

# Lecture 4: LCP Arrays Range Minimum Queries

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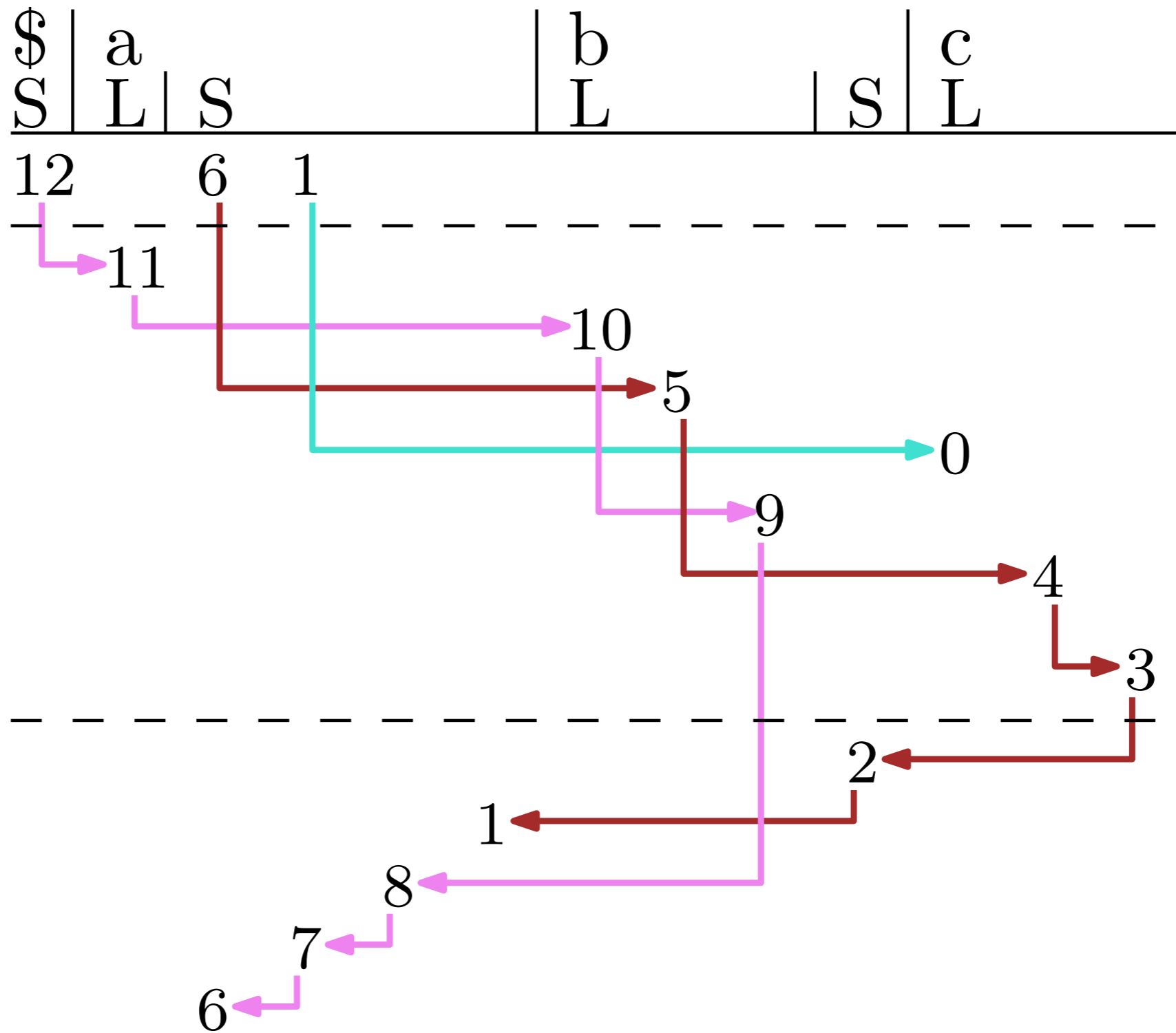


# Algorithm sais

- Definition: suffix  $T[i,n]$  called
  - ▶ **S-type** iff  $T[i..n] <_{\text{lex}} T[i+1..n]$  ( $T[n,n]='\$'$  always S)
  - ▶ **L-type** otherwise
- 1. Choose sample: leftmost  $S$  (call them  $S^*$ ),  $|S^*| < 1/2n$
- 2. Sort  $S^*$ -suffixes by **recursion**
  - ▶ on new text formed by sorted  $S^*$ -substrings
- 3. Scan  $A$  from left to right (say we're at pos.  $i$ ):
  - ▶ if  $T[A[i]-1]$  is **L**, write  $A[i]-1$  to 1st pos. in bucket
- 4. like (3), but sorting **S**-suffixes in a right-to-left scan

$T =$ 

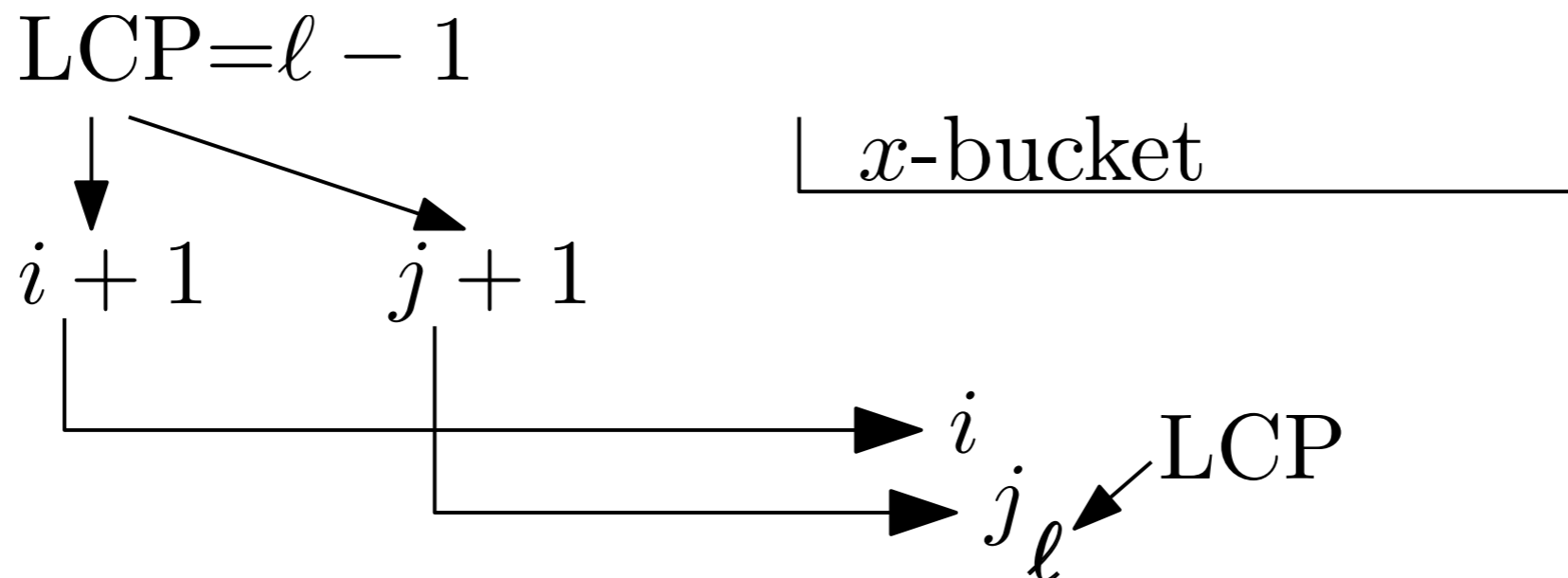
0	1	2	3	4	5	6	7	8	9	10	11	12
c	a	b	c	c	b	a	a	a	b	b	a	\$
L	S*	S	L	L	L	S*	S	S	L	L	L	S*



# Inducing LCPs

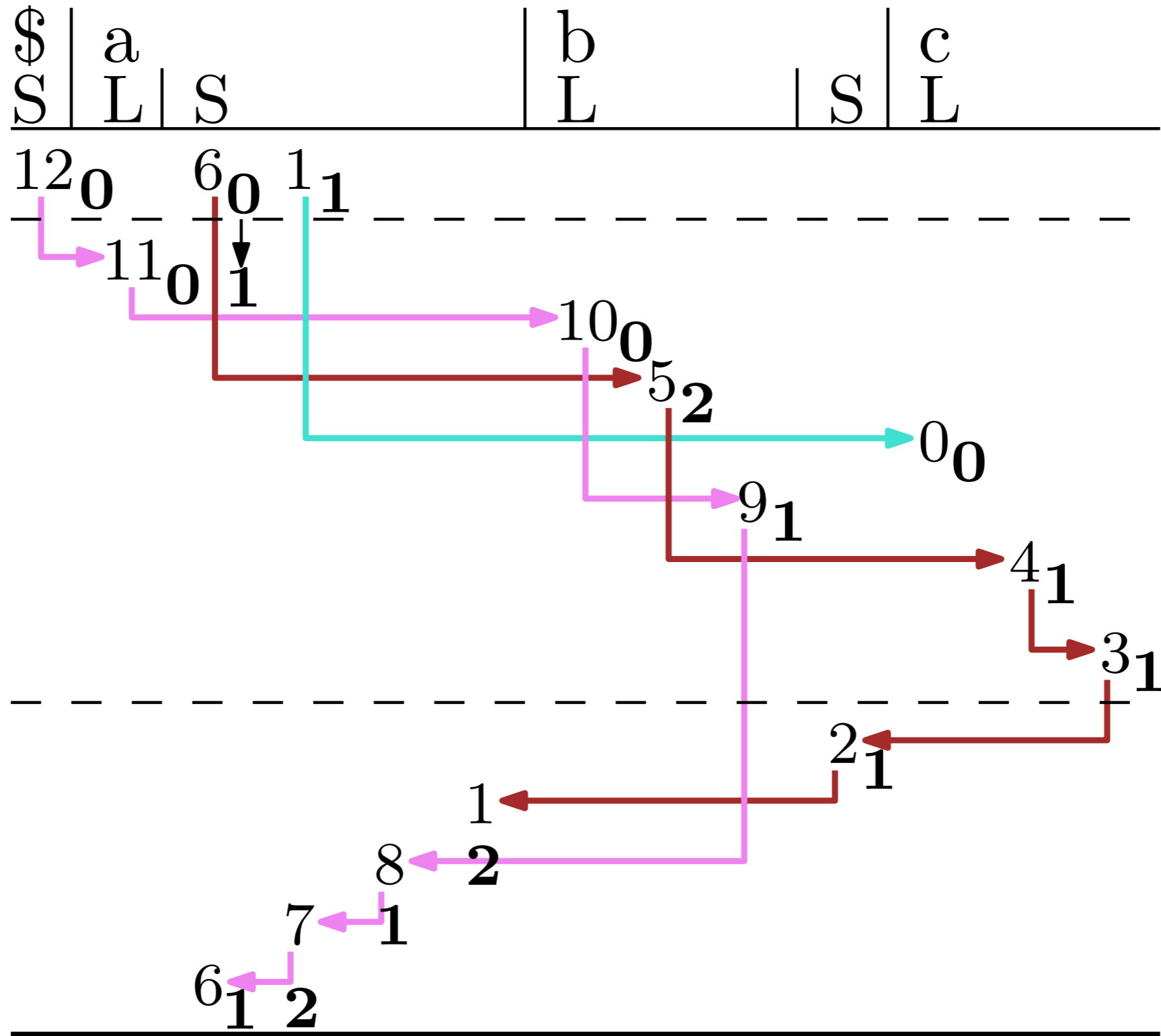
## ● Idea

- ▶ placing  $T[i..n]$  and  $T[j..n]$  at adjacent positions...
- ▶ ... allows us to induce LCP ...
- ▶ ... from LCP of  $T[i+1..n]$  and  $T[j+1..n]$ ...
- ▶ ... which is already known!



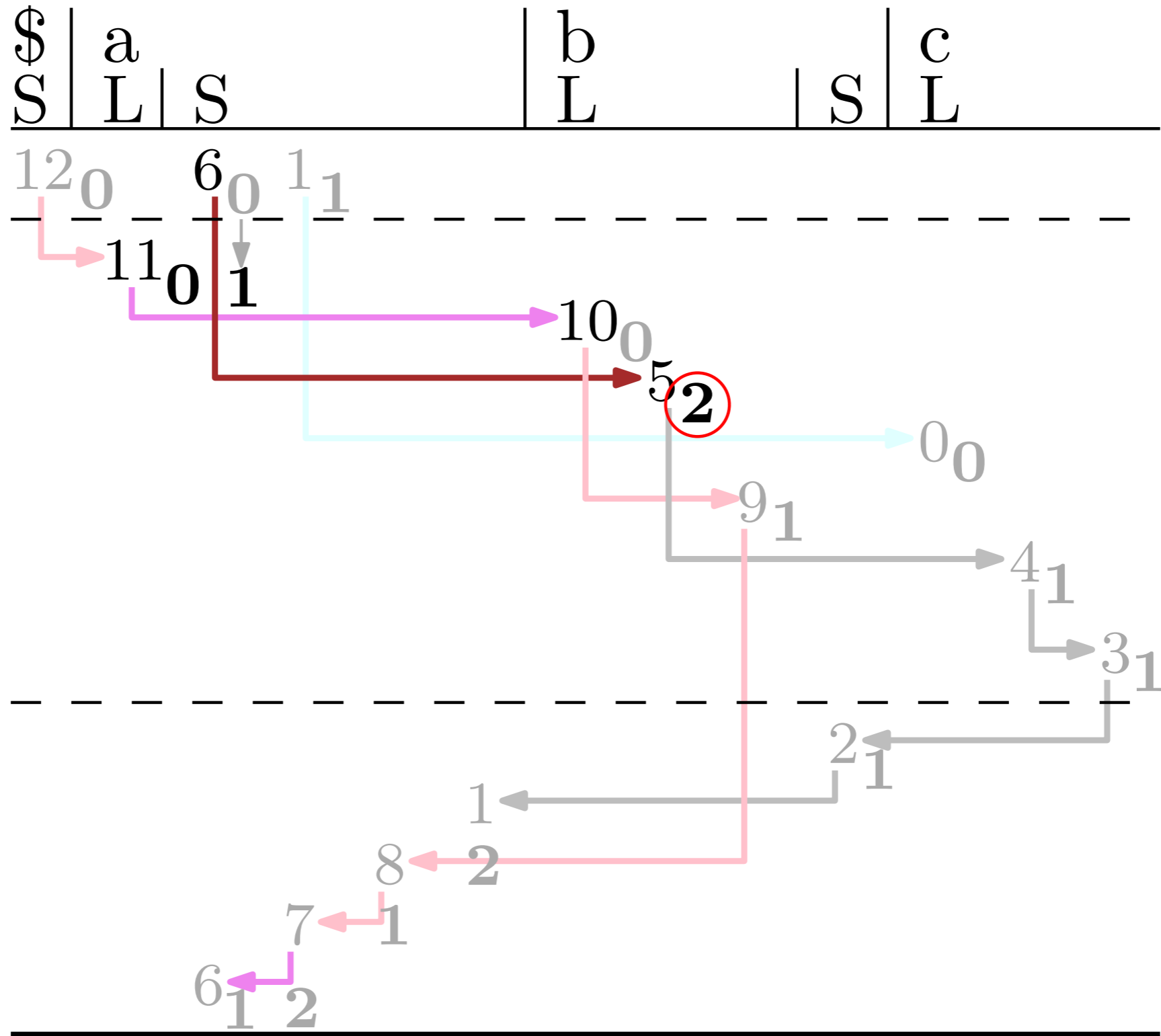
$T =$ 

0	1	2	3	4	5	6	7	8	9	10	11	12
c	a	b	c	c	b	a	a	a	b	b	a	\$
L	S*	S L L L S*					S S L L L S*					



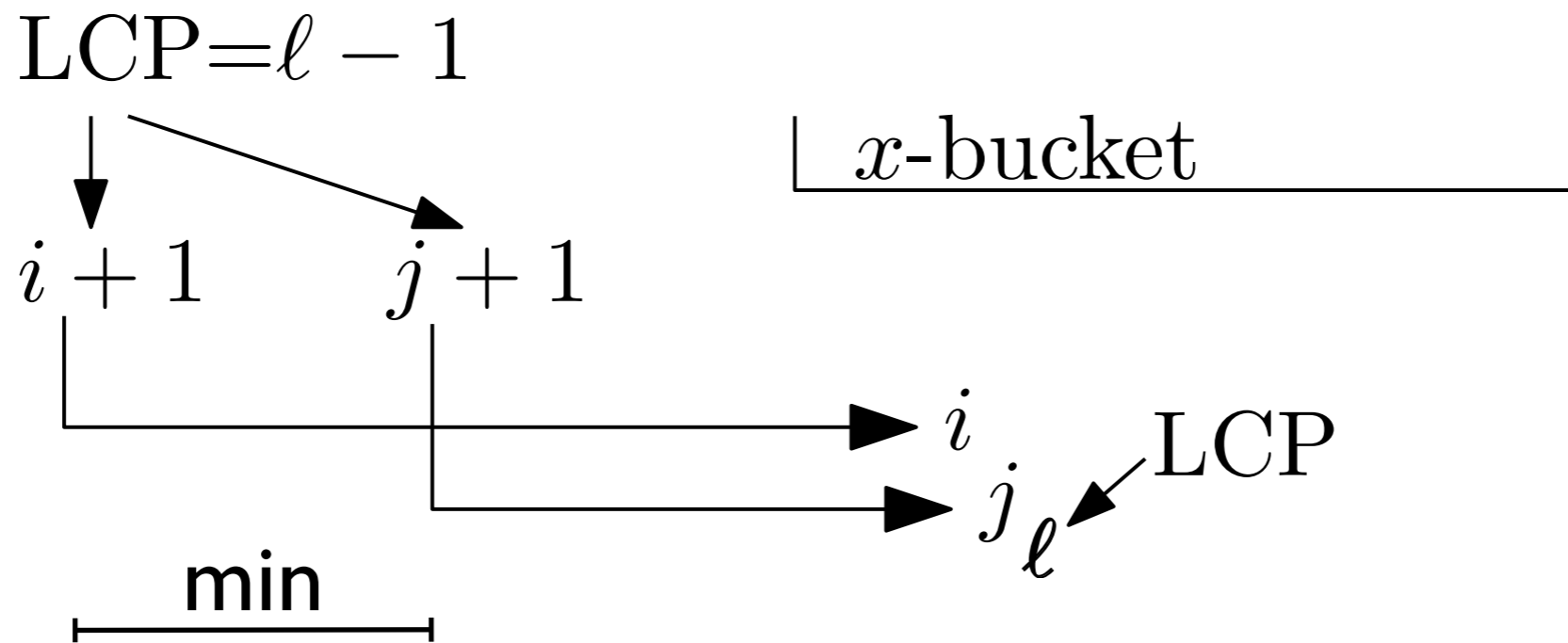
$T =$ 

0	1	2	3	4	5	6	7	8	9	10	11	12
c	a	b	c	c	b	a	a	a	b	b	a	\$
L	S*	S L L L S*					S S L L L S*					



# One Detail

- What is LCP of suffixes  $i+1$  and  $j+1$ ?



- In general:

$$LCP(T[x..n], T[y..n]) = LCP[\text{RMQ}_{LCP}(SA^{-1}[x], SA^{-1}[y])]$$