: Totally Random

- $\mathbb{P}[h(x) = t] = 1/m$
- independent of $h(y)$ for all $x \neq y \in U$
- requires $\Theta(u \log m)$ bits of space to store ℹ too big

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- $h(x) = (ax \mod u) \mod m$ for $0 < a < p$ and p being prime $> u$
- $h(x) = ax \gg (\log u - \log m)$ for m, u being powers of two

- Why is this family easier to store? **PINGO**

Hashing (2/2)

Definition: k-wise Independent

- choose h from family H with $\mathbb{P}[h(x_1) = t_1 \& \ldots \& h(x_k) = t_k] = O(1/m^k)$ for distinct $x_1, \ldots, x_k \in U$

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- $h(x) = ((\sum_{i=0}^{k-1} a_i x^i) \mod p) \mod m$ for $0 \leq a_i < p$ and $0 < a_{k-1} < p$

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Definition: Simple Tabulation Hashing

- view x as vector $x_1, \ldots, x_c$ of characters
- totally random hash table T_i for each character
- $h(x) = T_1(x_1) \texttt{ xor } \ldots \texttt{ xor } T_c(x_c)$

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- $O(cu^{1/c})$ space
- $O(c)$ time to compute
- 3-wise independent

Minimal Perfect Hashing

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Definition: Perfect Hash Function

- injective hash function
- maps n objects to m slots

- lower space bound for $m = (1 + \epsilon)n$ is

$$\log e - \epsilon \log \frac{1 + \epsilon}{\epsilon}$$

- for m close to n there are likely collisions

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Definition: Minimal Perfect Hash Function

- bijective hash function
- maps n objects to $m = n$ slots
- $h\colon N \to [0, n)$

- lower space bound as for PHF with $\epsilon = 0$:

$$\log e \approx 1.44$$

- no collisions

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- for m close to n there are likely collisions

Definition: Minimal Perfect Hash Function

- bijective hash function
- maps n objects to $m = n$ slots
- $h\colon N \to [0, n)$

- lower space bound as for PHF with $\epsilon = 0$:

$$\log e \approx 1.44$$

- no collisions

- can we make PHF to MPHF? **PINGO**

BDZ (RAM) Algorithm [BPZ13]

- for each object calculate three *potential* slots (h_0, h_1, and h_2)
- for each slot that contains only one object, remove the object from all its other slots
- one slot per object
- if that does not work use other hash functions
- use rank data structure to map slots to $[0, n)$

- example on the board

- 1.95 bits per object when $m = 1.23n$

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- how to check if hash function works
- interpret each slot as node in a hypergraph
- objects are edges
- if graph is peelable, we have a feasible mapping

Definition: Peelable

A hypergraph is peelable, if it is possible to obtain a graph without edges by iteratively taking away edges that contain a node with degree 1

- example on the board

Compress, Hash, and Displace [BBD09a]

- partition keys into buckets
- set $m = (1 + \epsilon)n$ ℹ $1.01n$
- sort partitions by size
- starting with largest bucket, find universal hash function mapping all keys to empty slots
- if key mapped to non-empty slot, try next hash function
- for each bucket store universal hash function
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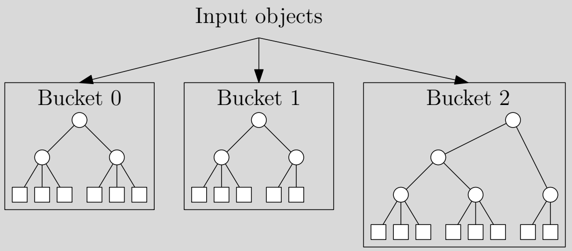
- example on the board

- can be used as PHF
- there are a lot of tricks w.r.t. bucket sizes and size distributions
- requires around 2.05 bits per object

RecSplit Overview [EGV20a]

- partition keys into buckets of size b
- for each bucket compute splitting trees
- split keys into smaller sets
- stop when sets have size ℓ



RecSplit Overview [EGV20a]

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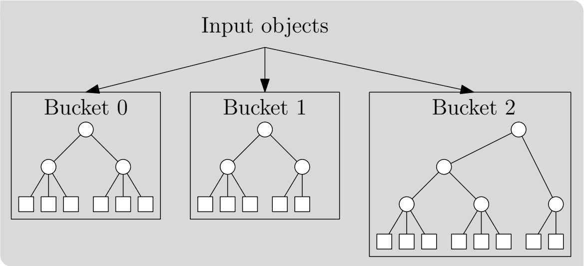
- upper aggregation levels have fanout 2
- lower two aggregation levels have fanout
  - $\max\{2, \lceil 0.35\ell + 0.55 \rceil\}$
  - $\max\{2, \lceil 0.21\ell + 0.9 \rceil\}$



RecSplit Overview [EGV20a]

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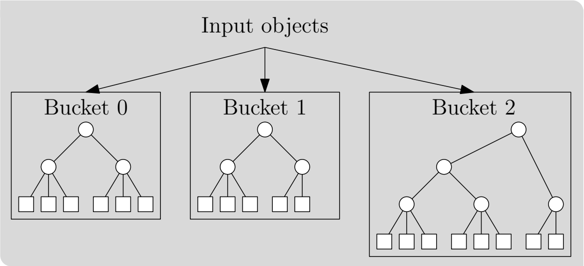
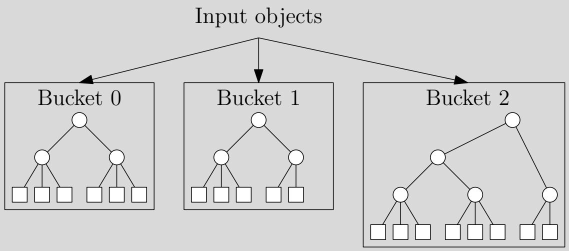
- last level is leaf level
- find bijections



RecSplit Splitting Tree

- tree structure is well defined
- store information for each node in preorder
- store hash function for each splitter
- encode function using Golomb-Rice



RecSplit Splitting Tree

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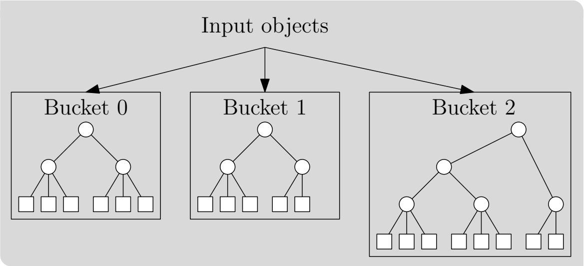
- encodings of splitting trees stored in one bit vector
- use Elias-Fano to store
  - size of buckets
  - starting position of bucket in bit vector



Golomb Encoding [Gol66]

Definition: Golomb Code

Given an integer $x > 0$ and a constant $b > 0$, the Golomb code consists of

- $q = \lfloor \frac{x}{b} \rfloor$
- $r = x - qb = x \,\%\, b$
- $c = \lceil \lg b \rceil$

with

$$(x)_{\mathrm{Gol}(b)} = (q)_1 (r)_2$$

where $(r)_2$ depends on its size

- $r < 2^{\lfloor \lg b \rfloor - 1}$: r requires $\lfloor \lg b \rfloor$ bits and starts with a 0
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- still variable length

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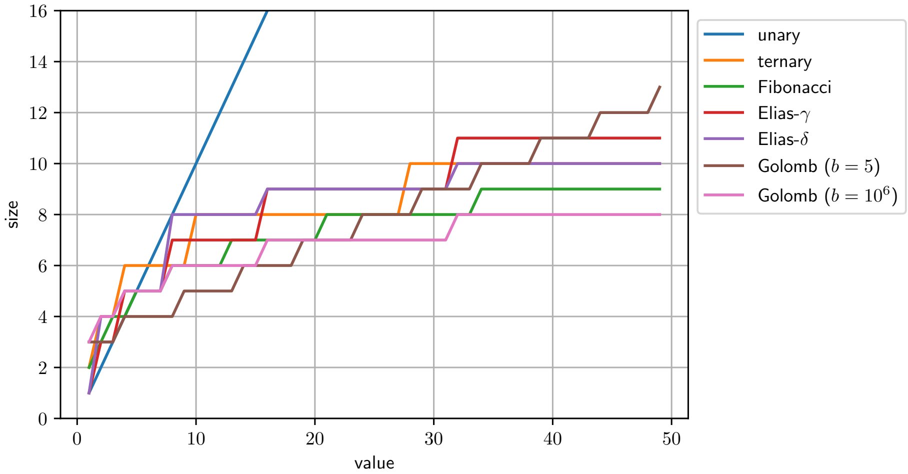
- b has to be fixed for all codes
- still variable length

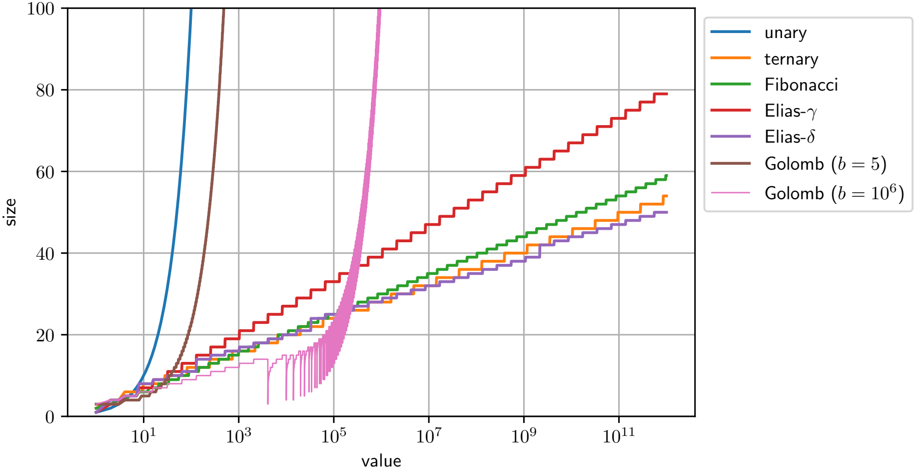
- Golomb-Rice is special case where r is power of two

- for $b = 5$, there are 4 remainders: $00, 01, 100, 101$, and 110
- $2^{\lfloor \lg 5 \rfloor - 1} = 2$
- $0, 1 < 2$: 00 and 01 require 2 bits
- $2, 3, 4 \geq 2$: require 3 bits and encode $0, 1, 2$ starting with 1

Comparison of Codes (1/2)

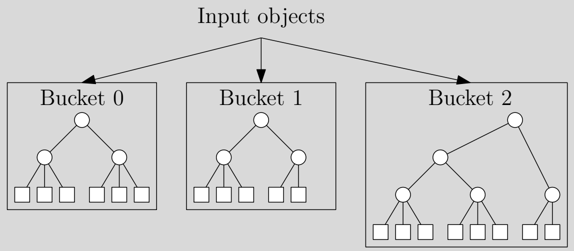


Comparison of Codes (2/2)



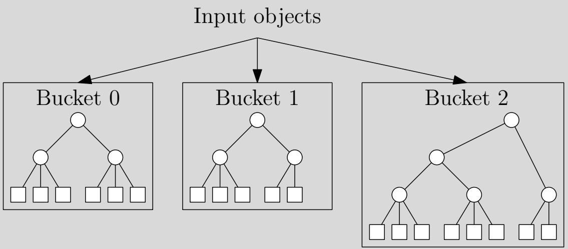
RecSplit Leaves

- find perfect hash function for keys in leaves
- test hash functions brute force
- use hash value modulo ℓ
- set bit in “bit vector” of length ℓ
- all bits set indicates bijection



RecSplit Queries

- find bucket
- follow splitting tree
- accumulate number of objects to the left
- use bijection in leaf
- result is sum of
  - objects in previous buckets
  - objects to the left in splitting tree
  - value of bijection



Parallel RecSplit

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- Dominik Bez, Florian Kurpicz, Hans-Peter Lehmann, and Peter Sanders. “High Performance Construction of RecSplit Based Minimal Perfect Hash Functions”. In: *ESA*. volume 274. LIPIcs. Schloss Dagstuhl - Leibniz-Zentrum für Informatik, 2023, 19:1–19:16. DOI: 10.4230/LIPICS.ESA.2023.19
- based on a Domink Bez’ Master’s thesis

- randomly distribute objects in leaf in two sets A and B
- hash objects in both set
- two “bit vectors”: cyclic shift one until all bits are set when `OR`ed
- store hash function *and* rotation

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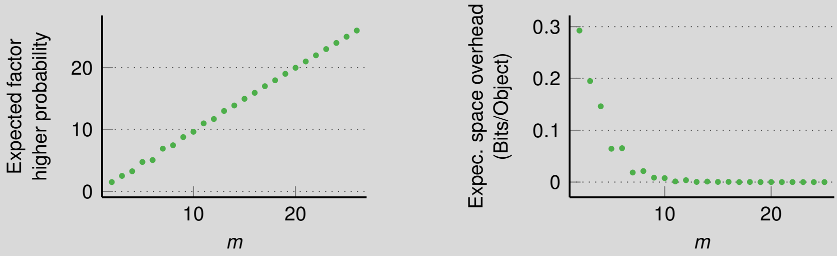

- randomly distribute objects in leaf in two sets A and B
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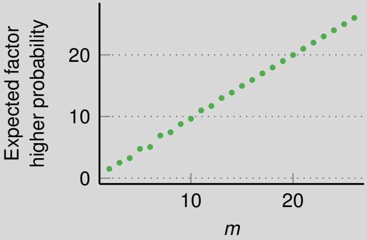
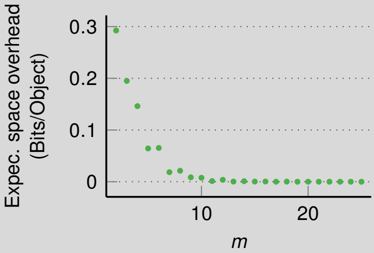
- store hash function *and* rotation

Lemma: Rotation Fitting

Let $|A| = \mathbb{A}$, $|B| = \mathbb{B}$, and $\mathbb{P}(R)$ be the probability of finding a bijection using rotation fitting. Let $\mathbb{P}(B)$ denote the probability of finding a bijection using RecSplit’s brute force strategy. Then, $\mathbb{P}(R) \to m\mathbb{P}(B)$ for $m \to \infty$.





Rotation Fitting (2/3)

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Proof (Sketch)

- consider number of different injective functions under cyclic shifts
- bit vector of length m with $\mathbb{B}$ set bits
- total number of equivalence classes under rotation is $\frac{1}{m}\sum_{d \text{ divides } \gcd(\mathbb{A},\mathbb{B})} \phi(d)\binom{m/d}{\mathbb{B}/d}$
- probability of the event $\mathcal{I}$ that there is a rotation has the m least significant bits set is

$$\mathbb{P}(\mathcal{I}) \geq m\frac{1}{\sum_{d \text{ divides } \gcd(\mathbb{A},\mathbb{B})} \phi(d)\binom{m/d}{\mathbb{B}/d}},$$

- $\phi(i) = |\{j \leq i \colon \gcd(i,j) = 1\}|$ is Euler's totient function

Proof (Sketch, cnt.)

- determine the probability $\mathbb{P}(R)$ using the events
  - $\mathcal{A}$: `popcount(a)=`$\mathbb{A}$
  - $\mathcal{B}$: `popcount(b)=`$\mathbb{B}$
  - B: found bijection using brute-force

Rotation Fitting (3/3)

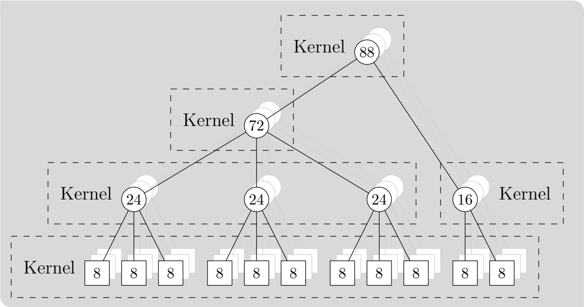
Proof (Sketch, ctn.)

$$
\begin{aligned}
\mathbb{P}(R) &= \mathbb{P}(\mathcal{A})\mathbb{P}(\mathcal{B})\mathbb{P}(\mathcal{I}) \\
&\geq \frac{m!}{(m-\mathbb{A})!m^{\mathbb{A}}} \cdot \frac{m!}{(m-\mathbb{B})!m^{\mathbb{B}}} \cdot \mathbb{P}(\mathcal{I}) = \frac{m!}{m^m} \cdot \frac{m!}{\mathbb{A}!\mathbb{B}!} \cdot \mathbb{P}(\mathcal{I}) = \mathbb{P}(B) \cdot \frac{m!}{\mathbb{A}!\mathbb{B}!} \cdot \mathbb{P}(\mathcal{I}) \\
&\geq \mathbb{P}(B) \cdot \frac{m!}{\mathbb{A}!\mathbb{B}!} \cdot m \frac{1}{\sum_{d|\gcd(\mathbb{A},\mathbb{B})} \phi(d)\binom{m/d}{b/d}} = \mathbb{P}(B) \cdot m \cdot \frac{m!}{m! + (\mathbb{A}!\mathbb{B}!)\sum_{d|gcd(\mathbb{A},\mathbb{B}),d\neq 1} \phi(d)\binom{m/d}{b/d}} \\
&= \mathbb{P}(B) \cdot m \cdot \frac{1}{1 + \sum_{d|\gcd(\mathbb{A},\mathbb{B}),d\neq 1} \phi(d) \frac{(m/d)!\mathbb{A}!\mathbb{B}!}{m!(\mathbb{A}/d)!(\mathbb{B}/d)!}} \\
&\sim \mathbb{P}(B) \cdot m \cdot \frac{1}{1 + \sum_{d|\gcd(\mathbb{A},\mathbb{B}),d\neq 1} \phi(d)\sqrt{d}\frac{\mathbb{A}^{\mathbb{A}-\mathbb{A}/d}\mathbb{B}^{\mathbb{B}-\mathbb{B}/d}}{m^{m-m/d}}} \\
&\to \mathbb{P}(B) \cdot m \text{ for } m \to \infty
\end{aligned}
$$

Parallel RecSplit on the GPU

Computing on the GPU

- several streaming multiprocessors (SMs)
- each SM contains many arithmetic logic units (ALUs)
- several threads operat in lock-step (warp)
- to hide latencies, each SM is oversubscribed with more threads than ALUs
- in CUDA, kernels are functions that can be executed on the GPU
- a kernel is executed on a grid of thread blocks



- use GPU to determine splitting and bijections

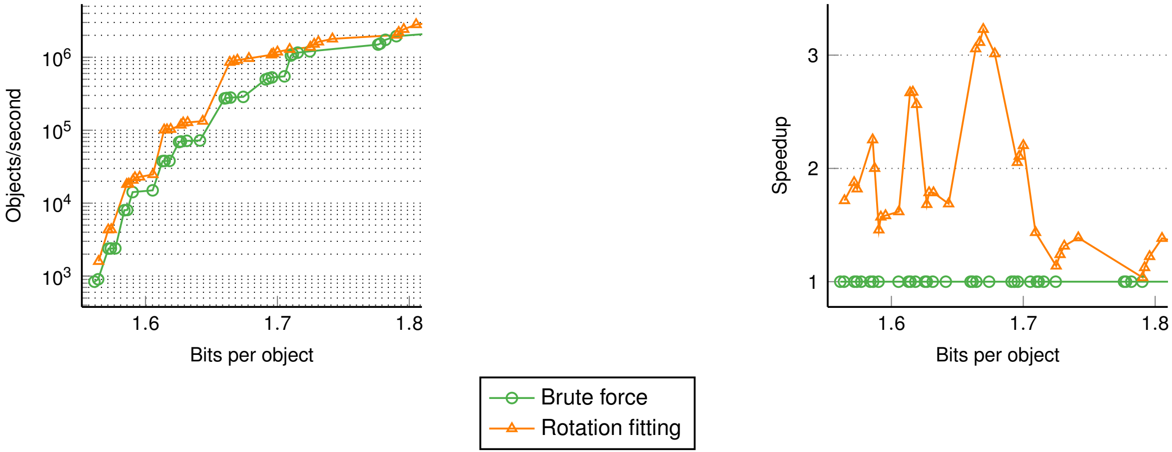
Experimental Evaluation

- Intel i7 11700 processor with 8 cores (16 hardware threads (HT)), base clock: 2.5 GHz
- AVX-512.
- Ubuntu 22.04 with Linux 5.15.0
- NVIDIA RTX 3090 GPU

- AMD EPYC 7702P processor with 64 cores (128 hardware threads), base clock: 2.0 GHz
- AVX2
- Ubuntu 20.04 with Linux 5.4.0

- GNU C++ compiler v.11.2.0 (`-O3 -march=native`)

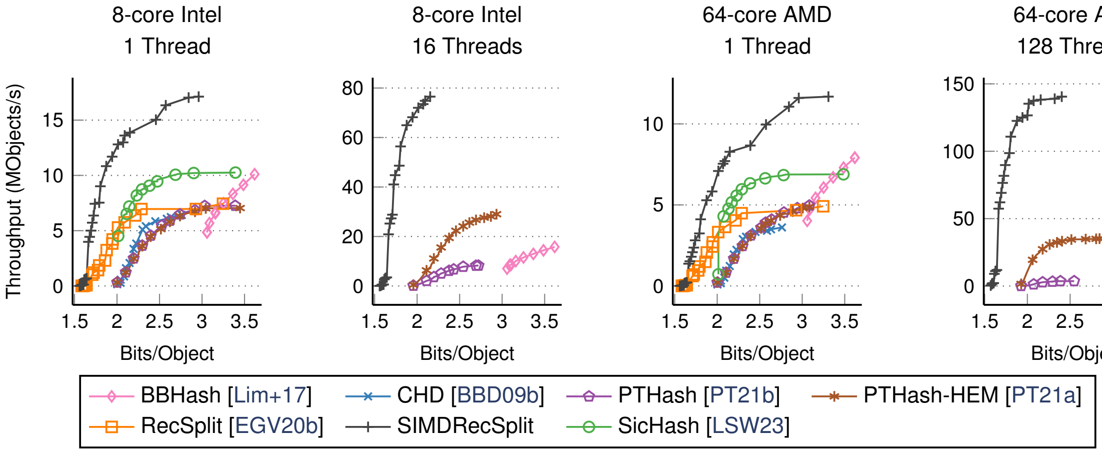
Rotation Fitting



Overview Results

Configuration	Method	Bijections	Threads	B/Obj	Constr.	Speedup
$\ell = 16, b = 2000$	RecSplit [EGV20b]	Brute force	1	1.560	1175.4	1
	RecSplit	Brute force	16	1.560	206.5	5
	SIMDRecSplit	Rotation fitting	1	1.560	138.0	8
	SIMDRecSplit	Rotation fitting	16	1.560	27.9	42
	GPURecSplit	Brute force	GPU	1.560	1.8	655
	GPURecSplit	Rotation fitting	GPU	1.560	1.0	1173
$\ell = 18, b = 50$	RecSplit [EGV20b]	Brute force	1	1.707	2942.9	1
	RecSplit	Brute force	16	1.713	504.0	5
	SIMDRecSplit	Rotation fitting	1	1.709	58.3	50
	SIMDRecSplit	Rotation fitting	16	1.708	12.3	239
	GPURecSplit	Brute force	GPU	1.708	5.2	564
	GPURecSplit	Rotation fitting	GPU	1.709	0.5	5438
$\ell = 24, b = 2000$	GPURecSplit	Brute force	GPU	1.496	2300.9	—
	GPURecSplit	Rotation fitting	GPU	1.496	467.9	—

Comparison with Competitors

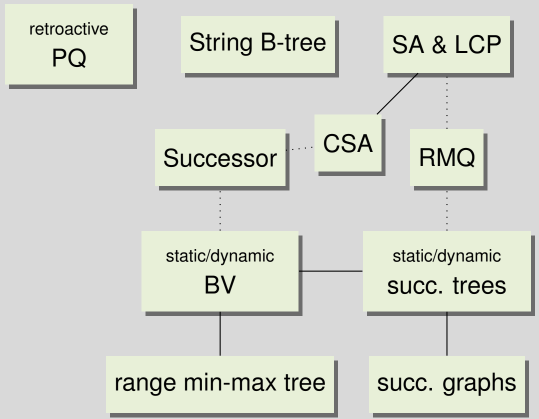


Conclusion and Outlook

This Lecture

- conclusion retroactive data structures
- minimal perfect hash functions

Advanced Data Structures



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Next Lecture (15.07.2024)

- **NO LECTURE ON 08.07.2024**
- learned data structures

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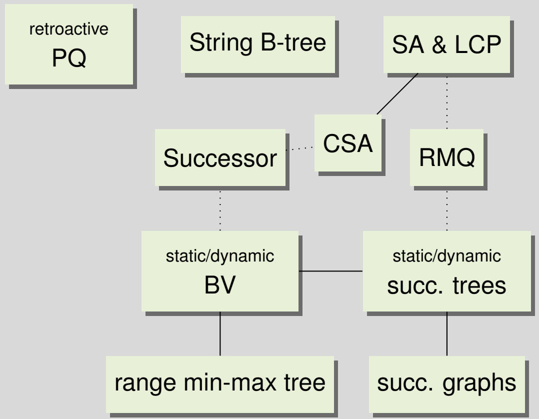

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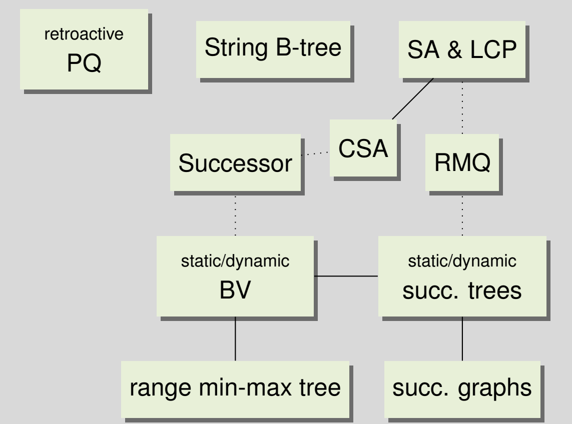
Oral Exams and Project

- registration exams and project will open this week
- exam dates: 19.08., 20.08., 26.08.,28.08., 30.08., 09.09, and 11.09.

Advanced Data Structures



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