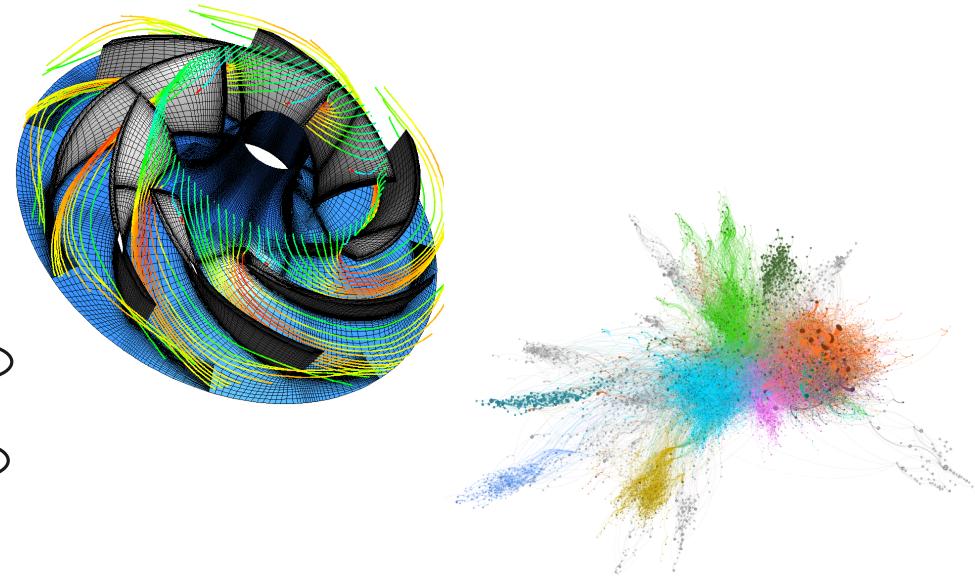
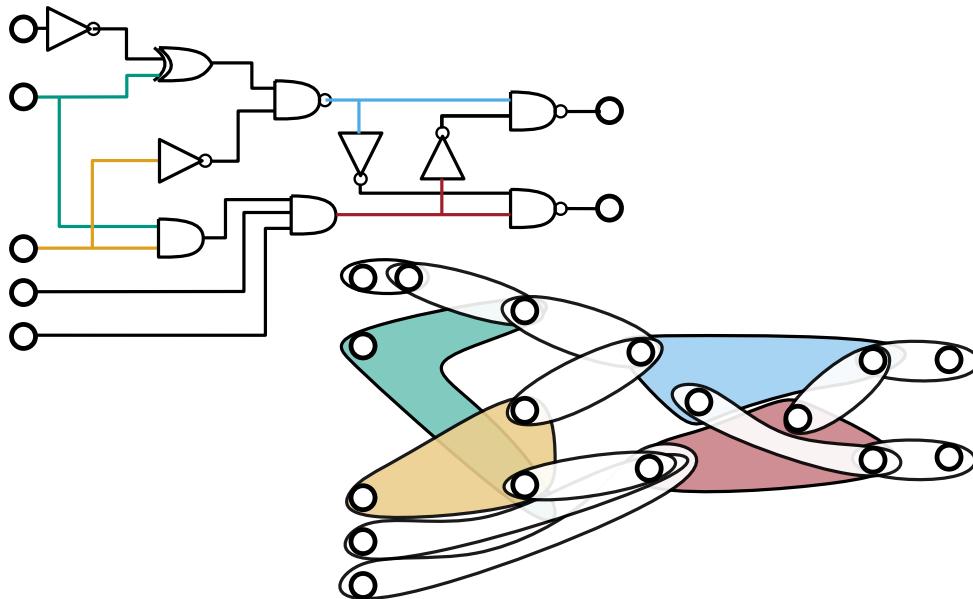


Recent Advances in (Hyper-)Graph Partitioning

Annual SPP Meeting · September 17, 2019

Peter Sanders and Tobias Heuer

INSTITUTE OF THEORETICAL INFORMATICS · ALGORITHMIC GROUP



News on (Hyper)graph Partitioning

■ Scalable Edge Partitioning

ALENEX'19, S. Schlag, C. Schulz, D. Seemaier, D. Strash

■ Faster Support Vector Machines

ALENEX'19, S. Schlag, M. Schmitt, C. Schulz

■ Data Distribution for Phylogenetic Inference with Site Repeats via Judicious Hypergraph Partitioning

IPDPSW'19, I. Baar, L. Hübner, P. Oettig, A. Zapletal, S. Schlag, A. Stamatakis, B. Morel

■ Network Flow-Based Refinement for Multilevel Hypergraph Partitioning

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■ Evaluation of a Flow-Based Hypergraph Bipartitioning Algorithm

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Scalable Edge Partitioning

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The Edge Partitioning Problem

Partition **edge set** of graph $G = (V, E, c, \omega)$ into **k** disjoint blocks

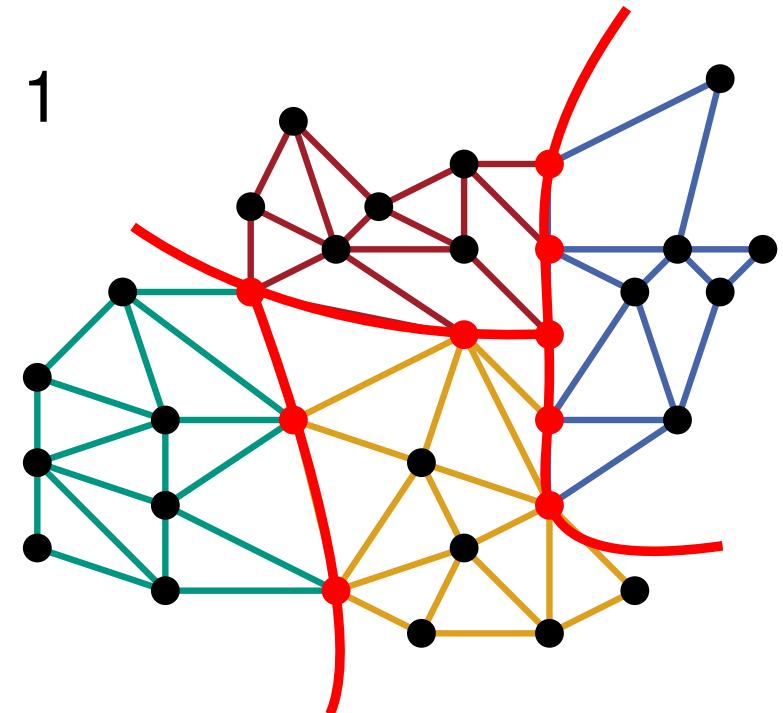
$\Pi = \{E_1, \dots, E_k\}$ such that

- Blocks E_i are roughly equal-sized:

$$\omega(E_i) \leq (1 + \varepsilon) \left\lceil \frac{\omega(E)}{k} \right\rceil$$

- minimize **vertex cut**:

$$\sum_{v \in V} |I(v)| - 1$$



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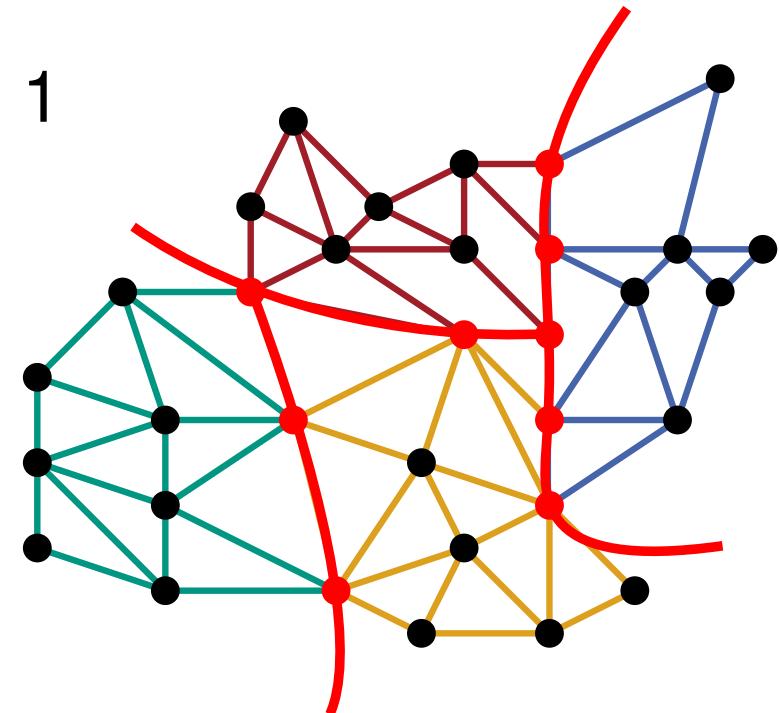
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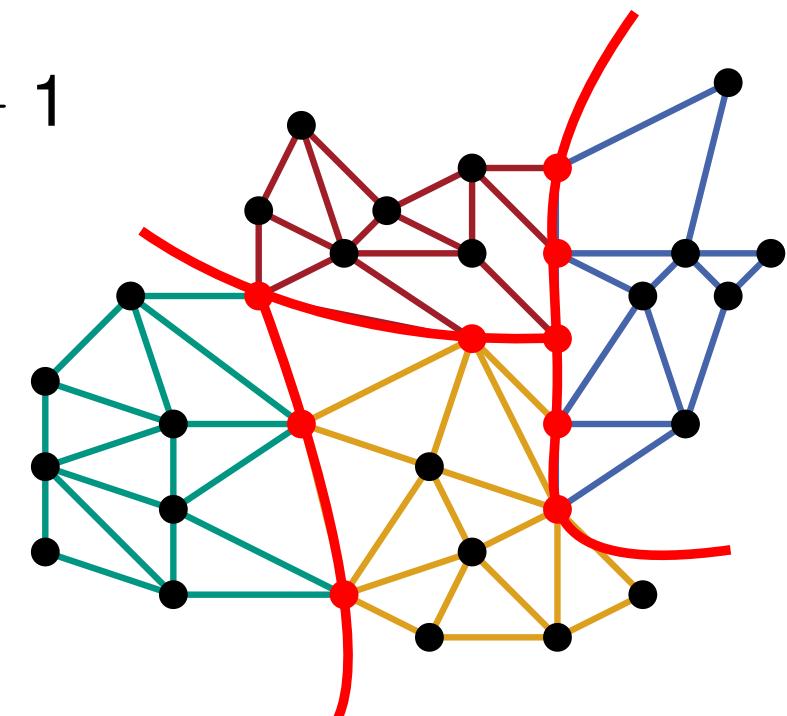
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blocks with edges incident to v

Motivation [Gonzalez et al.'12]:

- edge-centric distributed computations
- combat shortcomings of TLAV approaches
- duplicate node-centric computations



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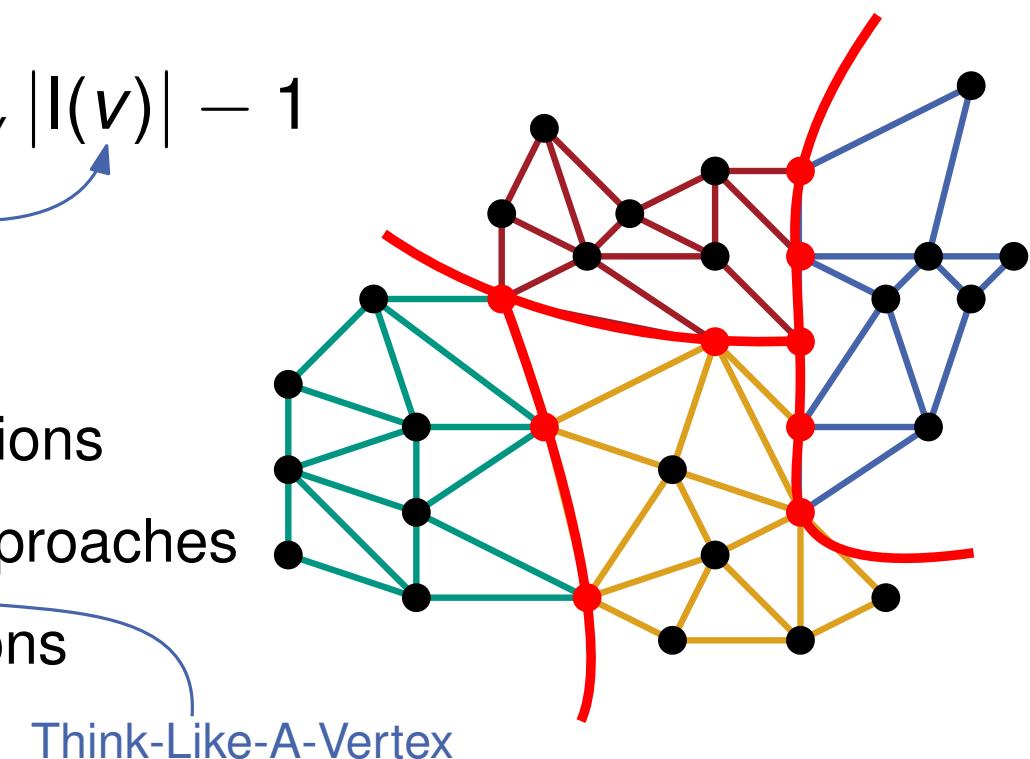
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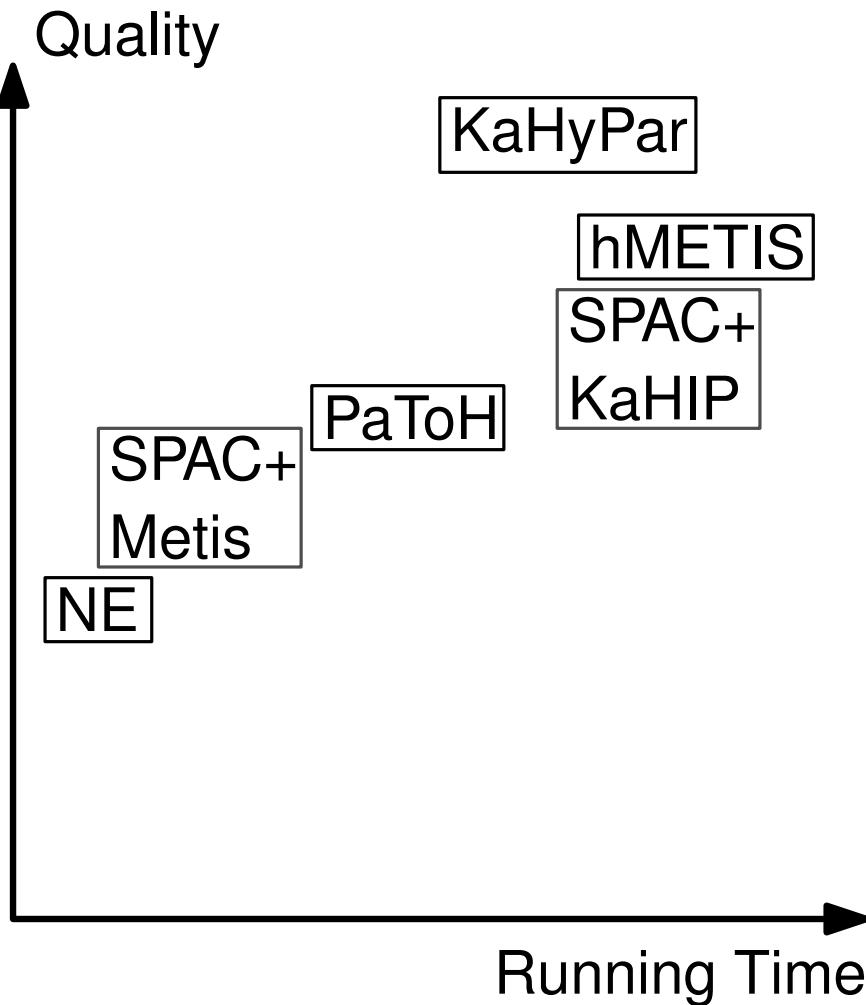
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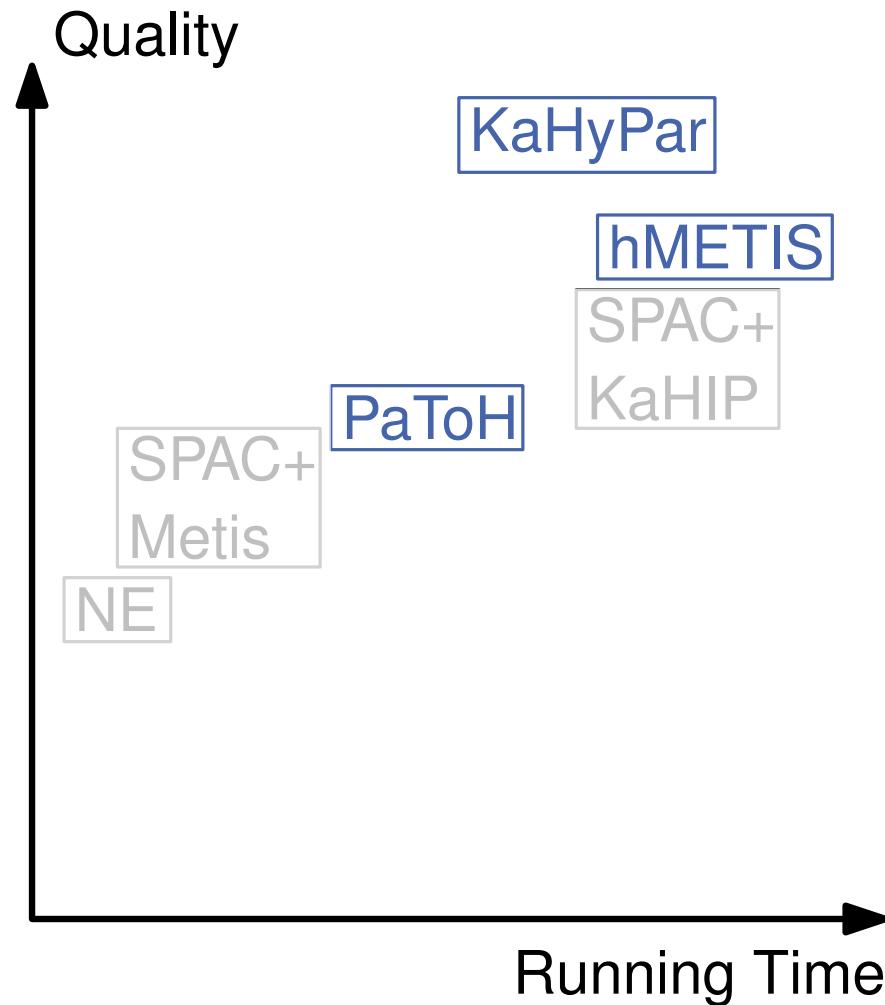
Edge Partitioning Algorithms

Sequential



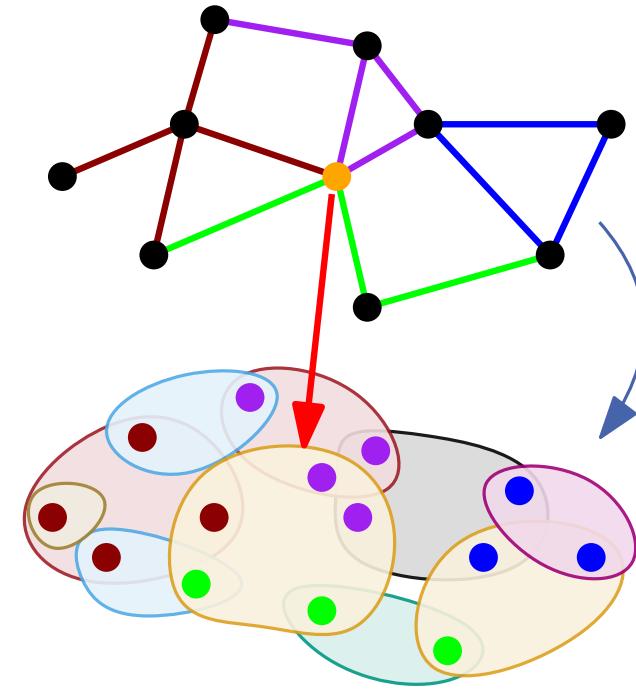
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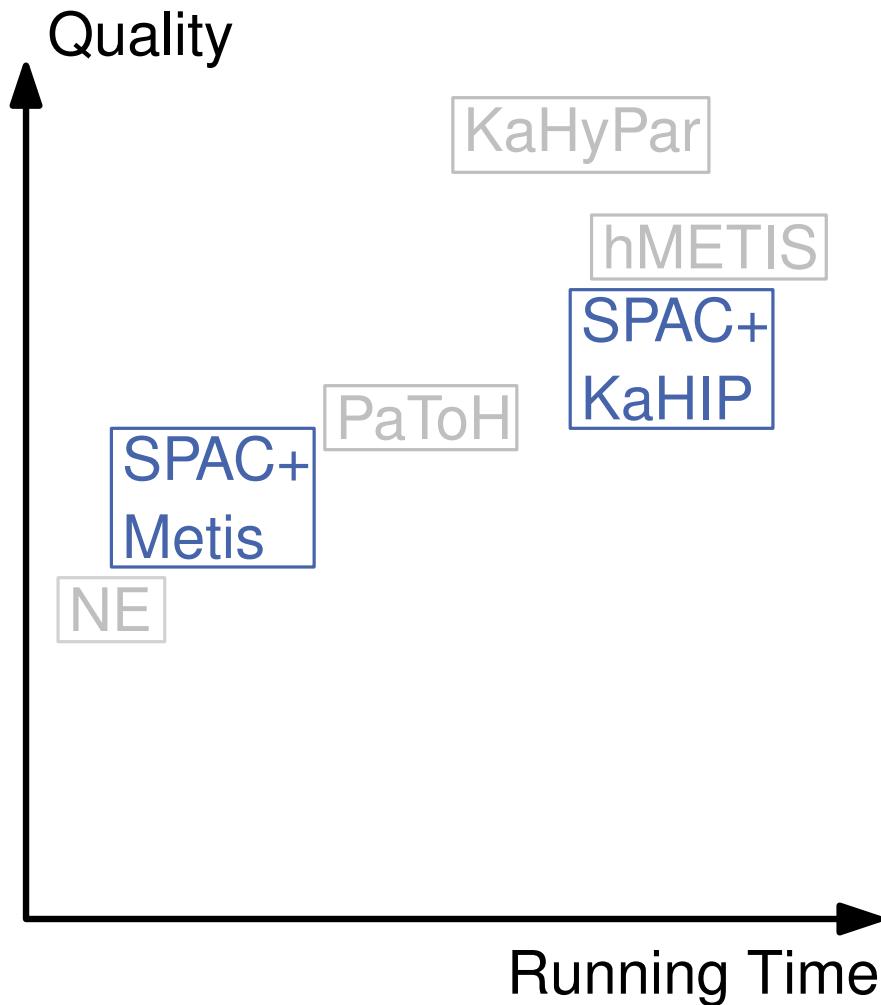
Hypergraph Model:

- Graph edge \rightsquigarrow vertex
- Graph node \rightsquigarrow hyperedge
- Optimize connectivity



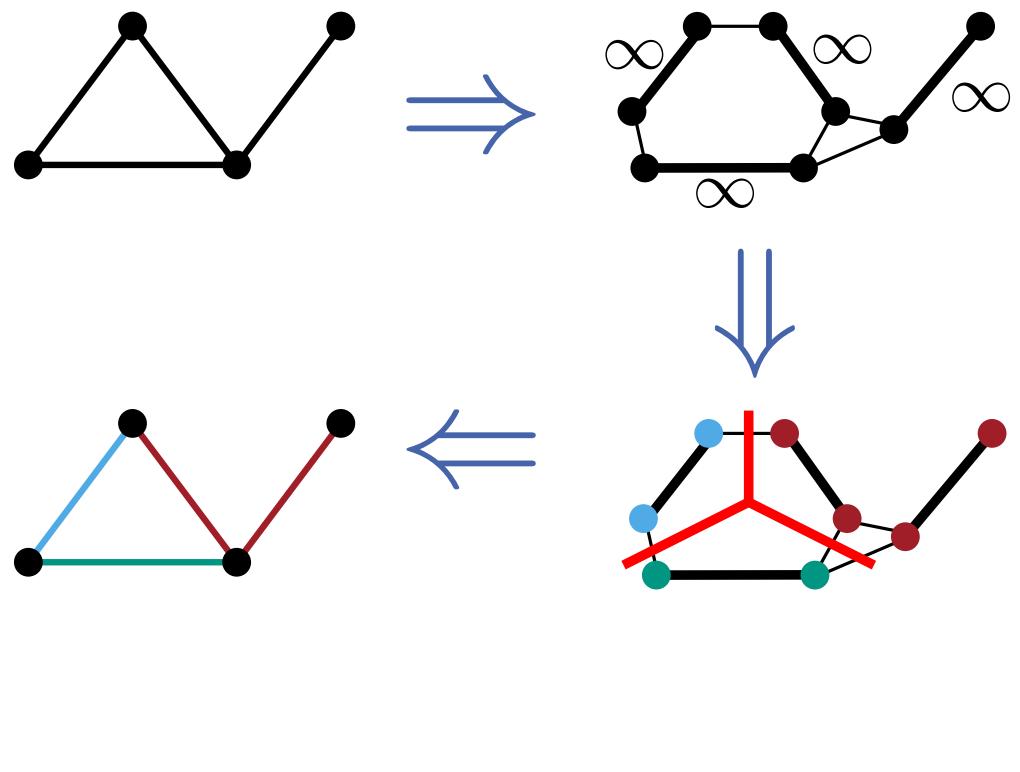
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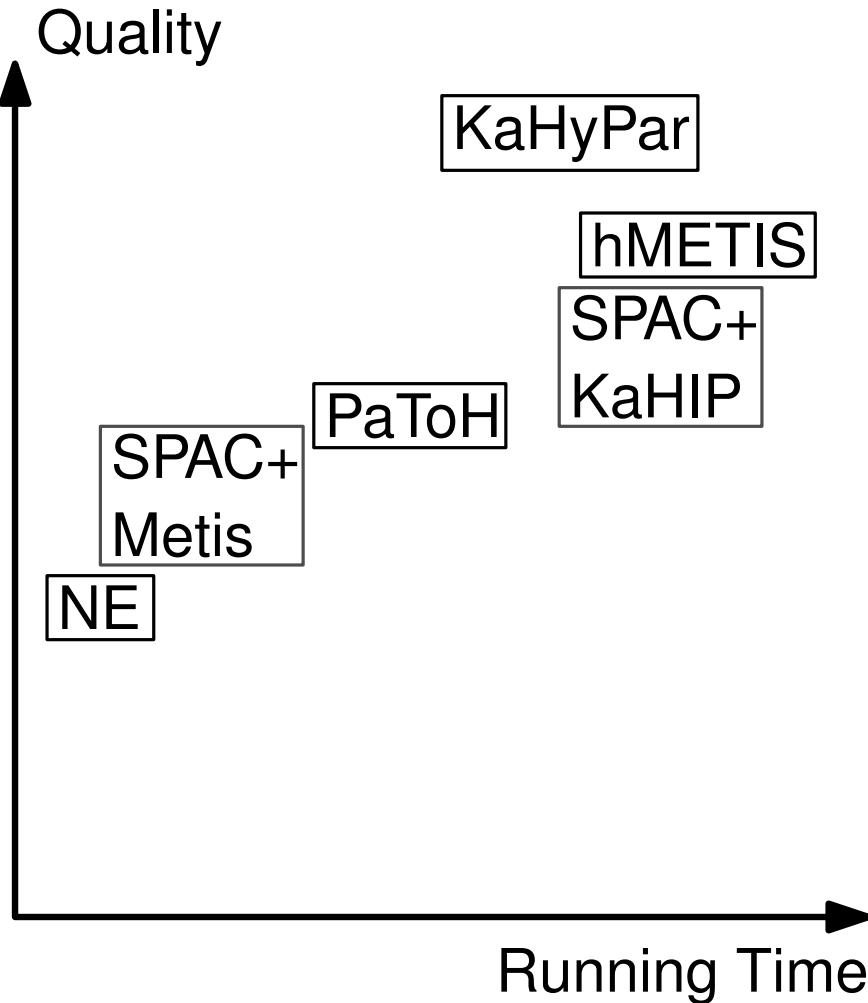
Split-And-Connect (SPAC) [Li et al.'17]:

- Build auxilary graph
- Use vertex partitioning algorithm

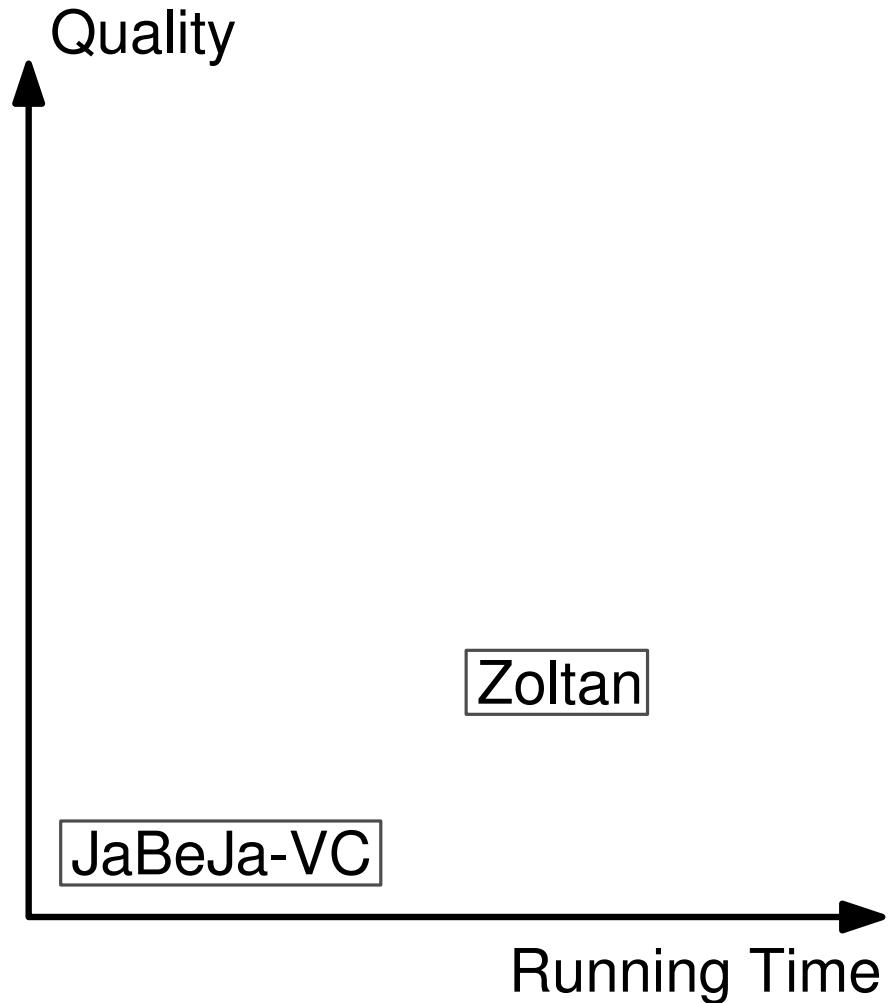


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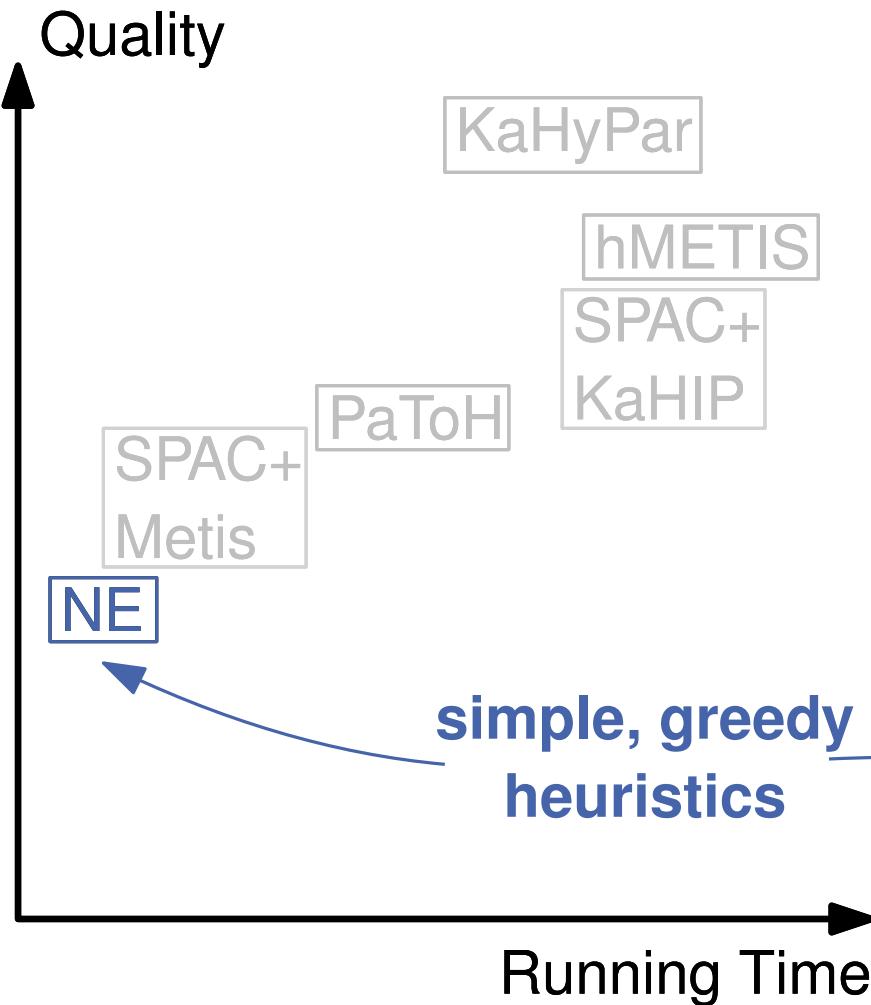


Distributed

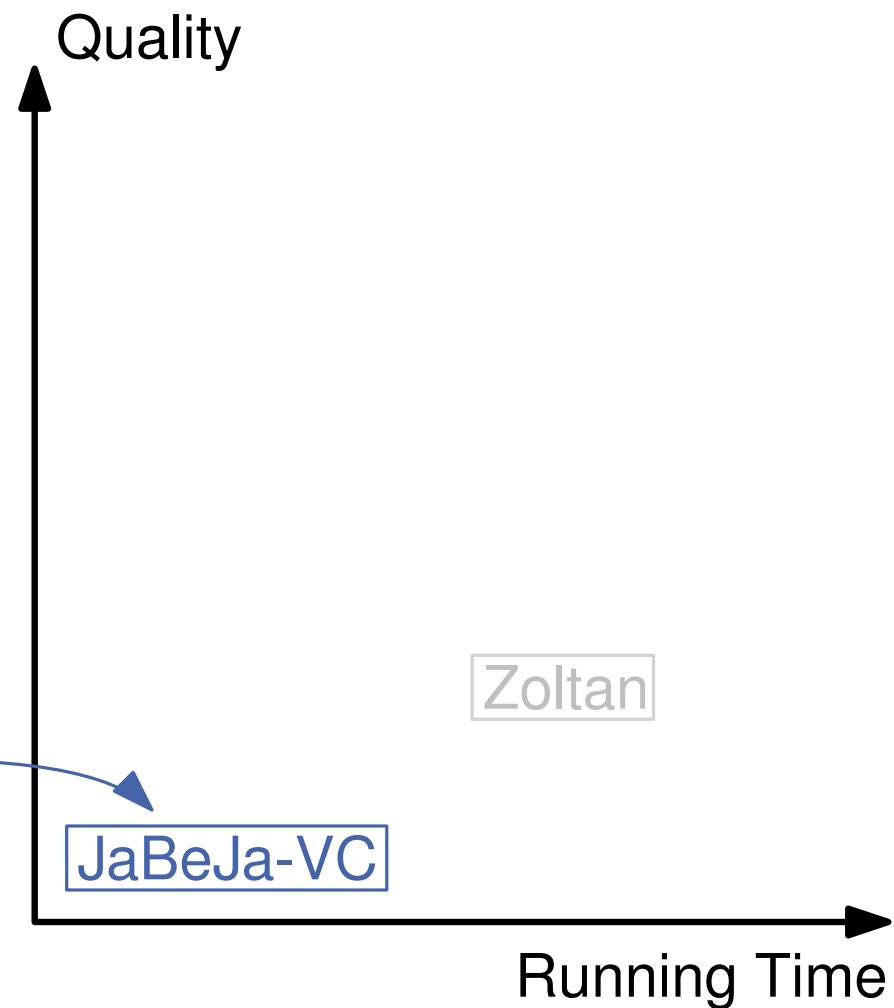


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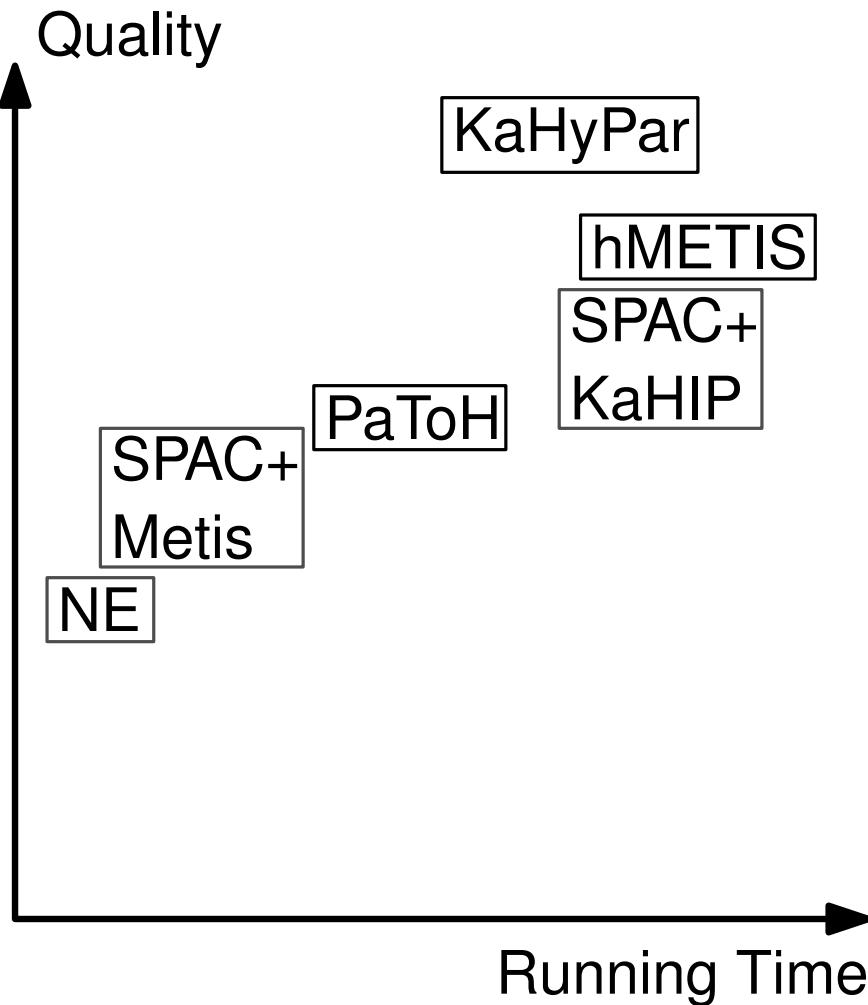


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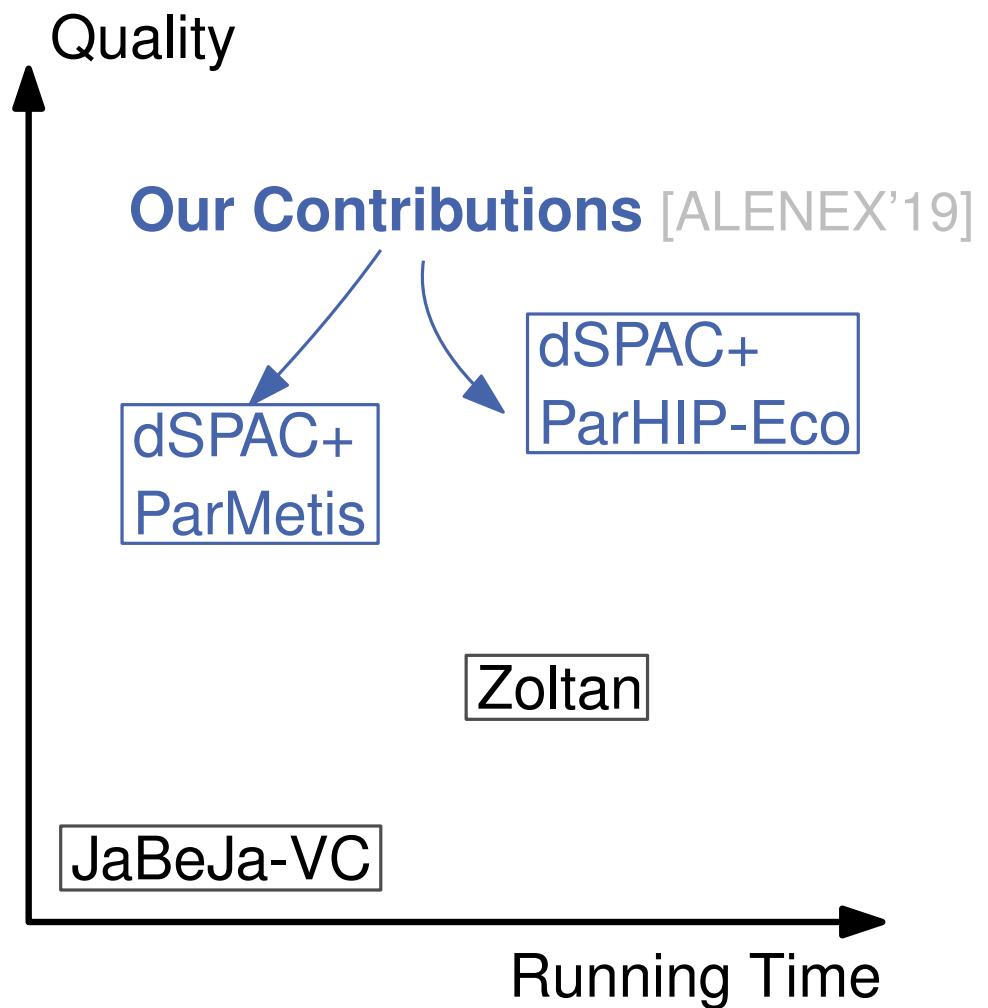


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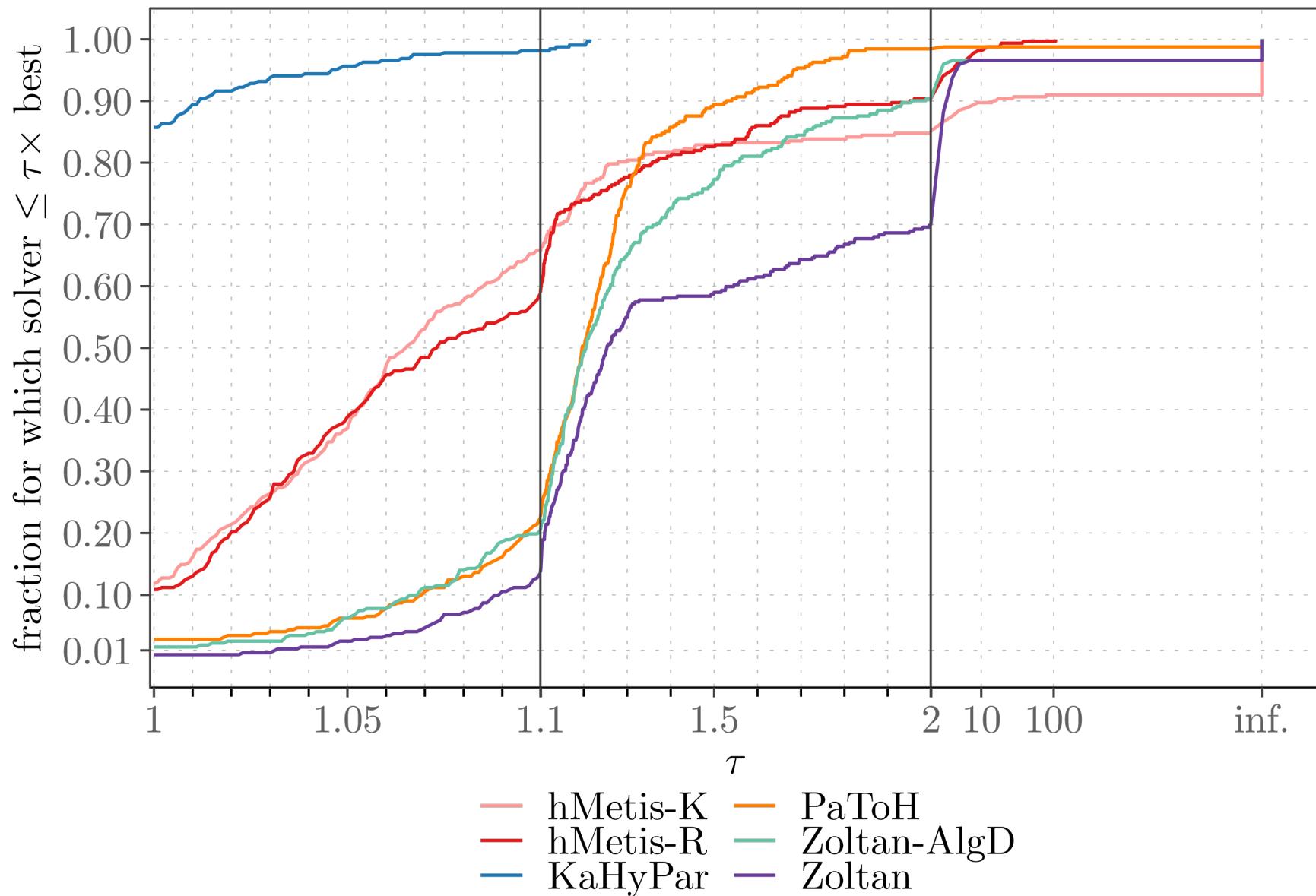
Experiments: Benchmark Setup

- Test suite: 70 graphs
 - Walshaw Graph Archive
 - Sparse Matrix-Vector Multiplication
 - Web & Social Graphs
 - Random Geometric Graphs
- $k \in \{2, 4, 8, 16, 32, 64, 128\}$
- Imbalance: $\epsilon = 3\%$
- Averages of 5 repetitions
- Sequential: 1 core
- Distributed: $32 * 20$ cores

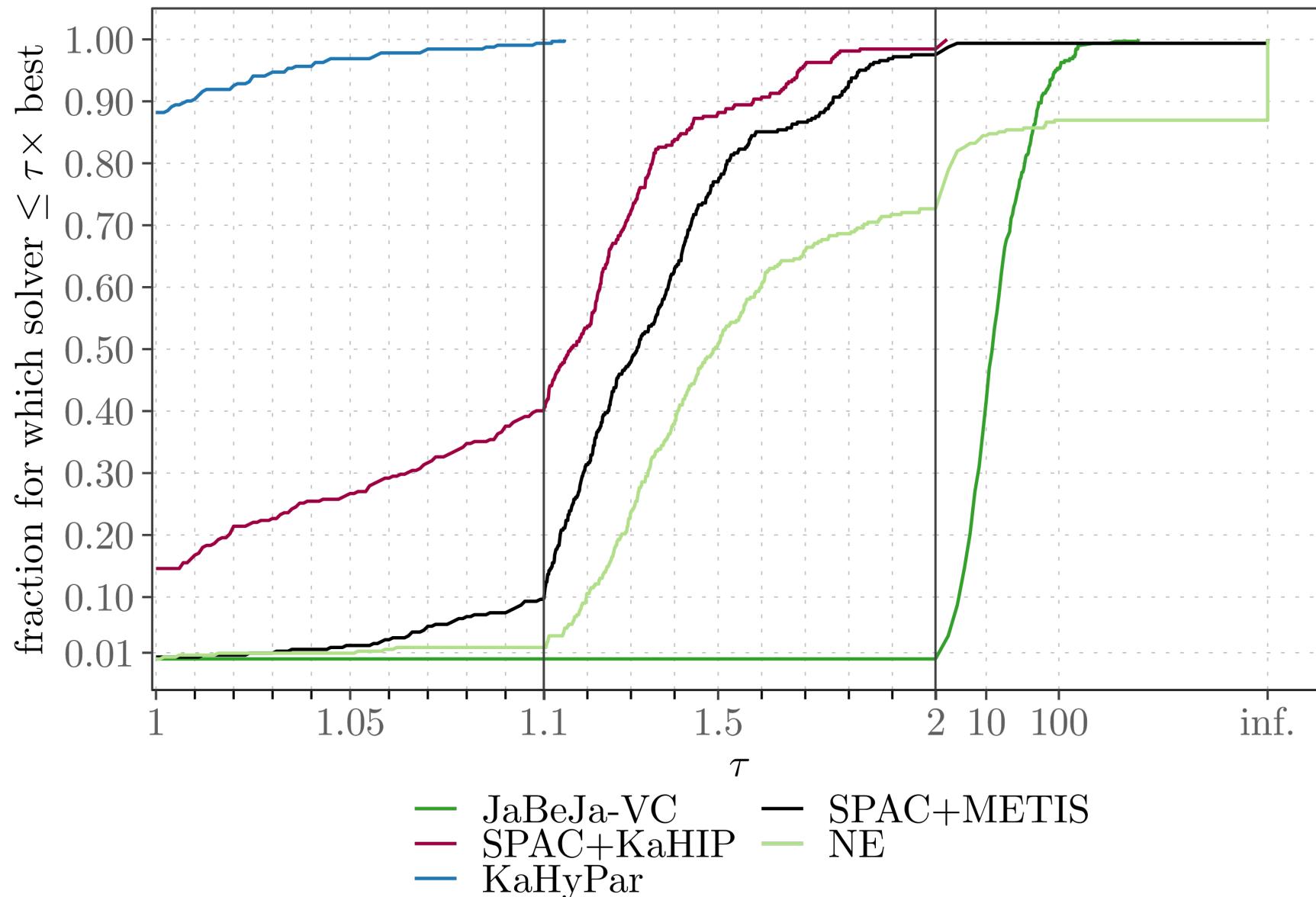
Competitors:

- | | |
|----------------------|-------|
| ■ KaHyPar-MF | HGP's |
| ■ PaToH | |
| ■ Zoltan | |
| ■ Zoltan-AlgD | |
| ■ hMetis- $\{R, K\}$ | |
| ■ JaBeJa-VC | |
| ■ NE | |
| ■ SPAC + KaHIP | |
| ■ SPAC + Metis | |
| ■ dSPAC + ParHIP | |
| ■ dSPAC + ParMetis | |

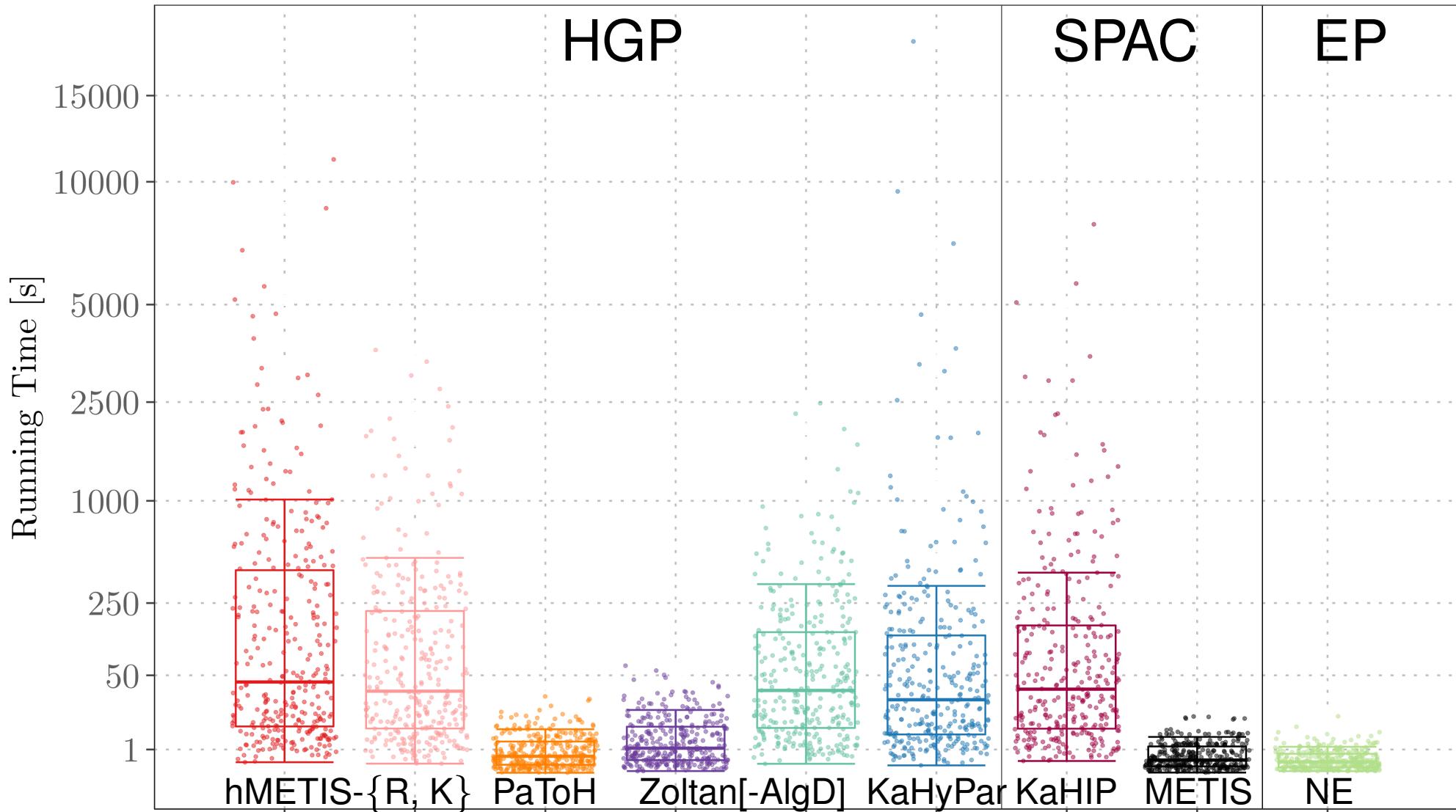
Experiments: Sequential HGP



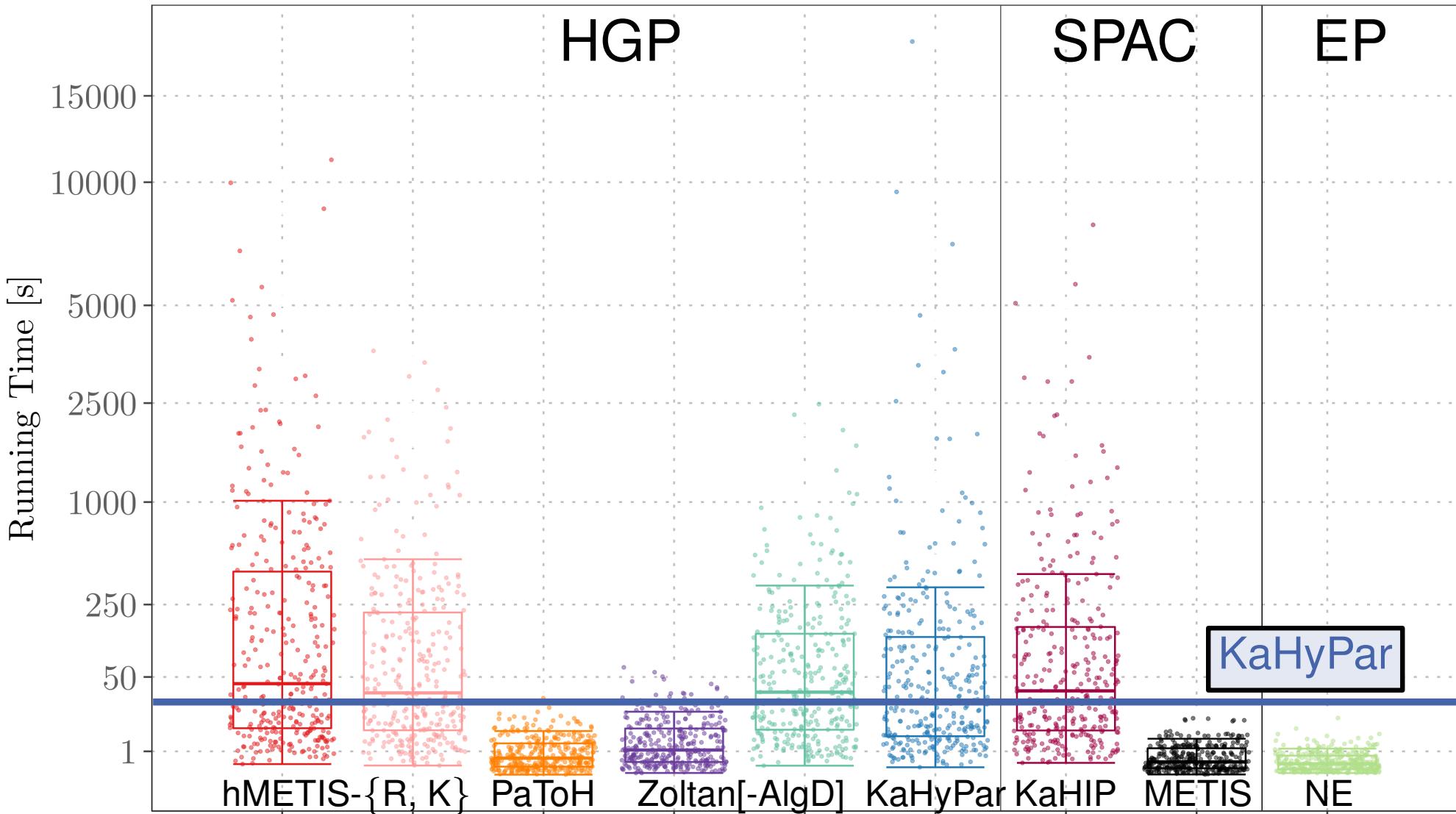
Experiments: Sequential HGP & SPAC+X



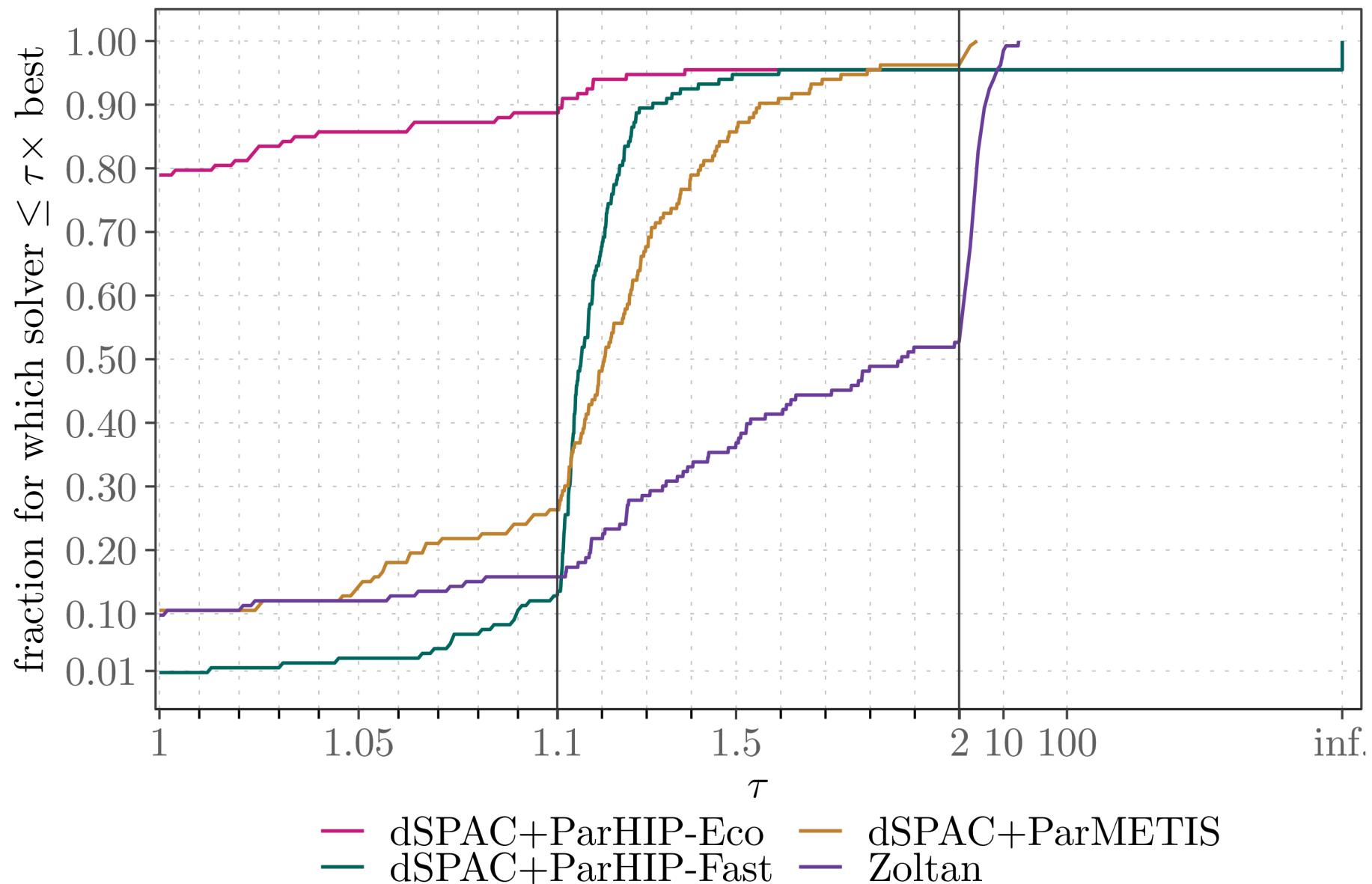
Experiments: Sequential Running Time



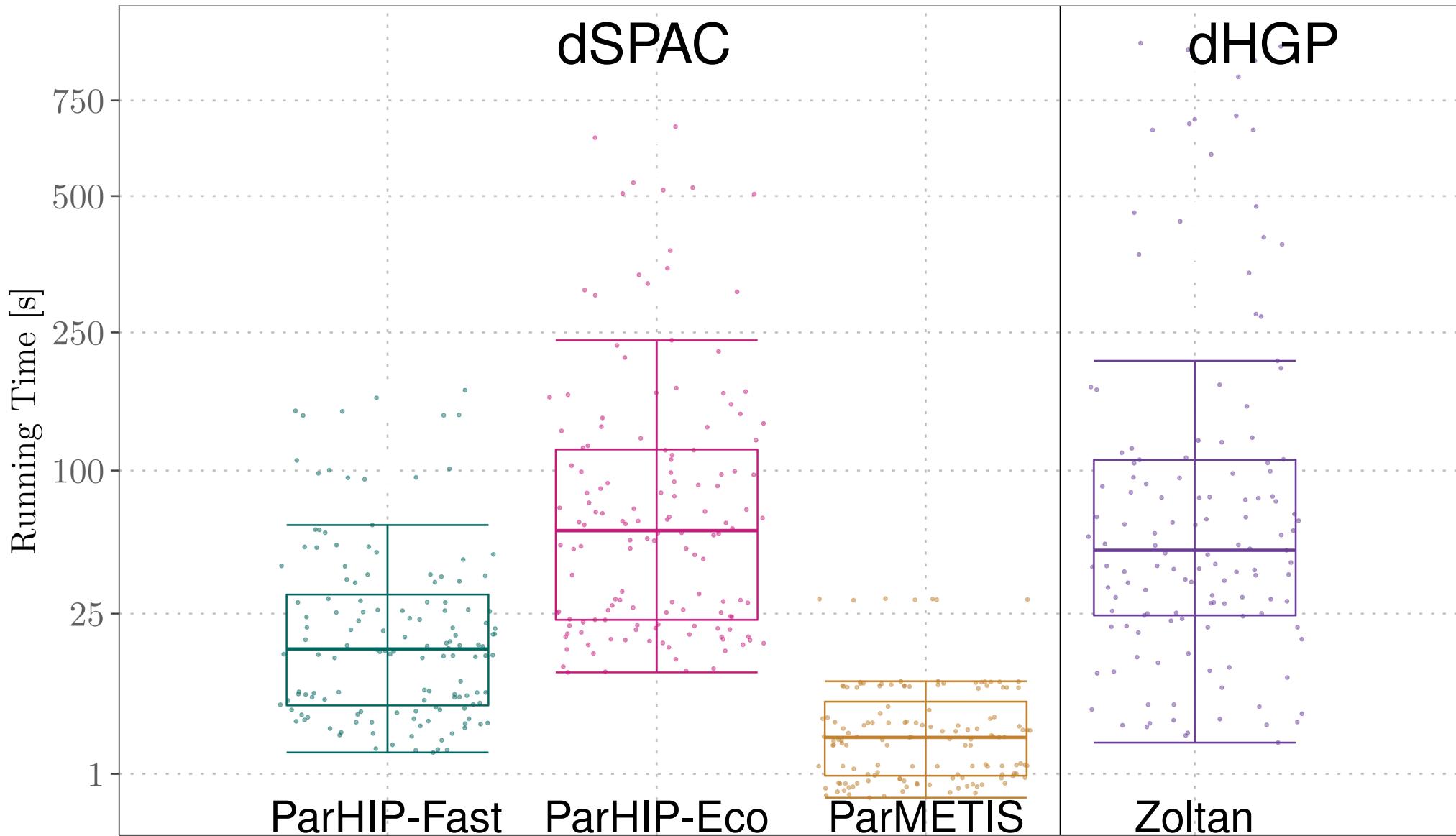
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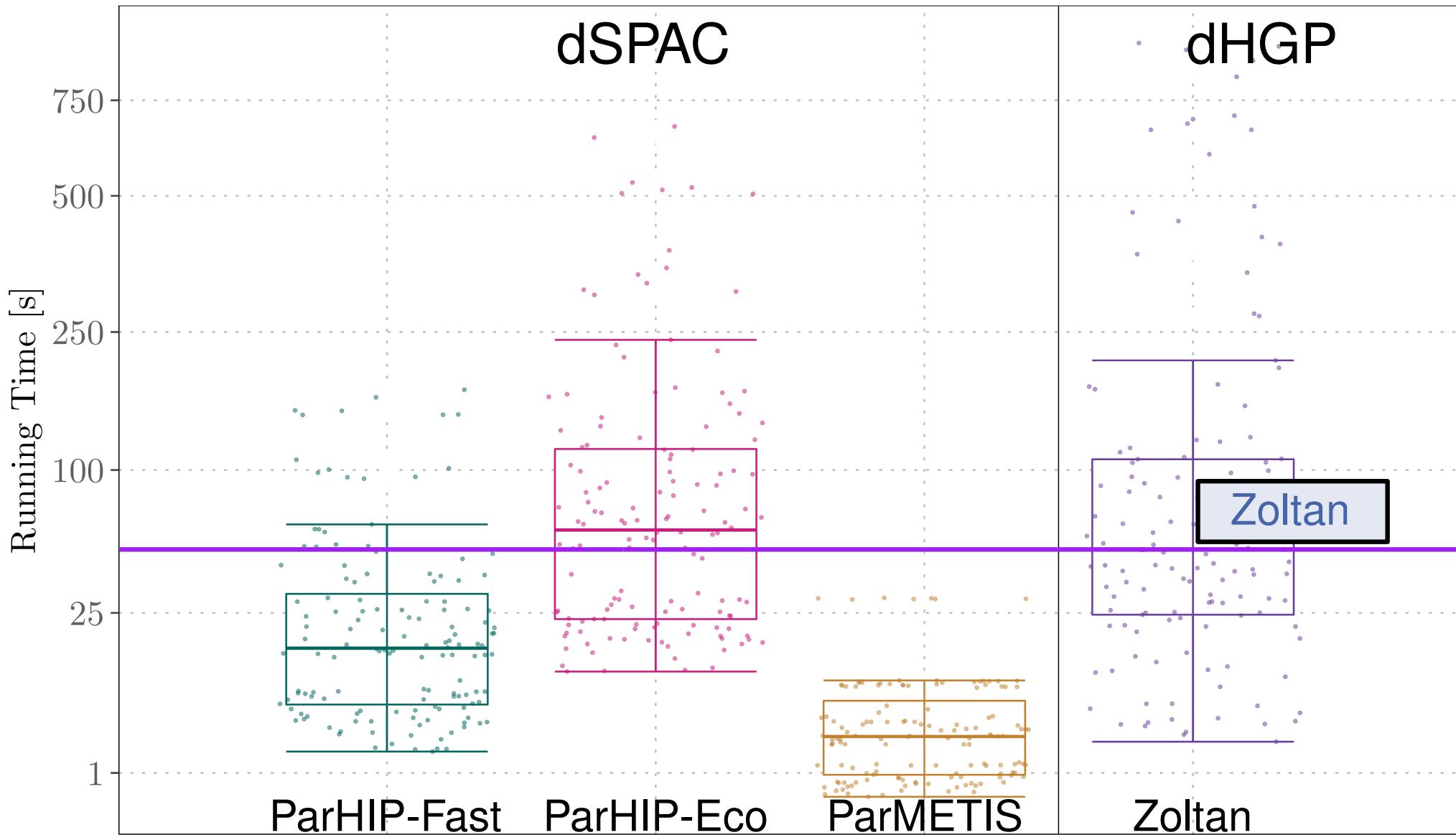
Experiments: Distributed HGP & dSPAC+X



Experiments: Distributed Running Time



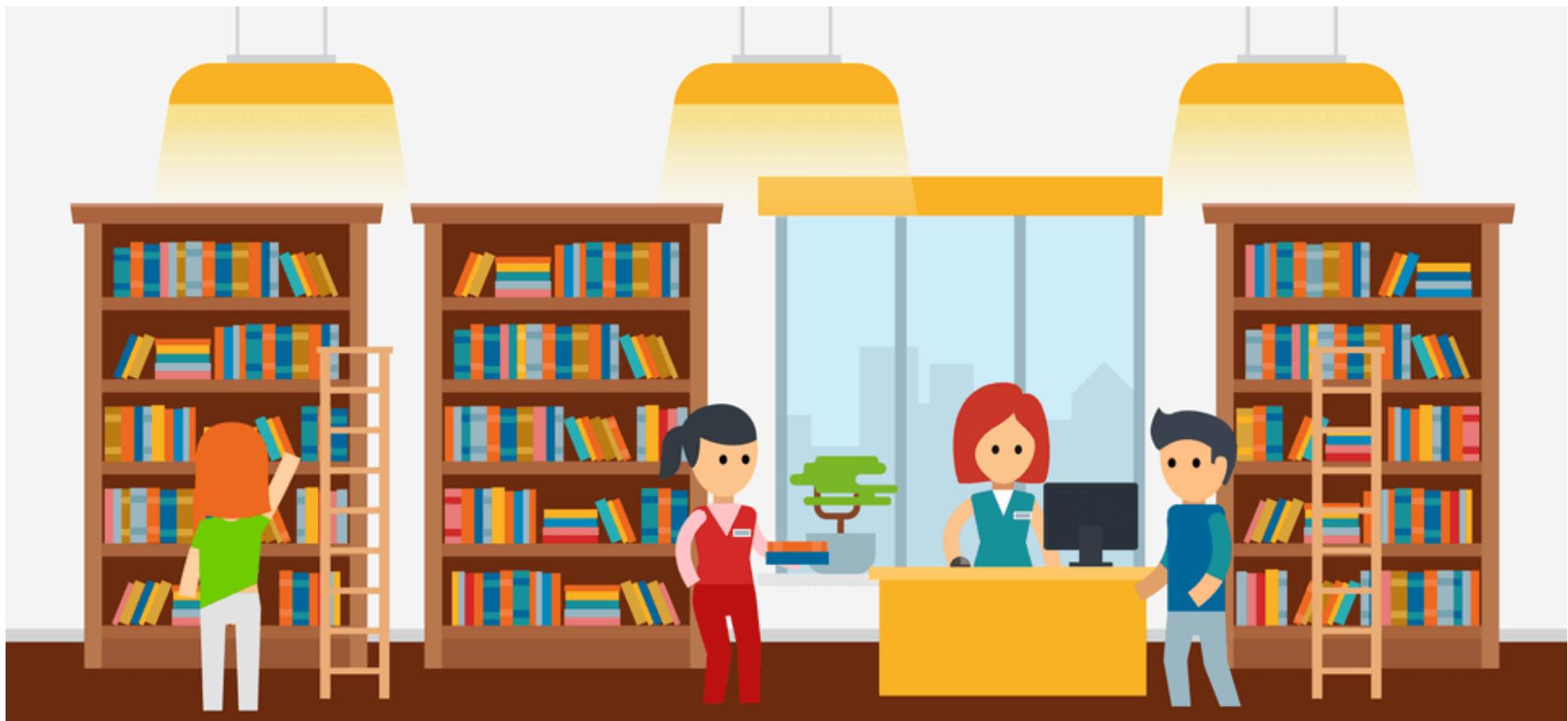
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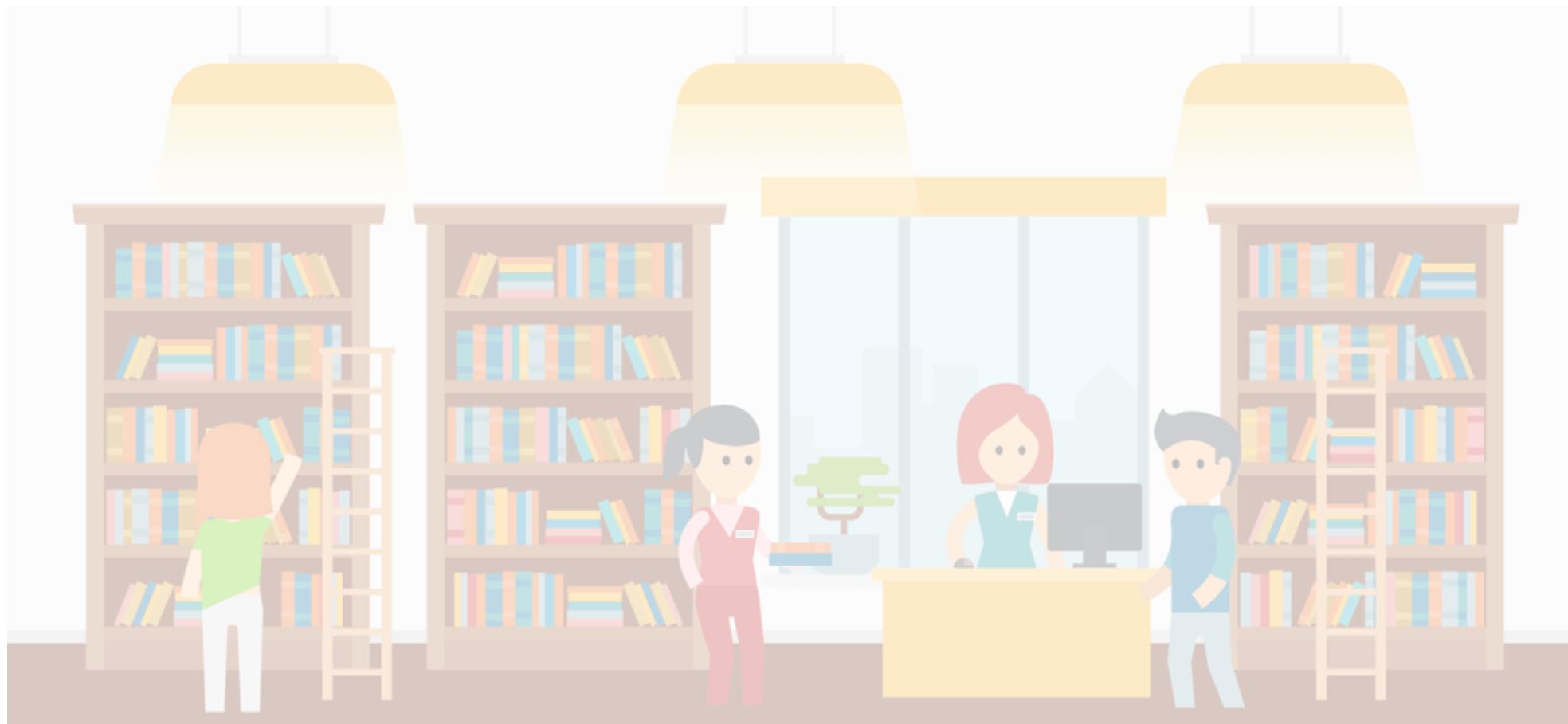
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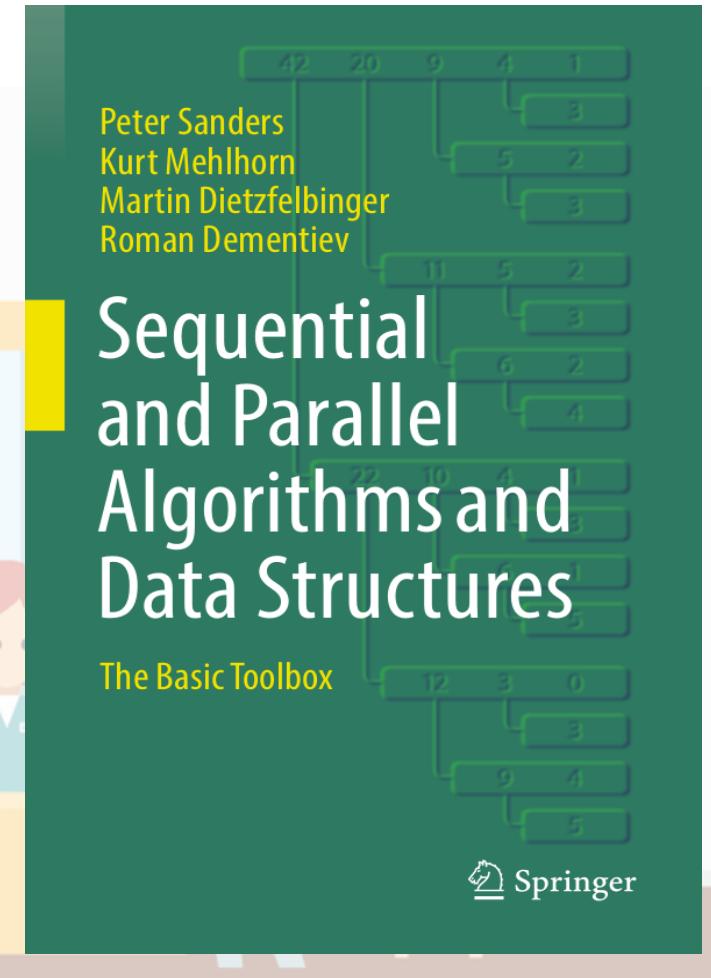
Binary Classification Problem



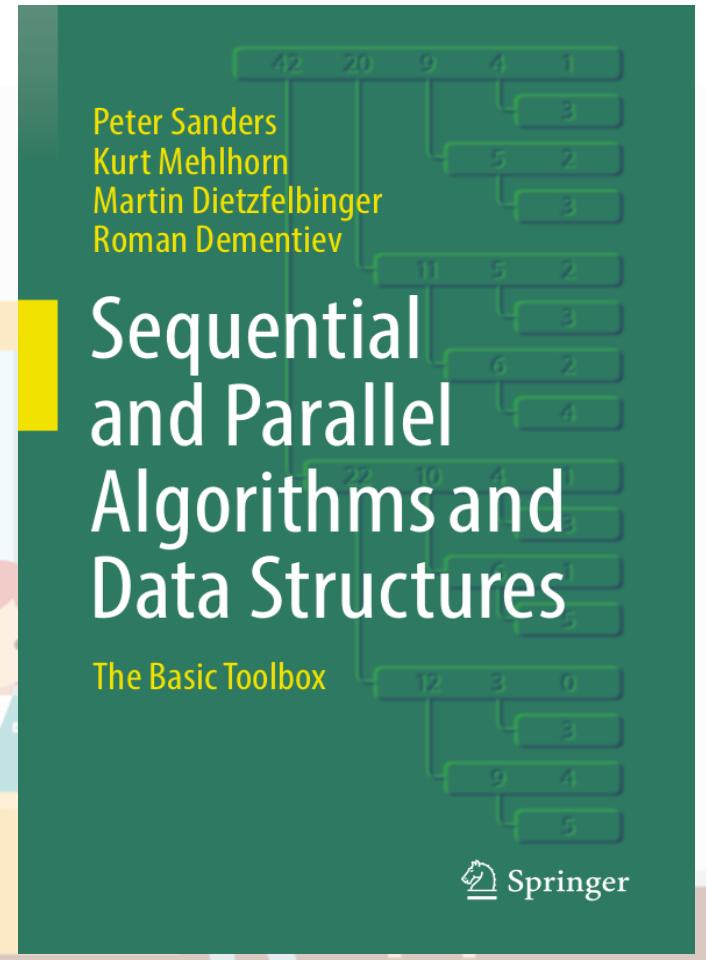
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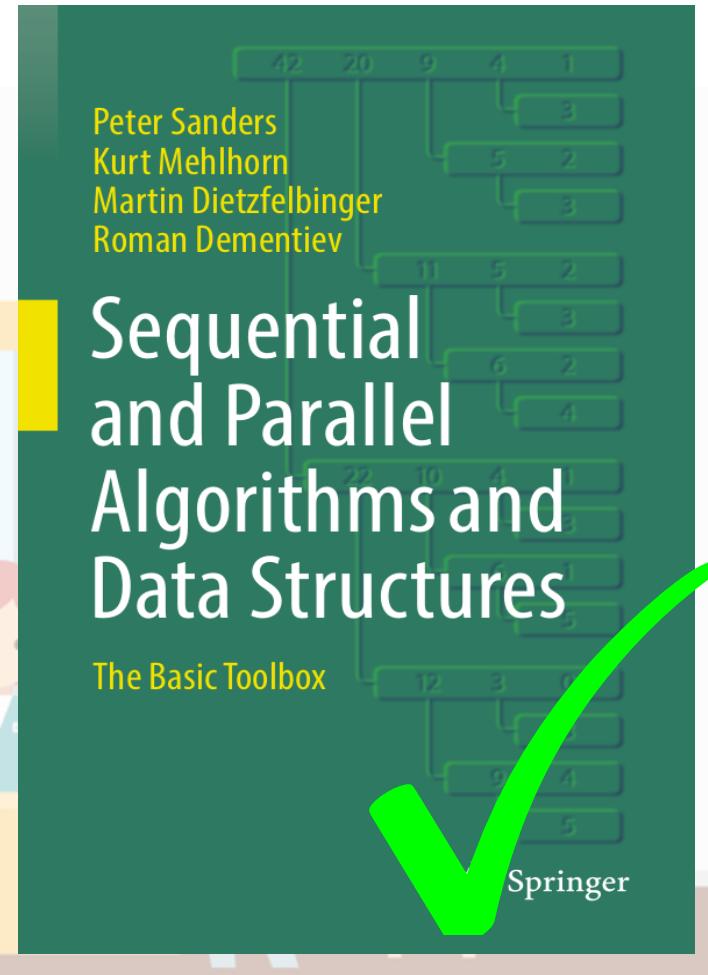
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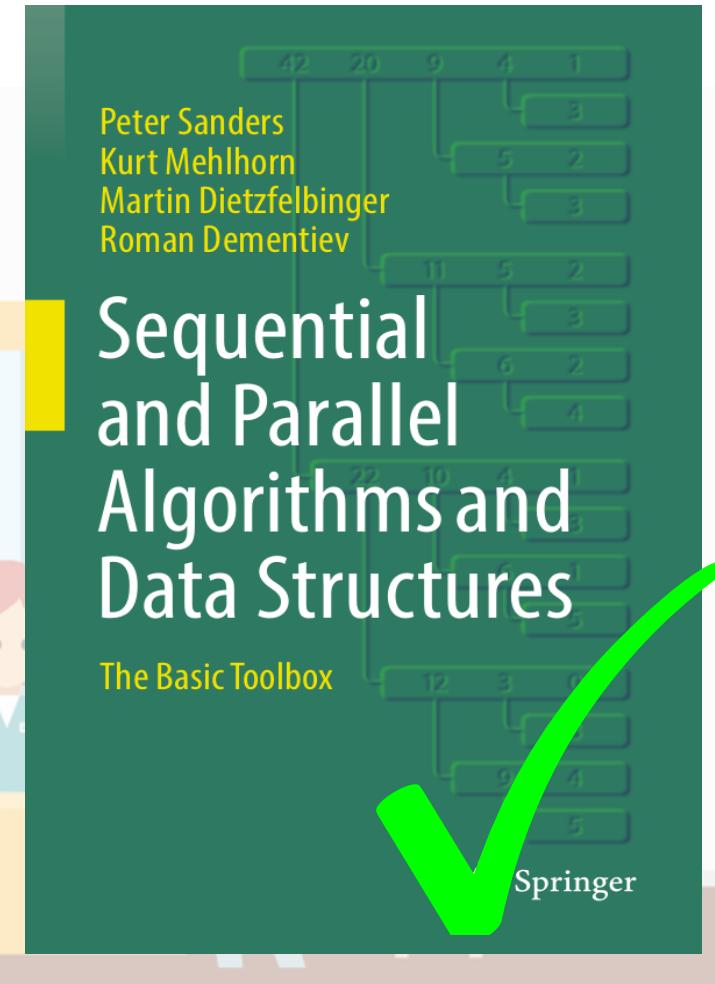
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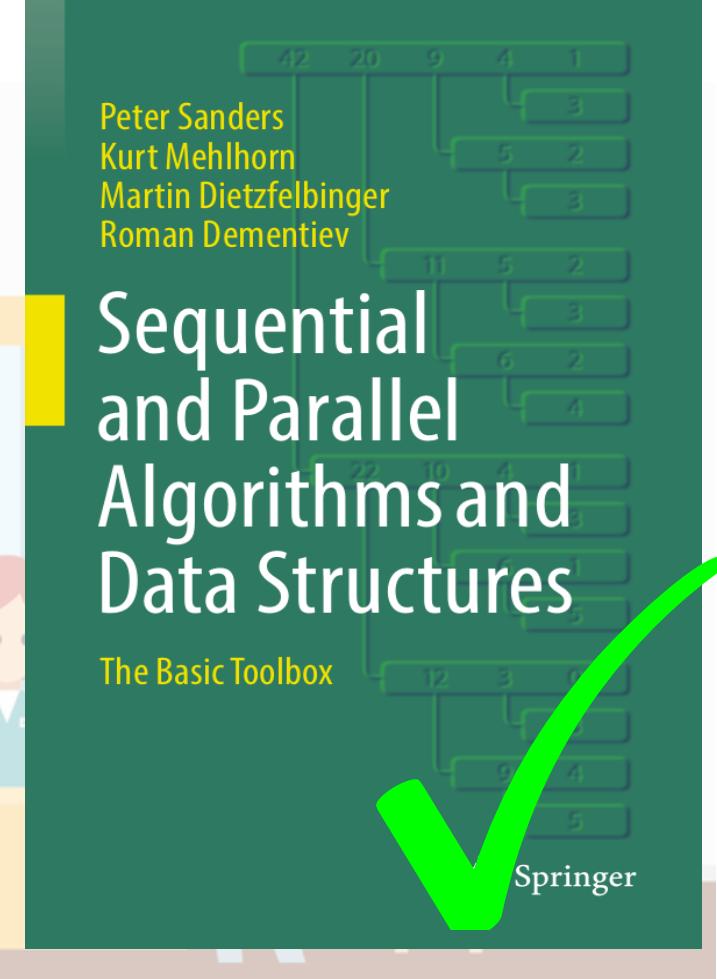
Binary Classification Problem

Very well written

Corollary 12.6. *The approximation ratio of the shortest-queue algorithm is $2 - 1/m$.*

Proof. Let $L_1 = \frac{1}{m} \sum_i t_i$ and $L_2 = \max_i t_i$. The makespan L^* of the optimal solution is at least $\max(L_1, L_2)$. The makespan of the shortest-queue solution is bounded by

$$L_1 + \frac{m-1}{m} L_2 \leq L^* + \frac{m-1}{m} L^* = \left(2 - \frac{1}{m}\right) L^*. \quad \square$$



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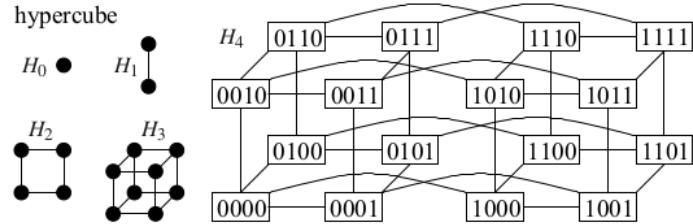
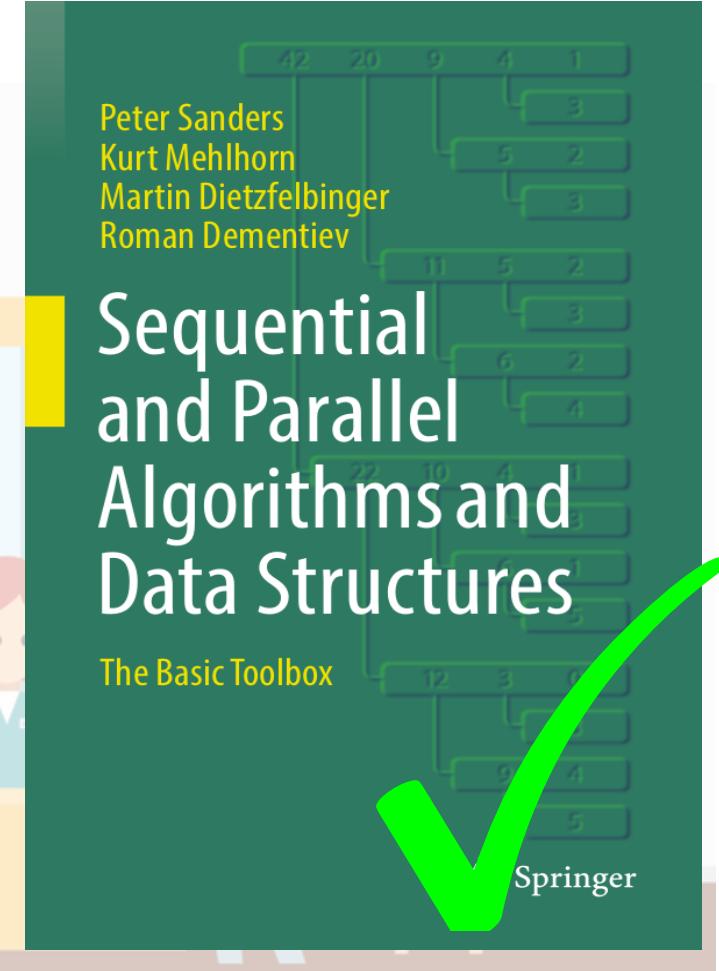


Fig. 13.2. Basic communication topologies for collective communications



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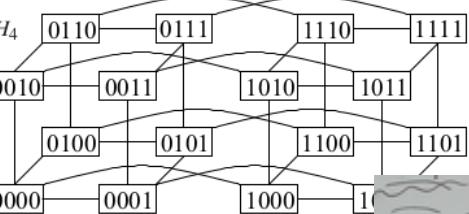
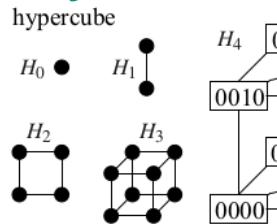
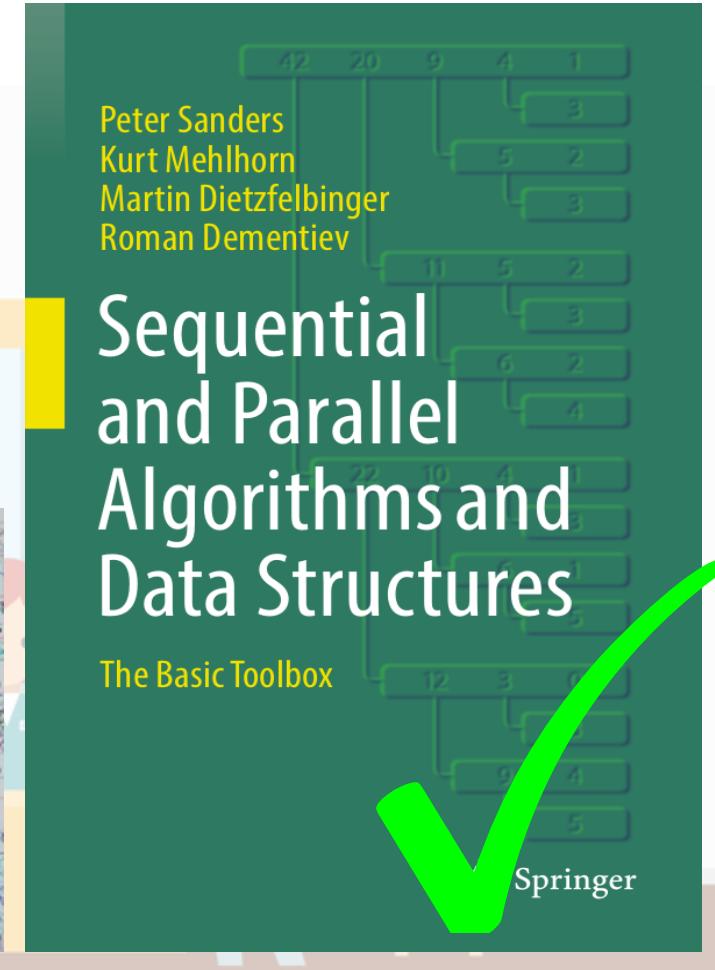


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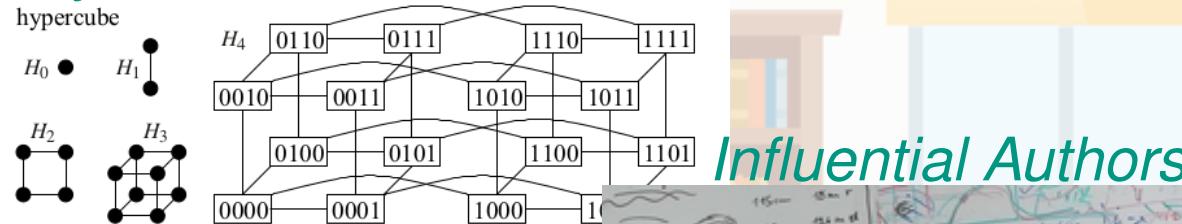
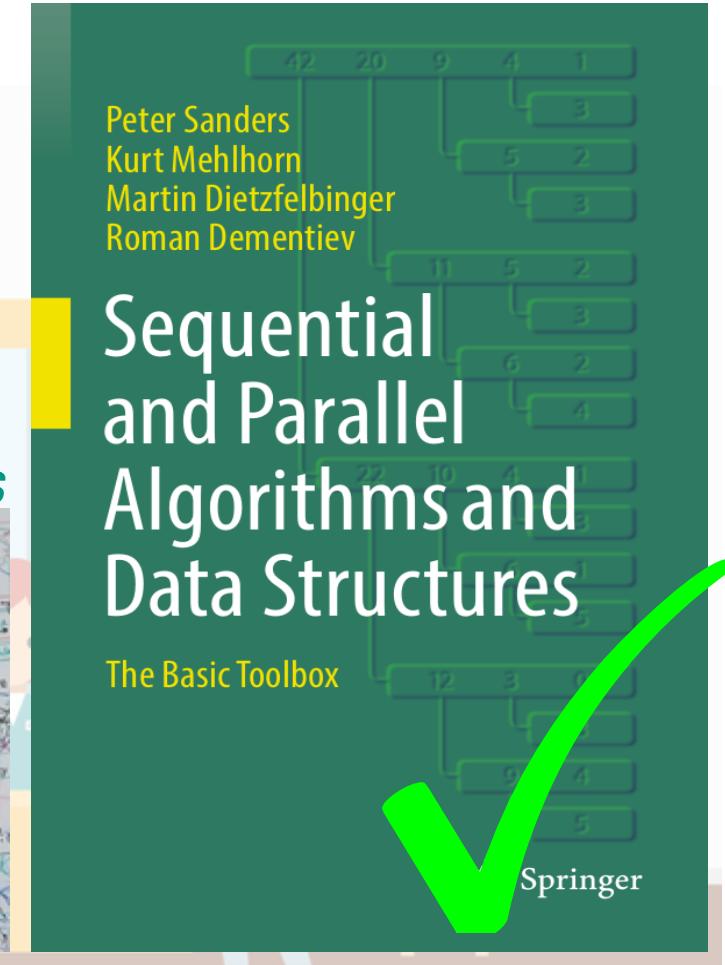
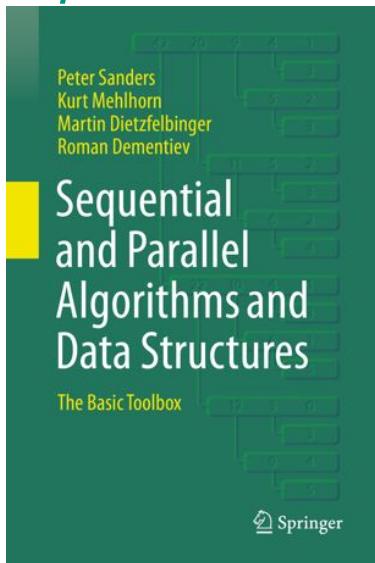


Fig. 13.2. Basic communication topologies for collective communication



Binary Classification Problem

Example of an excellent book



Train classifier on n labeled data points

(x_i, y_i)

Data point $x \in \mathbb{R}^d$

Features:

text patterns, illustration quality,
authors, ...

Label $y \in \{-1, +1\}$

Result:

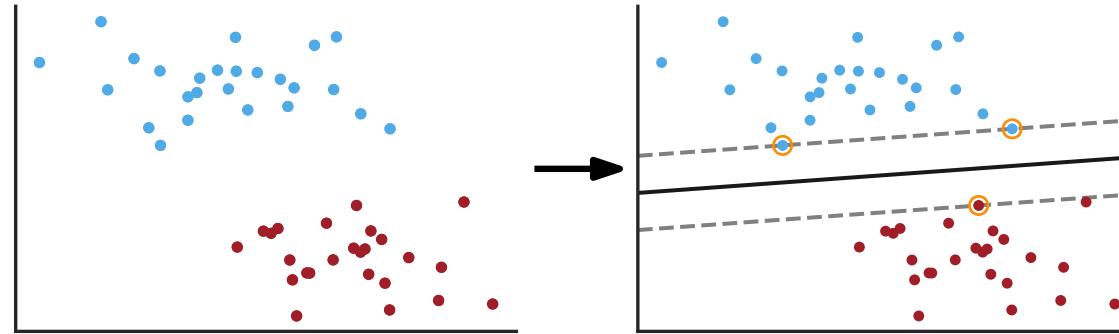
+1: good book
−1: bad book

Goal: Assign label y_{n+1} to new data points x_{n+1}

Support Vector Machines [CV'97]

Find **hyperplane** with **maximum margin** between classes C^- & C^+

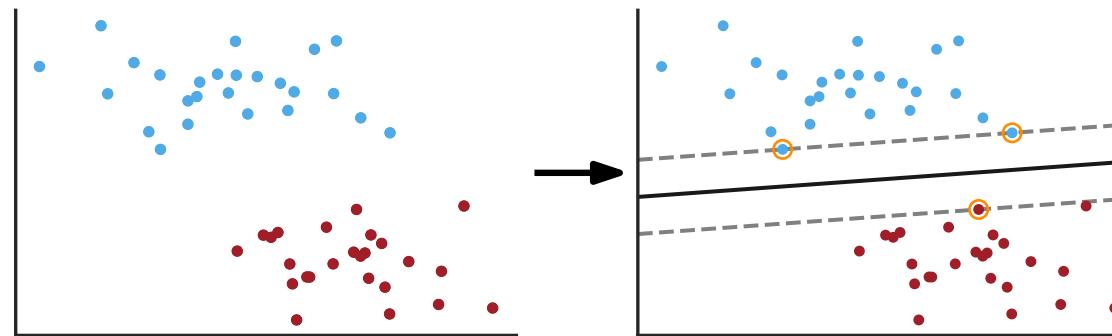
Linear
SVM



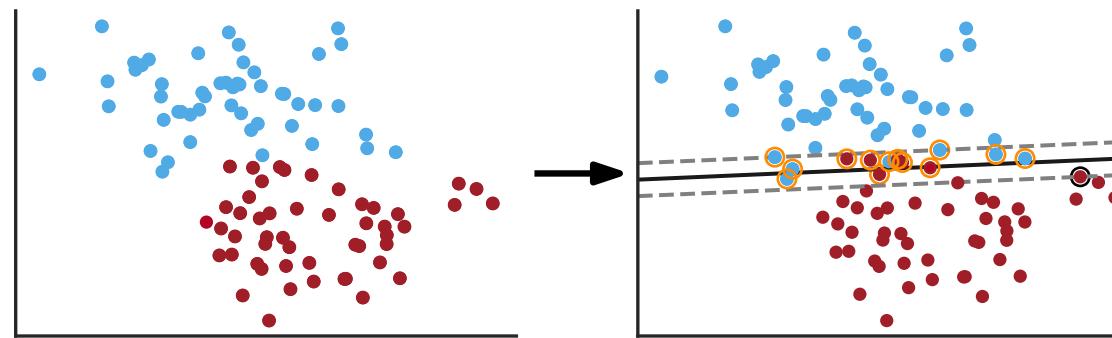
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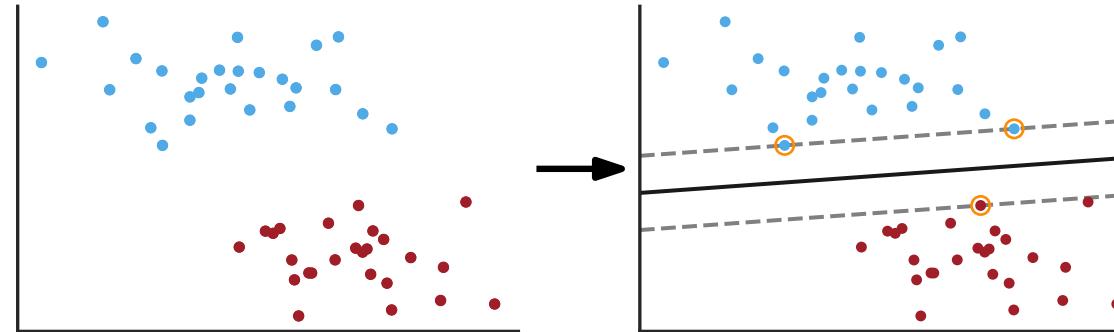
Soft Margin
SVM



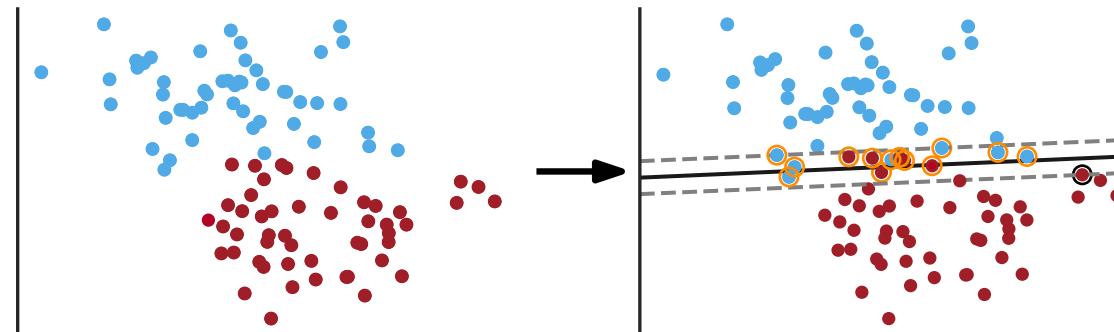
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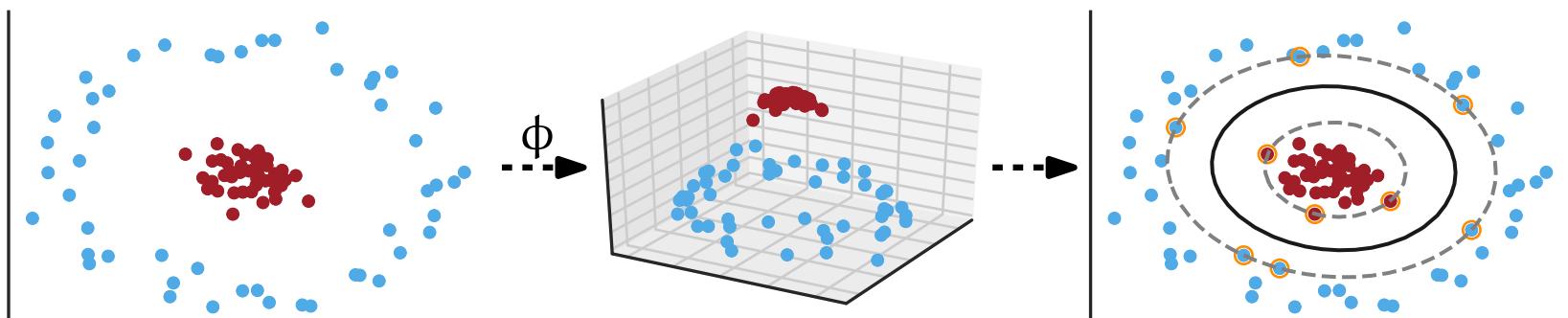
Linear
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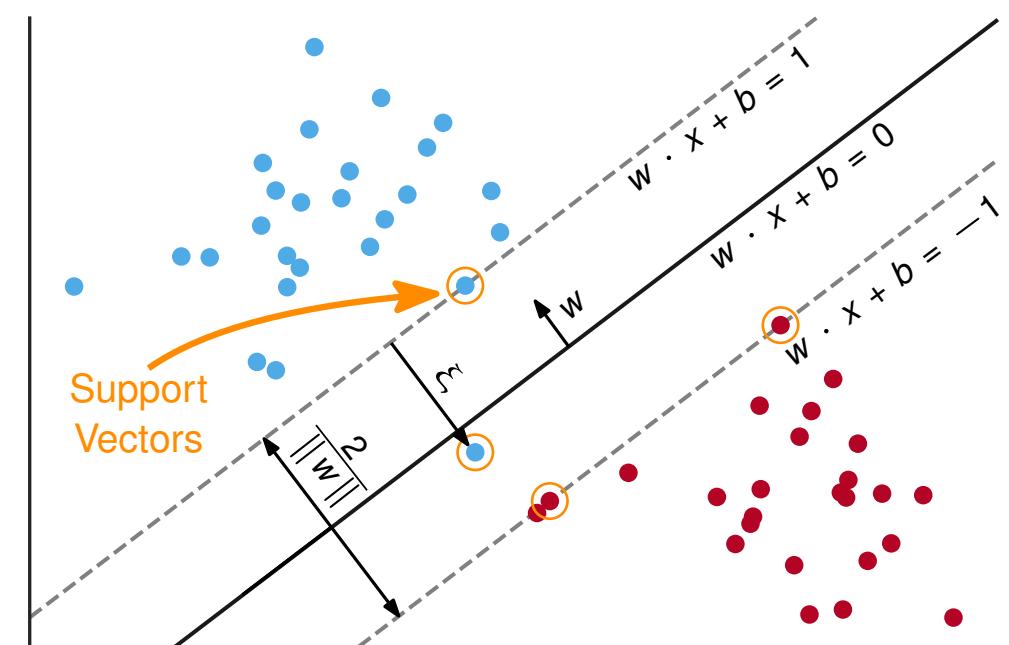
Nonlinear
SVM



Support Vector Machines [CV'97]

SVM Optimization Problem:

$$\begin{aligned} \text{minimize} \quad & \frac{1}{2} \|w\|^2 + C \sum_{i=1}^n \xi_i \\ \text{subject to} \quad & y_i(w \cdot \phi(x_i) + b) \geq 1 - \xi_i, \quad \xi_i \geq 0 \end{aligned}$$

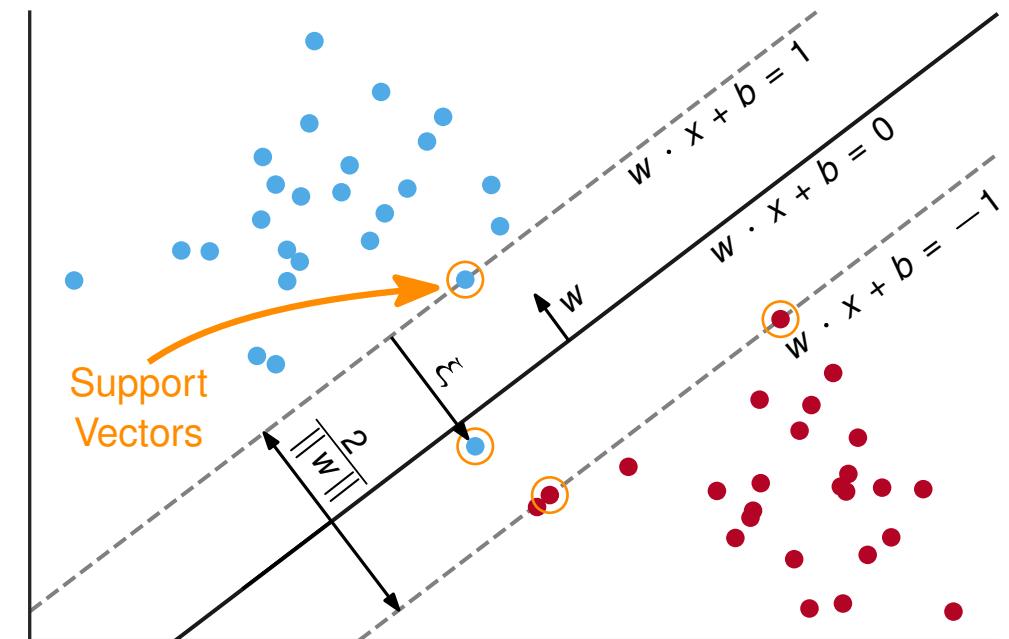


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Slack variables



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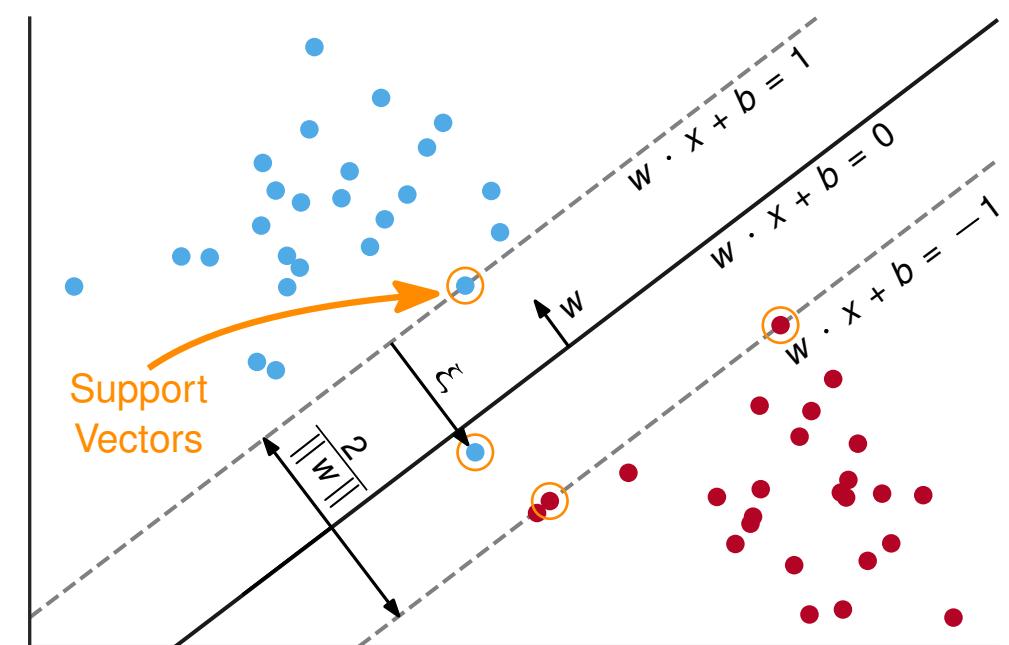
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Mapping to higher dimensional space

$$\phi : \mathbb{R}^d \rightarrow \mathbb{R}^p (d \leq p)$$

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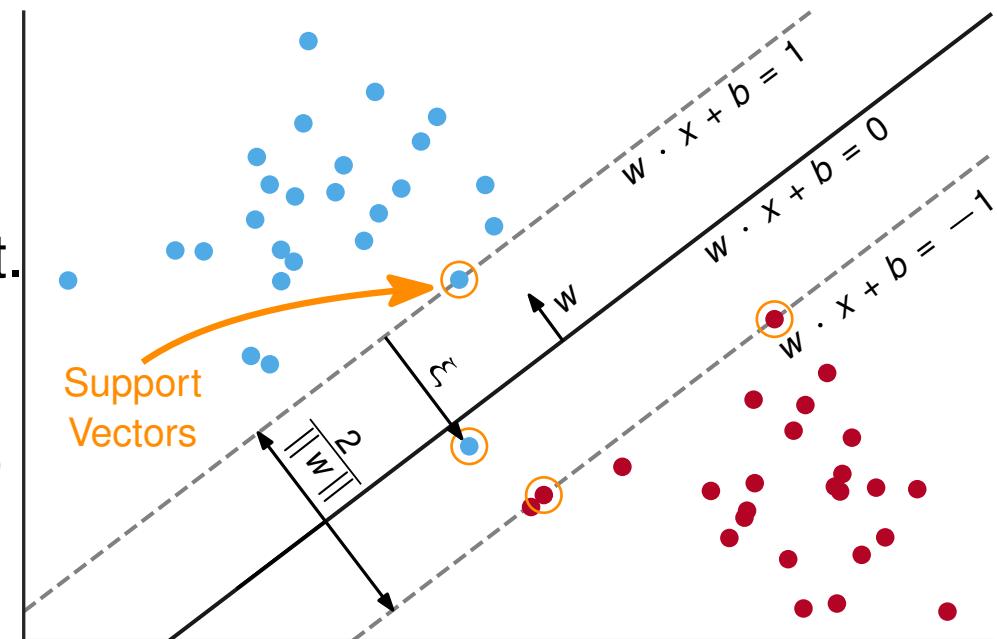
$$\phi : \mathbb{R}^d \rightarrow \mathbb{R}^p (d \leq p)$$

via ...

Kernel Trick:

- Replace inner product by kernel fct.
- $k(x_i, x_j) = \phi(x_i)^T \phi(x_j)$
- Here: $k(x_i, x_j) = \exp(-\gamma \|x_i - x_j\|^2)$

Gaussian kernel (RBF)



Support Vector Machines – Training

Model Selection: instance-specific tuning of several parameters

- Slack penalty C
- Kernel parameters (here: γ)

Complexity:

- Solver running time between $\mathcal{O}(n^2)$ & $\mathcal{O}(n^3)$ [GCBV'04]
- Model selection \leadsto train many models

Support Vector Machines – Training

Model Selection: instance-specific tuning of several parameters

- Slack penalty C
- Kernel parameter (γ , σ)

Training on **large** data sets becomes **infeasible!**

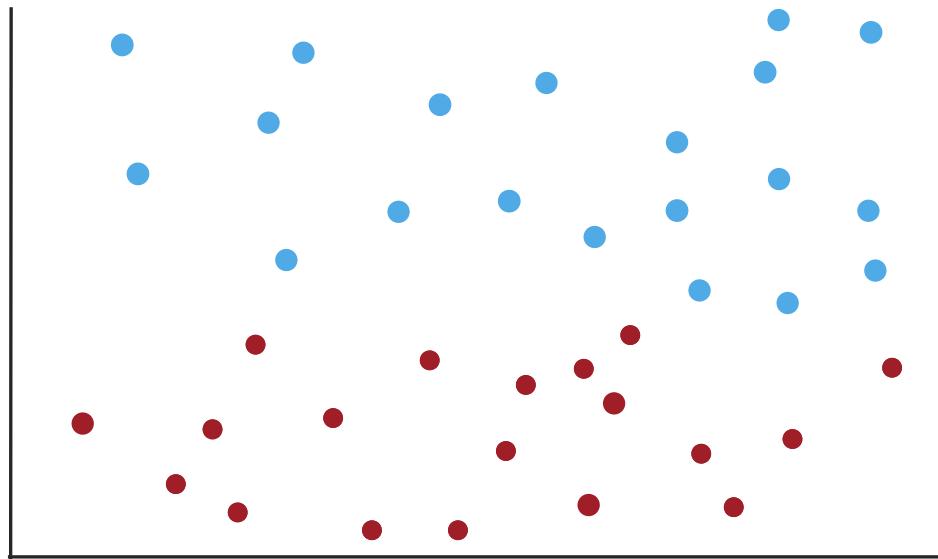
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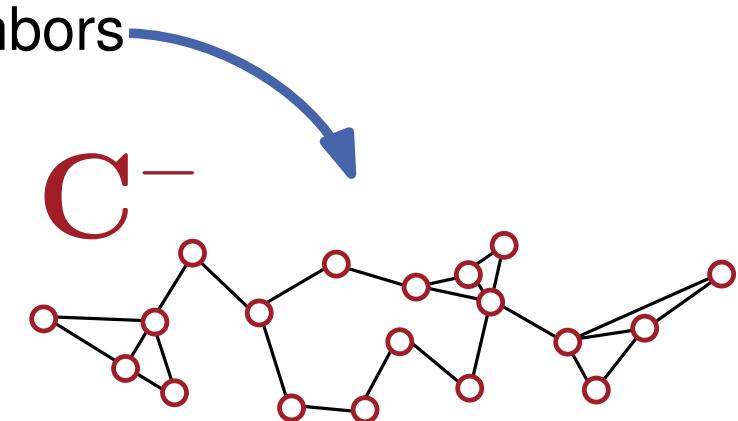
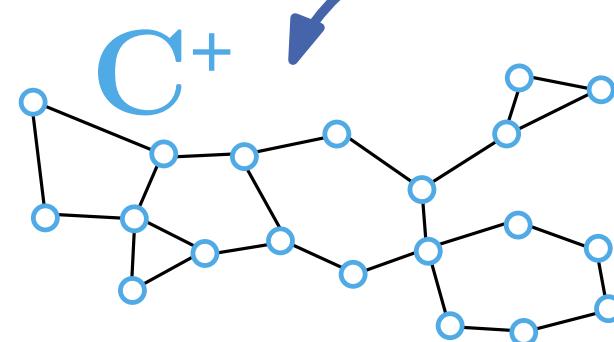
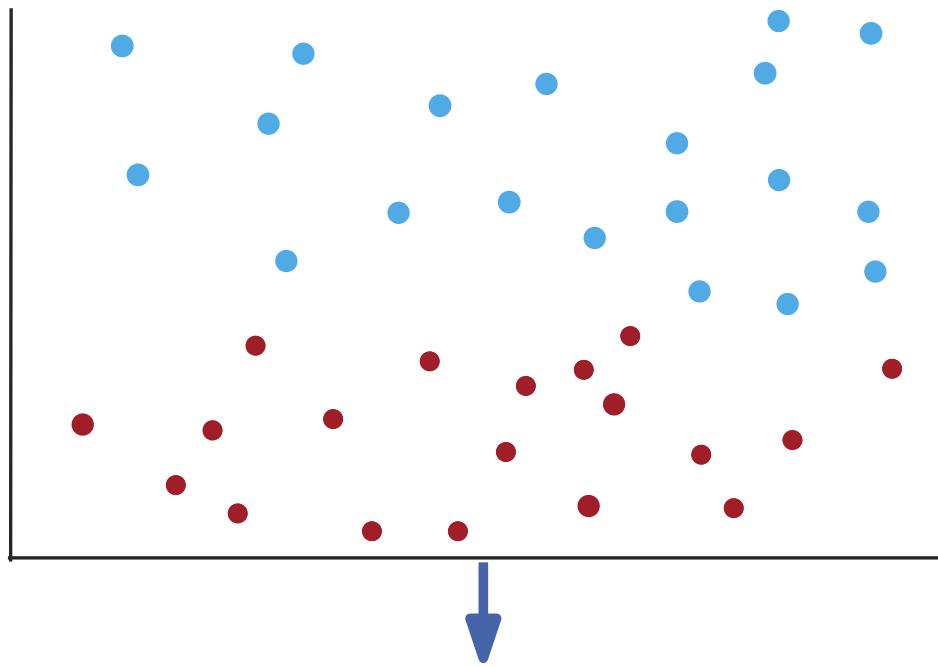
Performance Improvements:

- Sampling
- Parallelization
- Hierarchical techniques
 - Input space
 - Graph representation

From Feature Vectors to Graphs

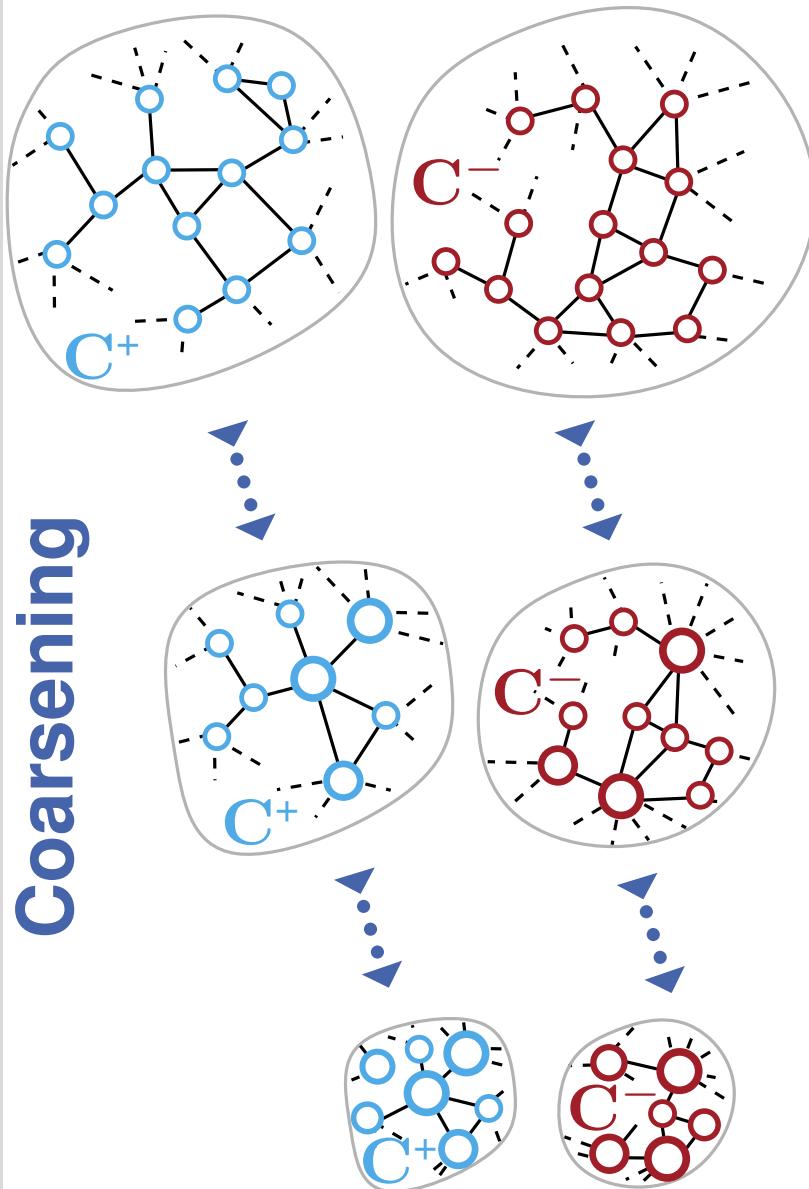


From Feature Vectors to Graphs

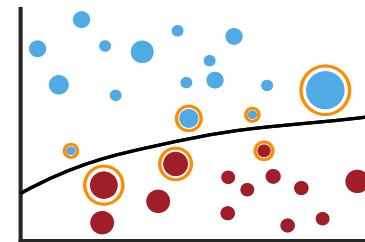
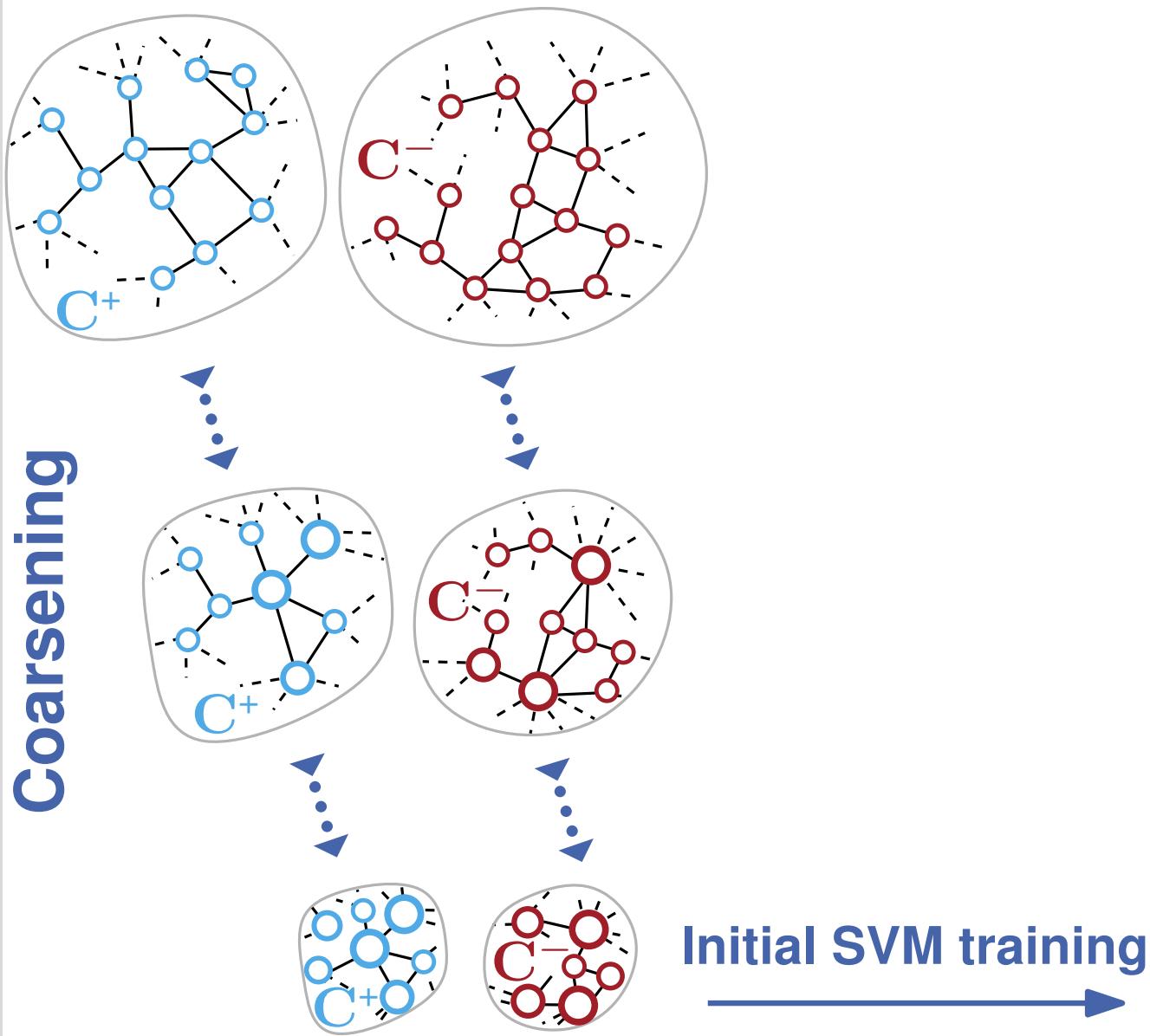


$\omega(e) = \frac{1}{\text{dist}(p,q)}$ \Rightarrow encode **proximity** information into **edge weights**

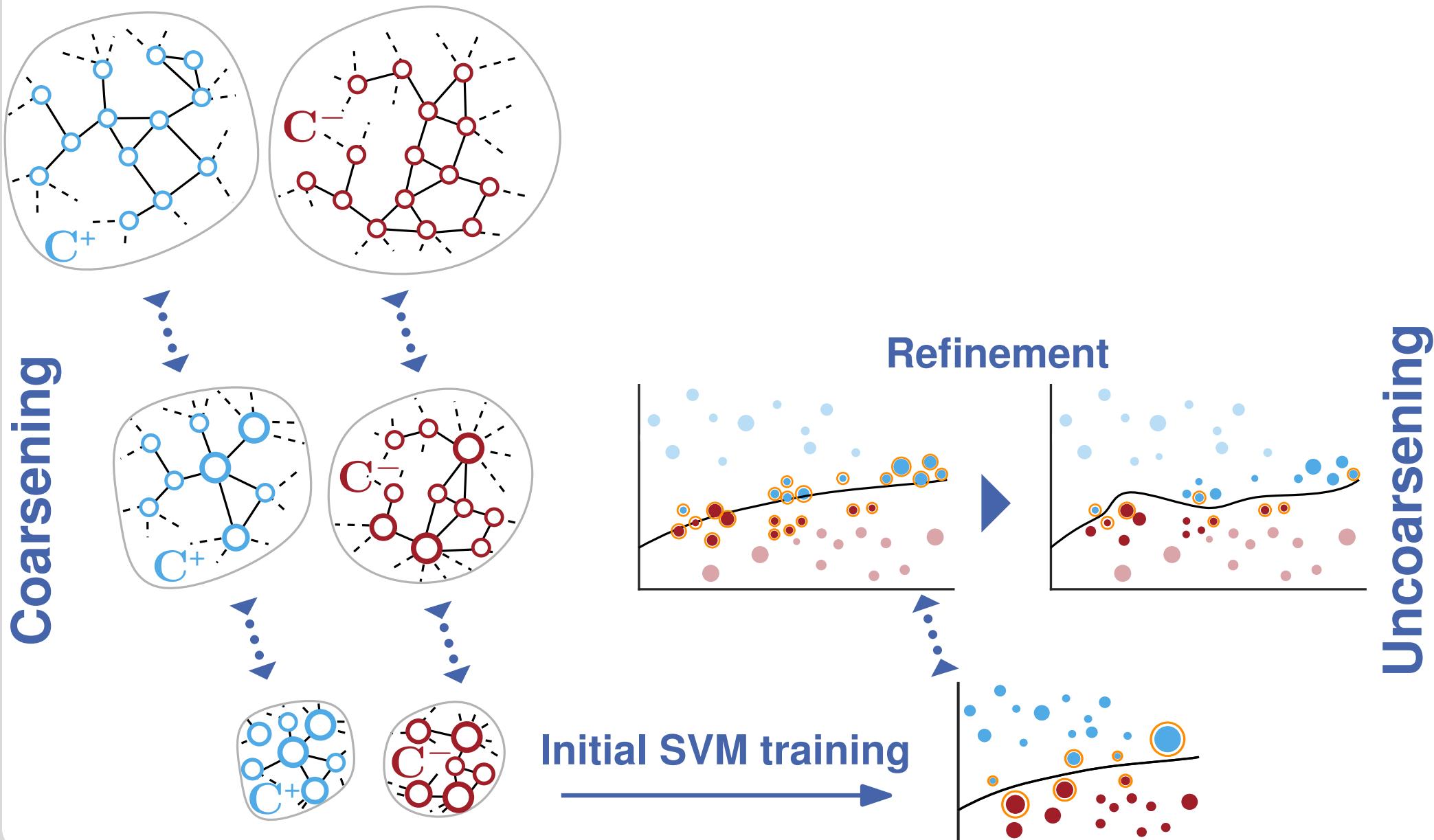
Multilevel Support Vector Machines [RS'15]



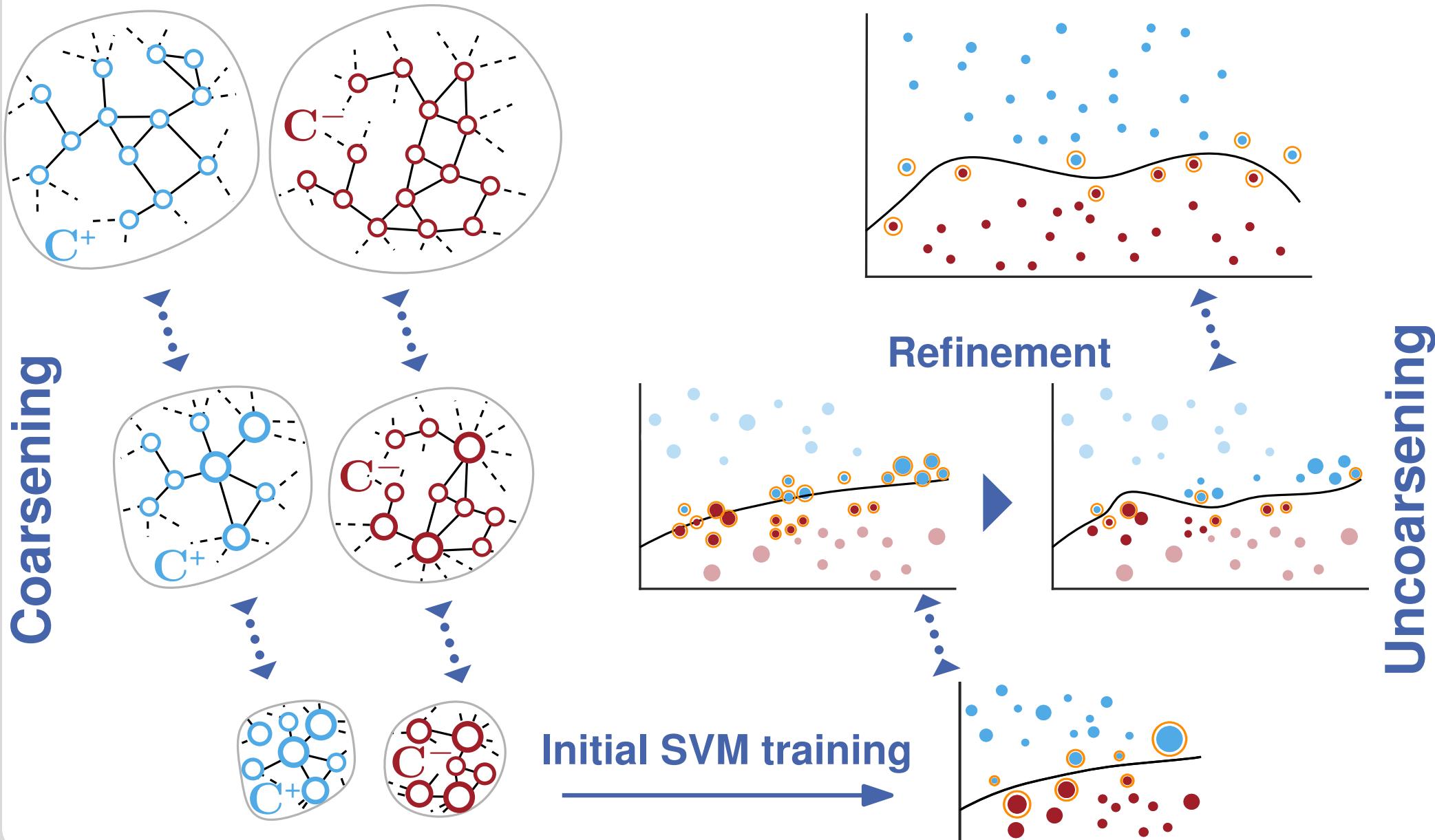
Multilevel Support Vector Machines [RS'15]



Multilevel Support Vector Machines [RS'15]



Multilevel Support Vector Machines [RS'15]



Experimental Setup

Machine: AMD Opteron 6168 with 1.9 GHz, 256 GB of RAM

Implementation:

- approx. k -nearest neighbors: FLANN 1.8.4 [ML'14]
- SVM training: LibSVM 3.22

Configuration:

- $k = 10$ nearest neighbors
- $\ell = 10$ label propagation iterations
- stop coarsening $|C^{+/-}| \approx 500$ nodes

Algorithms:

- KaSVM / KaSVM_{fast}
- mlsvm-AMG (outperforms DC-SVM & EnsembleSVM) [SJKLRS'17]
- LibSVM

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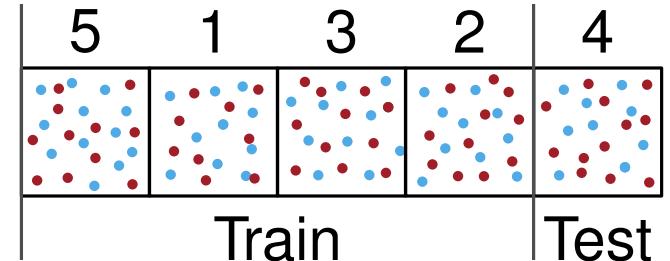
Algorithms:

- KaSVM / KaSVM_{fast}
 - mlsvm-AMG (outperforms DC-SVM & EnsembleSVM) [SJKLRS'17]
 - LibSVM
- initially trained model / no refinement
- 

Experimental Methodology

k-fold cross validation:

- Shuffle data set → split into $k = 5$ parts
- k training repetitions:
 - ⇒ **Training** set: $k-1$ parts
 - ⇒ **Test** set: 1 part
- 5 k-folds per instance



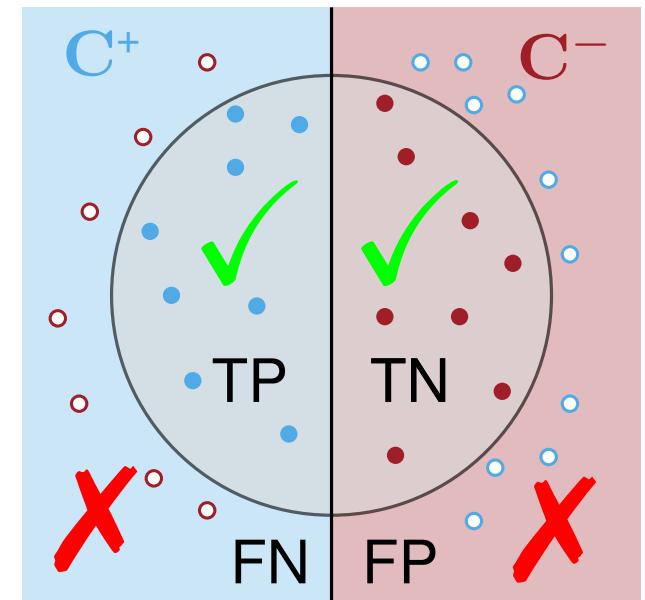
Performance Measures:

$$\text{Accuracy (ACC)} = \frac{TP + TN}{FP + TN + TP + FN}$$

$$\text{Sensitivity (SN)} = \frac{TP}{TP + FN}$$

$$\text{Specificity (SP)} = \frac{TN}{TN + FP}$$

$$\text{Geometric mean} = \sqrt{SP \cdot SN}$$



Running Time

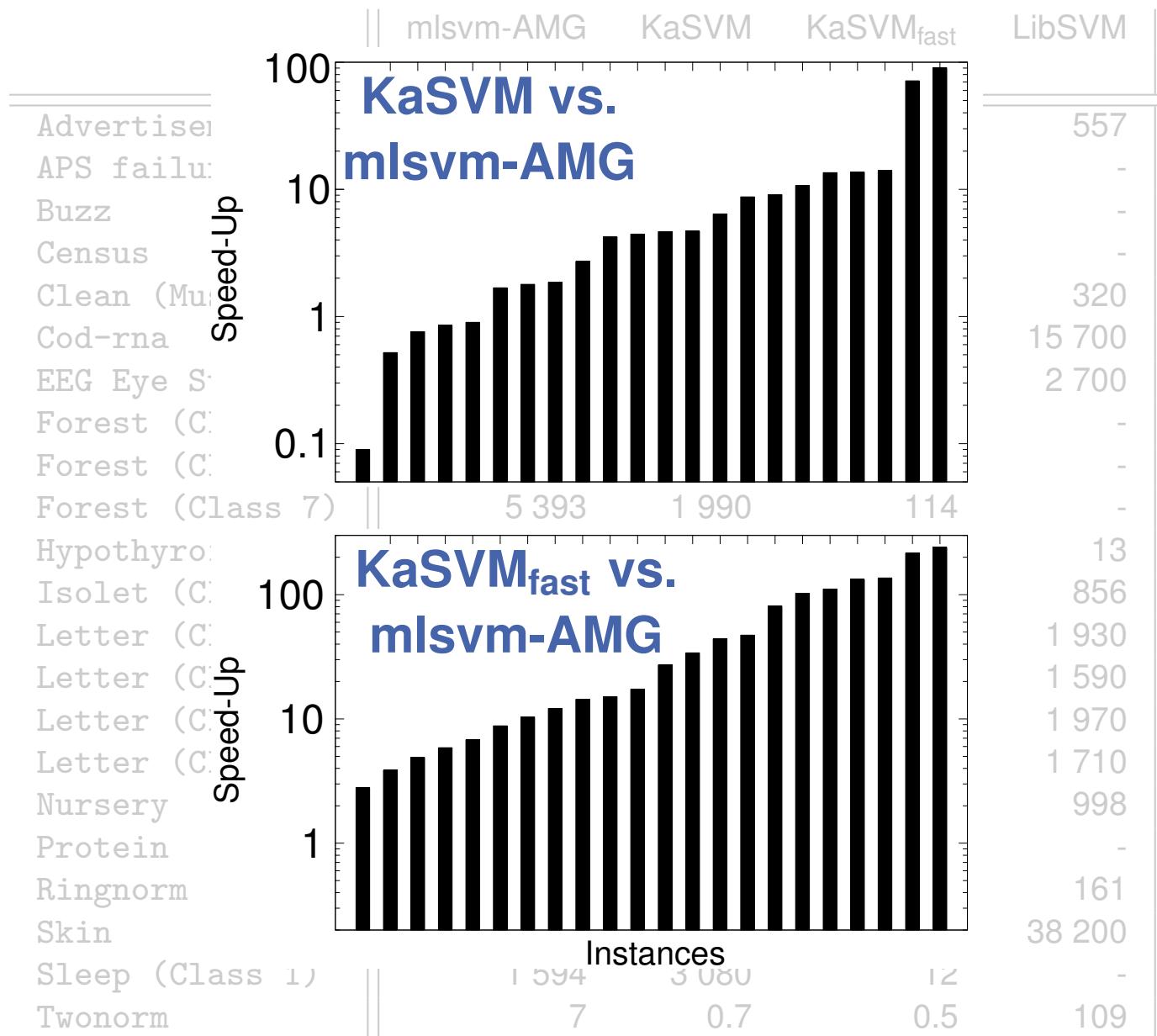
	mlsvm-AMG	KaSVM	KaSVM _{fast}	LibSVM
	running time [s]			
Advertisement	343	192	70	557
APS failure	1 473	109	13	-
Buzz	110	121	19	-
Census	3 047	657	38	-
Clean (Musk)	14	8	4	320
Cod-rna	80	43	7	15 700
EEG Eye State	123	1 320	0.9	2 700
Forest (Class 3)	10 156	744	99	-
Forest (Class 5)	6 986	1 090	158	-
Forest (Class 7)	5 393	1 990	114	-
Hypothyroid	2	3	0.9	13
Isolet (Class A)	1 627	23	7	856
Letter (Class A)	17	4	2	1 930
Letter (Class B)	55	4	2	1 590
Letter (Class H)	74	9	2	1 970
Letter (Class Z)	31	3	2	1 710
Nursery	7	2	0.7	998
Protein	3 654	41	17	-
Ringnorm	10	13	0.6	161
Skin	81	18	12	38 200
Sleep (Class 1)	1 594	3 080	12	-
Twonorm	7	0.7	0.5	109

Running Time

	mlsvm-AMG	KaSVM	KaSVM _{fast} running time [s]	LibSVM
Advertisement	343	192	70	557
APS failure	1 473	109	13	■
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Twonorm	7	0.7	0.5	109

Running
time
 $> 24h$

Running Time



Classification Quality

Instance	mlsvm-AMG G-mean	KaSVM G-mean	KaSVM _{fast} G-mean	LibSVM G-mean
Advertisement	0.91	0.91	0.87	0.941
APS failure	0.94	0.92	0.94	-
Buzz	0.95	0.94	0.93	-
Census	0.81	0.83	0.81	-
Clean (Musk)	0.88	0.96	0.88	0.996
Cod-rna	0.94	0.94	0.93	0.855
EEG Eye State	0.75	0.83	0.63	0.929
Forest (Class 3)	0.94	0.95	0.94	-
Forest (Class 5)	0.79	0.90	0.90	-
Forest (Class 7)	0.86	0.93	0.91	-
Hypothyroid	0.94	0.90	0.94	0.927
Isolet (Class A)	0.00	0.99	0.88	0.990
Letter (Class A)	0.94	0.96	0.95	0.995
Letter (Class B)	0.88	0.92	0.92	0.979
Letter (Class H)	0.76	0.89	0.86	0.970
Letter (Class Z)	0.92	0.95	0.95	0.991
Nursery	1.00	1.00	1.00	1.000
Protein	0.92	0.93	0.91	-
Ringnorm	0.98	0.97	0.82	0.987
Skin	0.99	1.00	1.00	1.000
Sleep (Class 1)	0.69	0.70	0.41	-
Twonorm	0.97	0.96	0.96	0.981

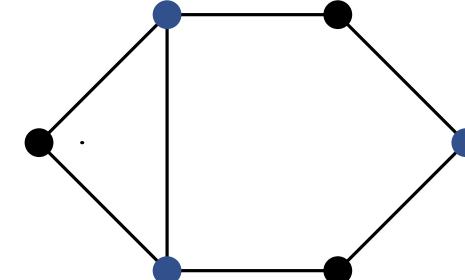
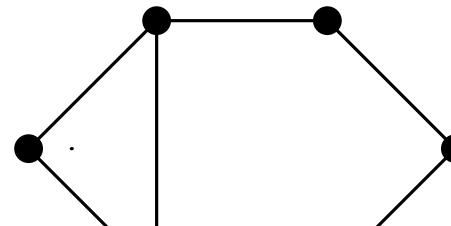
PACE Challenge 2019

Vertex Cover Track

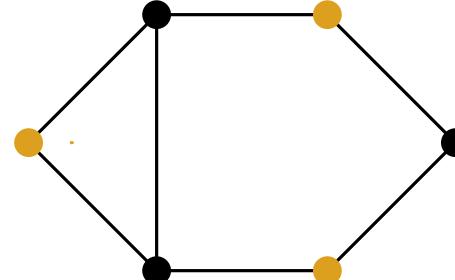
D. Hespe, S. Lamm, C. Schulz, D. Strash

Vertex Cover and Complementary Problems

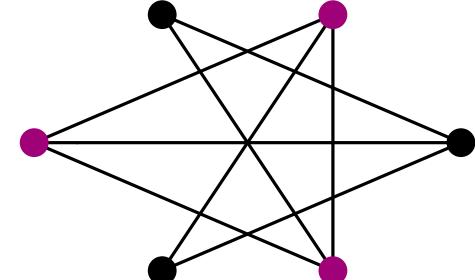
Input graph



Vertex cover



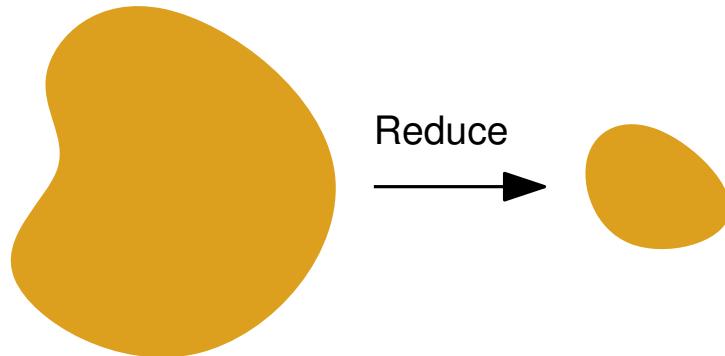
Independent Set



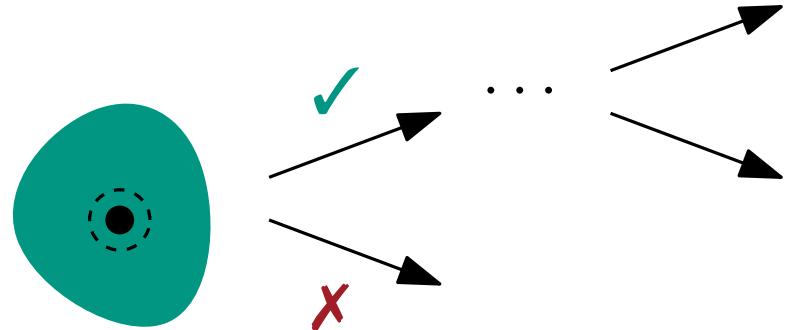
Clique

Techniques

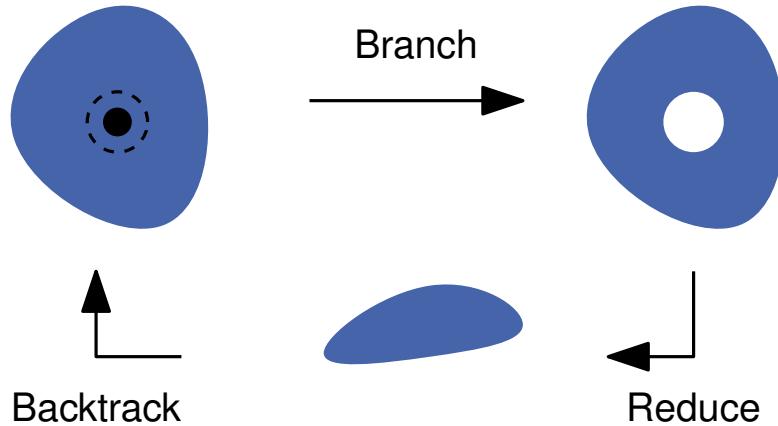
Kernelization



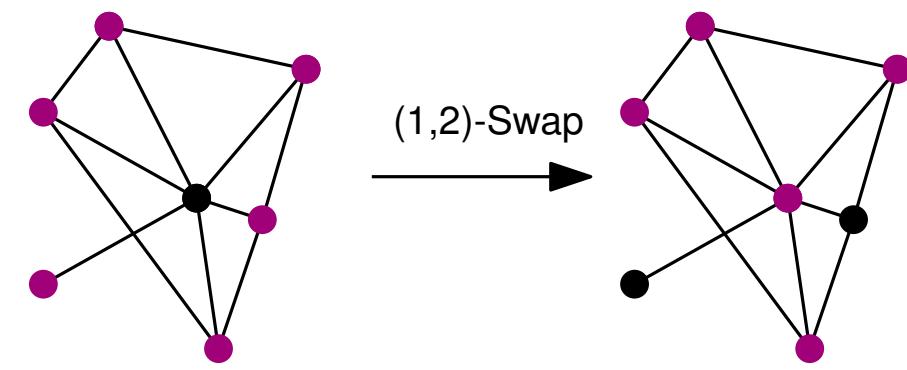
Branch-and-Bound



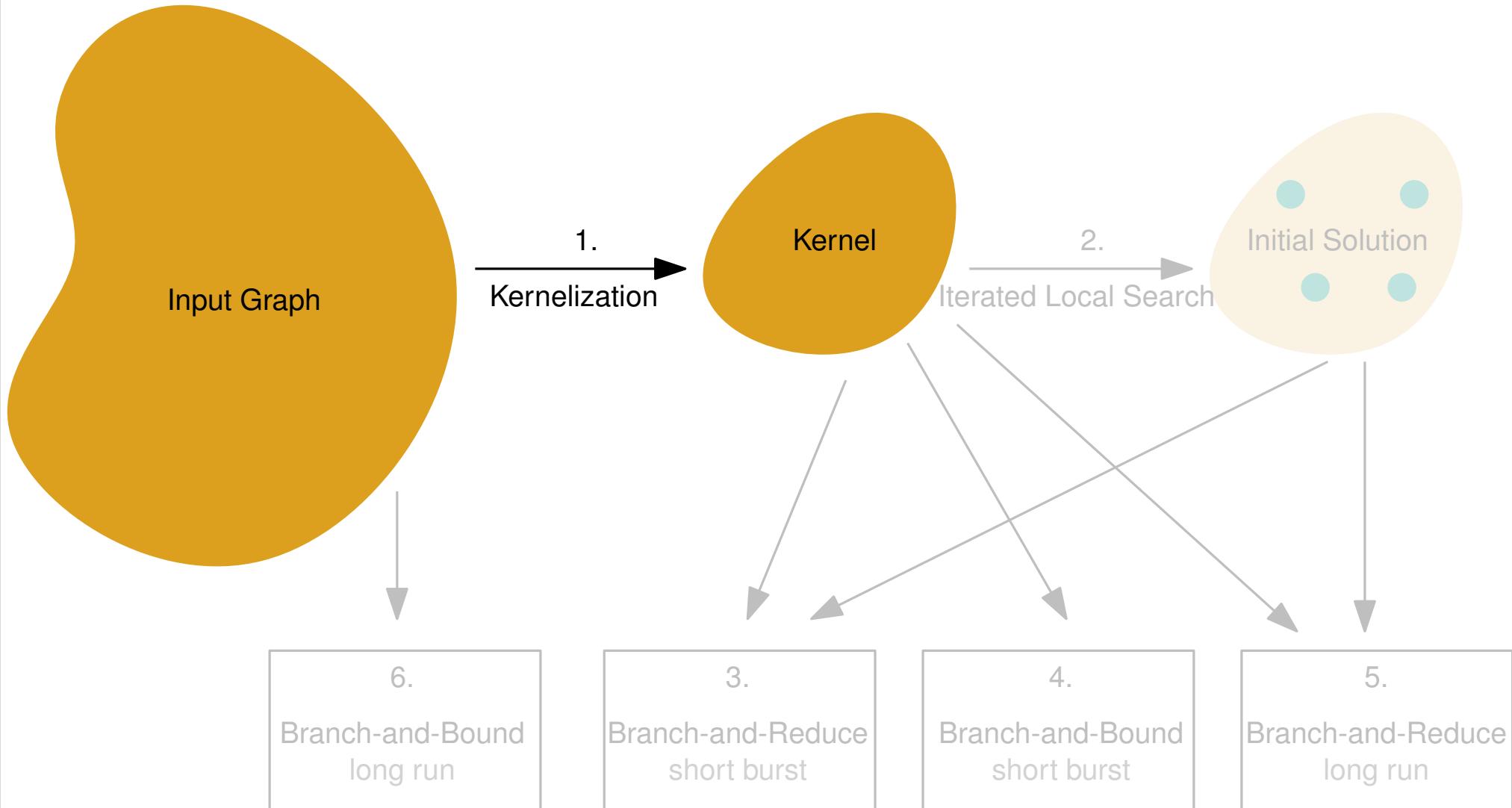
Branch-and-Reduce



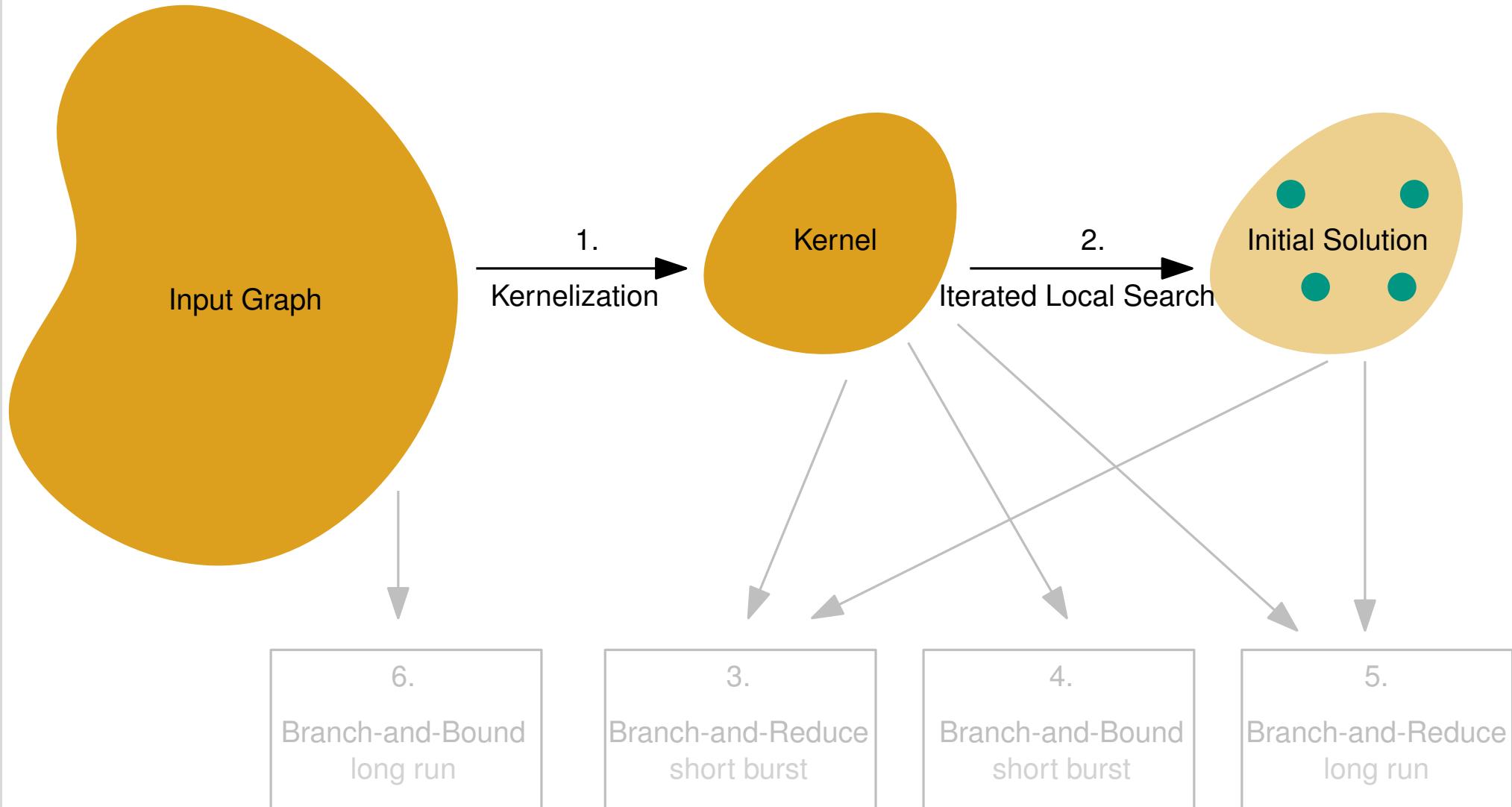
Iterated Local Search



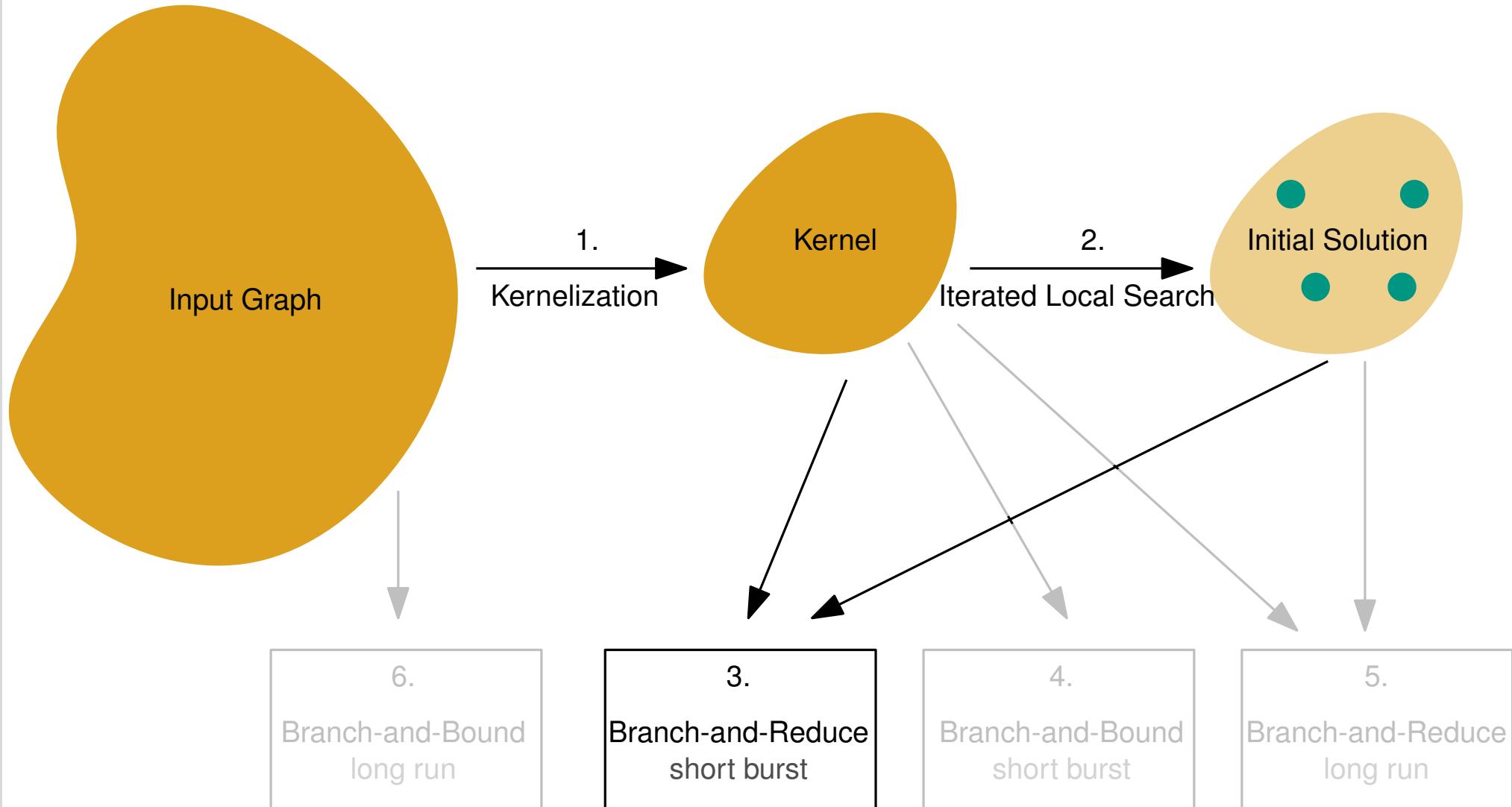
Algorithm Overview



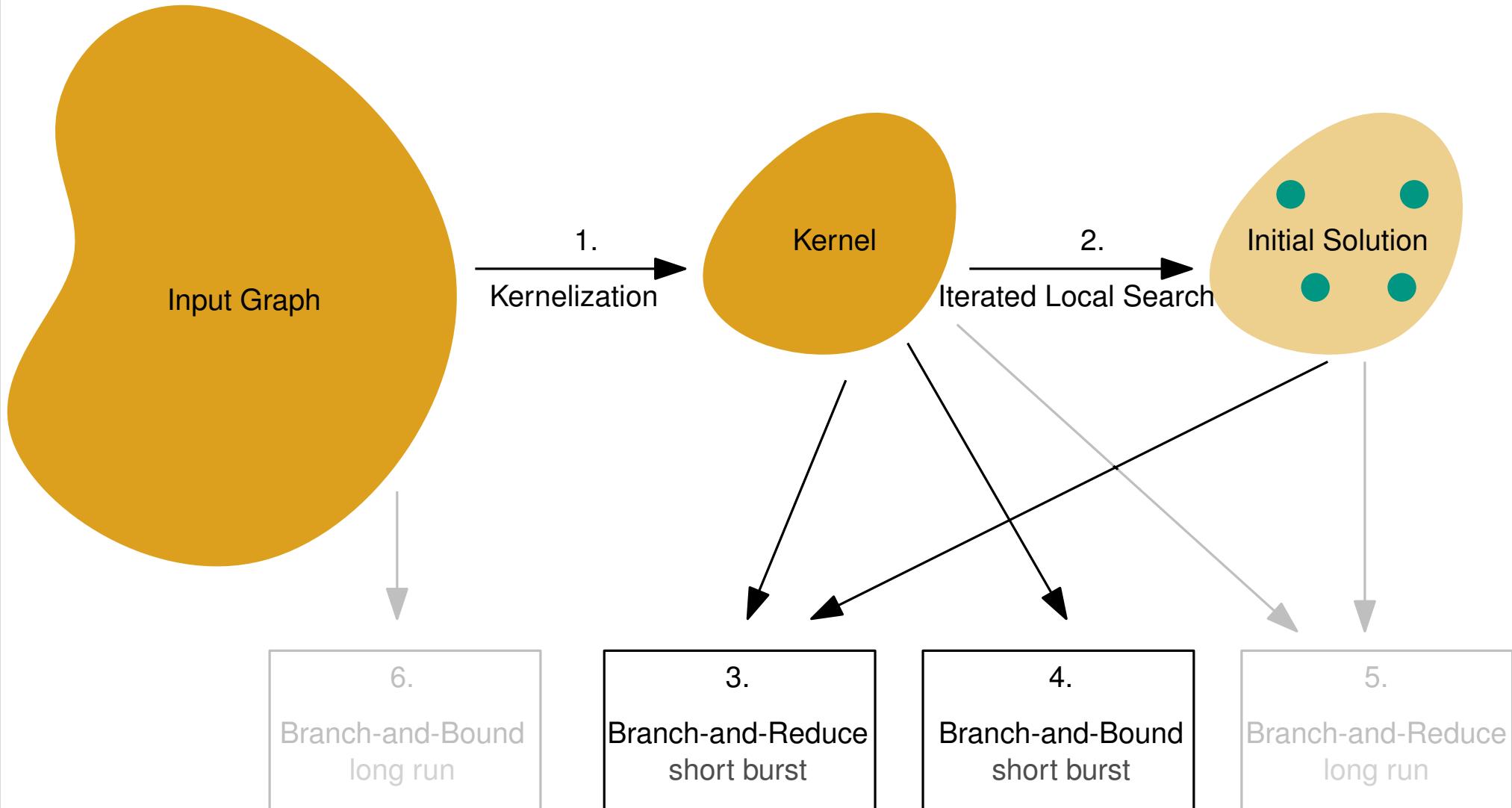
Algorithm Overview



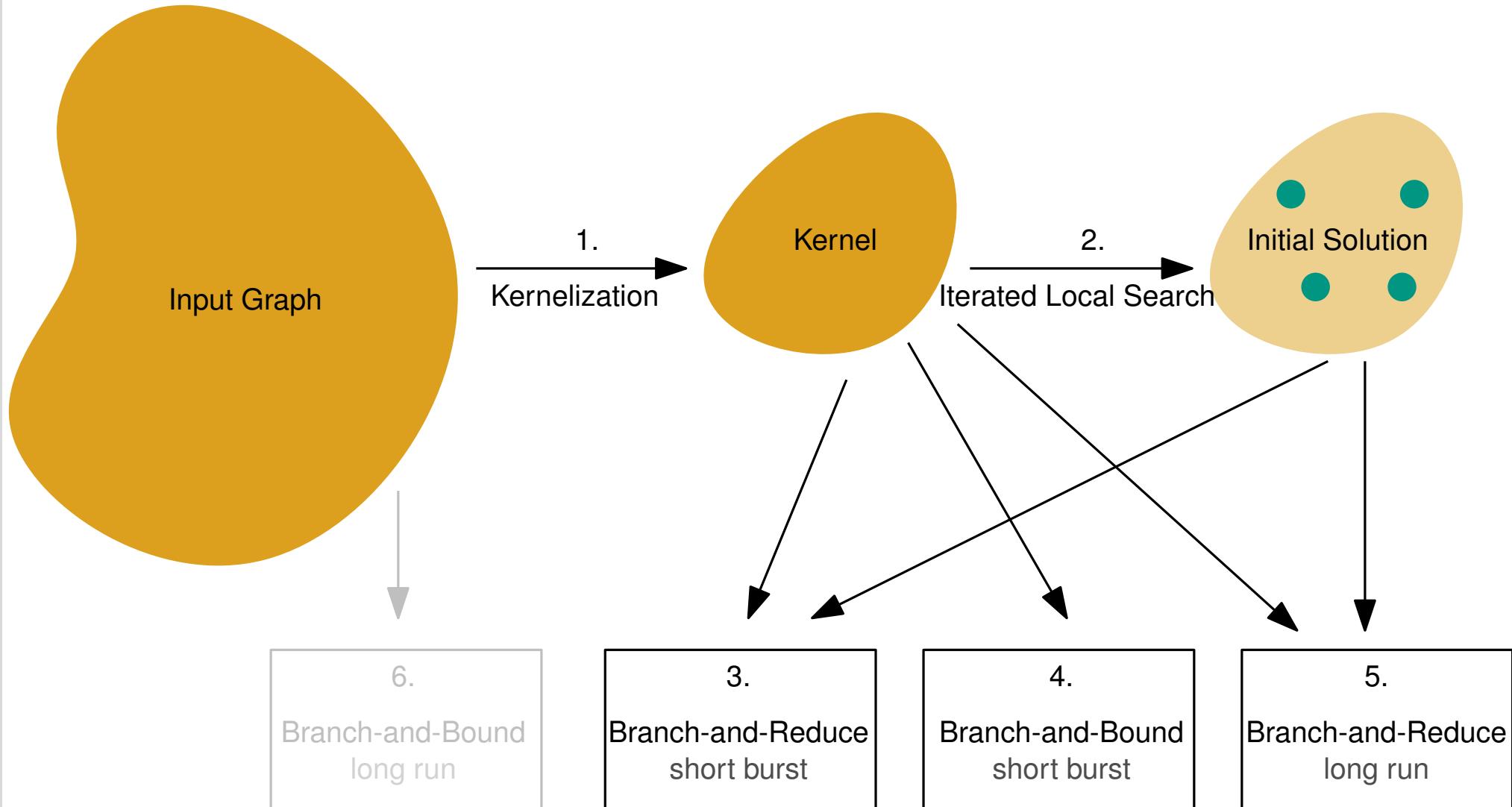
Algorithm Overview



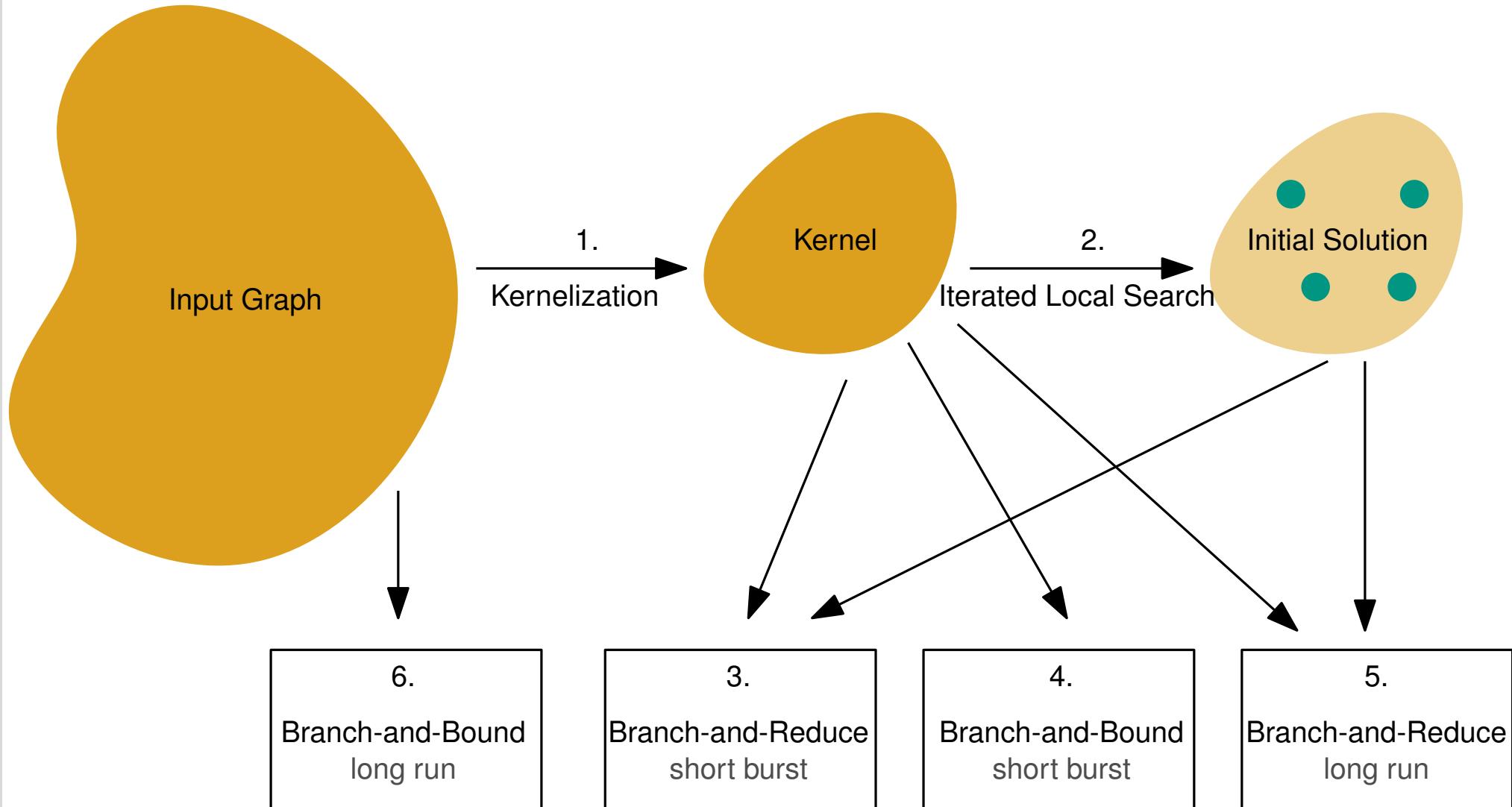
Algorithm Overview



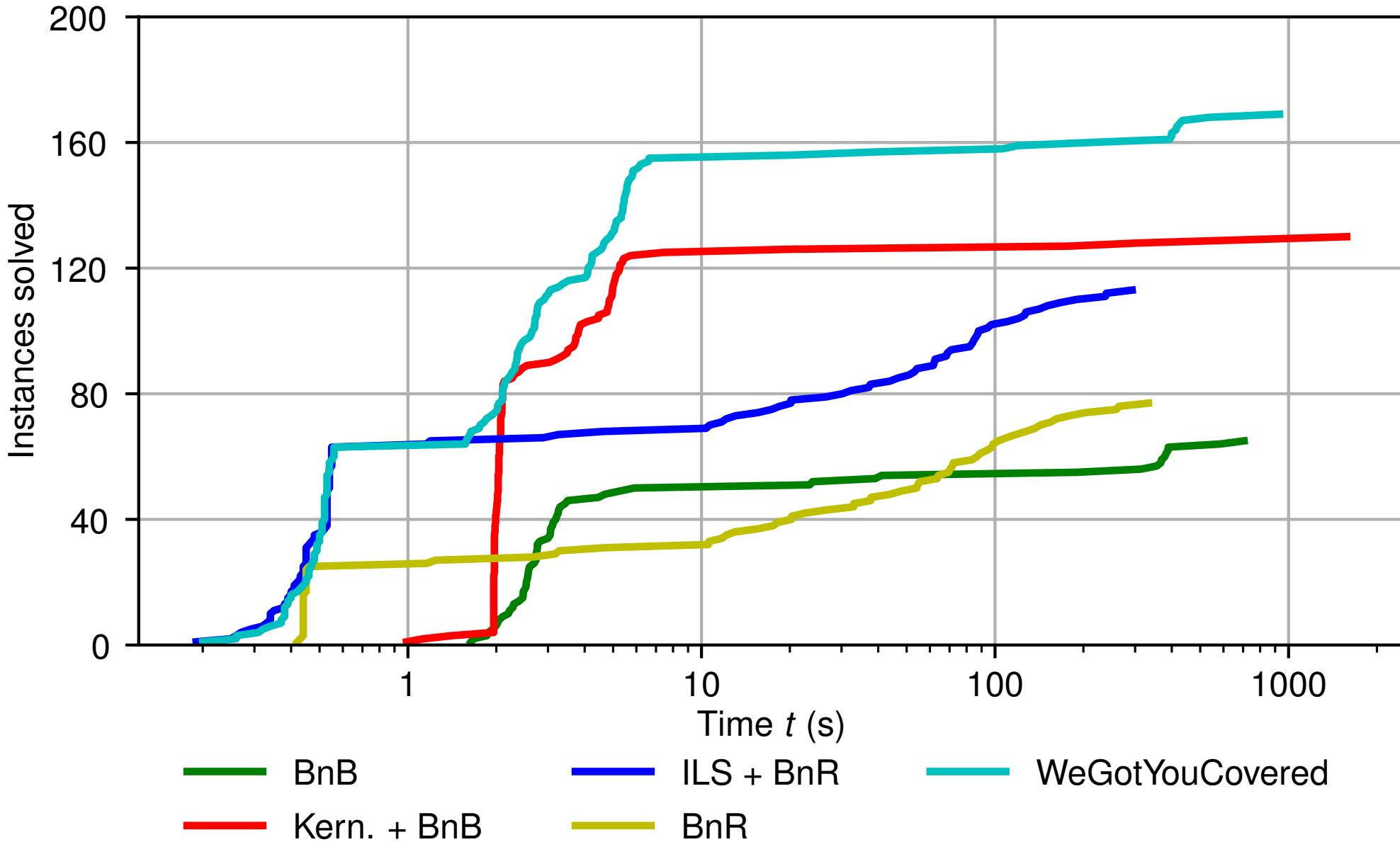
Algorithm Overview



Algorithm Overview



Instances Solved Over Time



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- [ML'14]
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